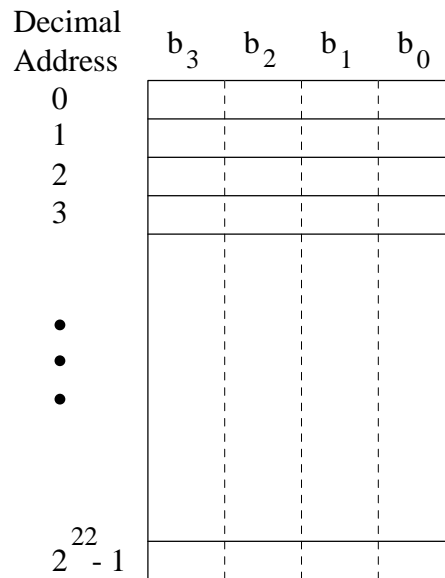


Implementation of Large Memory Chips

Consider a 4M x 4-bit chip which has a 22-bit addresses since $4M = 2^{22}$.

Logically/Externally, we view a 4M x 4-bit memory as pictured below with:

- each memory word made up of four bits: $b_3 b_2 b_1 b_0$



When we want to read a word, we supply a 22-bit address, and receive the corresponding 4-bit word.

The register-file implementation (see handout), does not scale well for large memories for several reasons:

- the number of gates in the address decoder (and MUXs) grows exponentially with the number of bits in the address.
- lots of wires into/out of the memory chip for address, data, and control

These problems are solved by

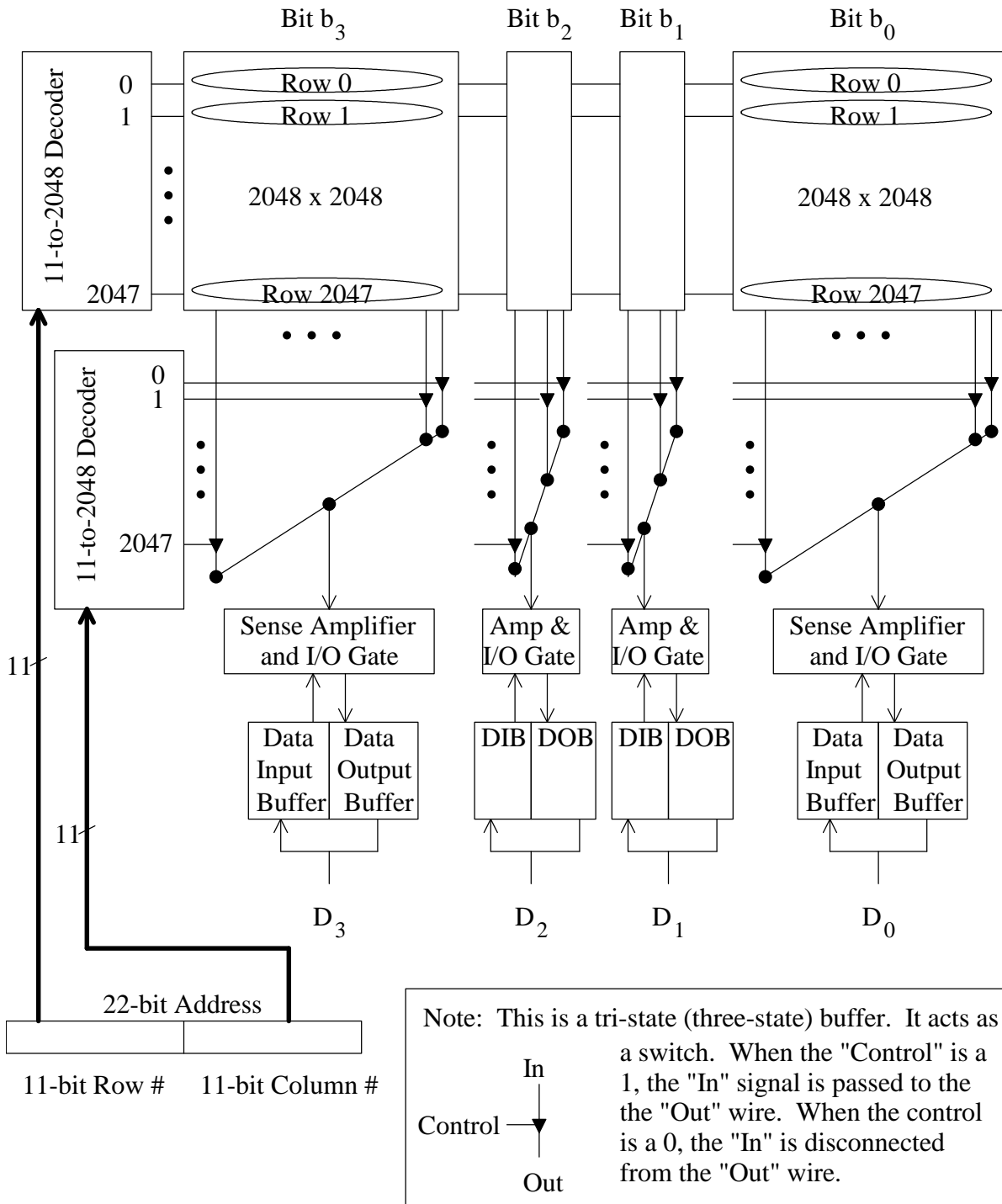
- using square-array of bits and decoding the address in two parts (row then column number)
- eliminate MUX's by using tri-state buffers
- single-port RAM memory - data wires shared for reading and writing

To help us see how the 4M x 4-bit memory gets mapped to the 2048 x 2048 x 4 memory array on the next page consider splitting memory into 2048 word blocks as shown below.

Decimal Address	Binary Address				Binary Address	
	b ₃	b ₂	b ₁	b ₀	Row #	Location in Row
0	Row 0	Row 0	Row 0	Row 0	0000000000	0000000000
1					0000000000	0000000001
2					0000000000	0000000010
3					0000000000	0000000011
•						
•						
•						
2047					0000000000	1111111111
2048	Row 1	Row 1	Row 1	Row 1	0000000001	0000000000
2049					0000000001	0000000001
•					0000000001	0000000010
•					0000000001	0000000011
•						
4095					0000000001	1111111111
4096	Row 2	Row 2	Row 2	Row 2	0000000010	0000000000
4097					0000000010	0000000001
4098					0000000010	0000000010
4099					0000000010	0000000011
•						
•						
•						
					0000000010	1111111111
					0000000011	0000000000
•						
•						
•						
	Row 2047	Row 2047	Row 2047	Row 2047	1111111111	0000000000
					1111111111	0000000001
					1111111111	0000000010
					1111111111	0000000011
$2^{22} - 1$					1111111111	1111111111

Each bit of a word is split into a separate 2048 x 2048 square memory “array”. Each of these square memory arrays is $2048 \times 2048 = 2^{11} \times 2^{11} = 2^{22} = 4\text{M}$ bits.

The 22-bit address of the 4M memory is split into two 11-bit parts. The upper 11-bits is first used to activate the correct row within the square memory arrays. Of the 2048-bit row that is read from each memory array, we are interested in only one bit. The lower 11-bits of the address specifies the location of the desired bit within the 2048-bit row.



Implementing Large Memory with Smaller Chips

Consider for example, implementing 4M x 32 bits memory with 256KB x 1 bit chips. The 256KB x 1 chips are implemented as square arrays of 512 x 512.

We will use an two-dimensional array of the 256KB x 1 chips to implement the larger memory.

The number of chips per row would be 32 bits / 1 bit = 32 chips.

The number of chips per column would be $4M / 256K = 2^{22} / 2^{18} = 2^4 = 16$ chips per column.

The 22-bit address of the 4M x 32 bit memory would be split up as follows:

