1) Draw the logic circuit using ANDs, ORs, and NOT gates for $\mathrm{F}=\bar{A} \bar{B} C+A \bar{B} \bar{C}+\bar{A} B C+A \bar{B} C$. What is the complexity (sum of \# inputs and \# gates)?
$\qquad$
Absent:

| Identity Name | AND Form | OR Form |
| :--- | :--- | :--- |
| Identity Law | $1 x=x$ | $0+x=x$ |
| Null (or Dominance) Law | $0 x=0$ | $1+x=1$ |
| Idempotent Law | $x x=x$ | $x+x=x$ |
| Inverse Law | $x \bar{x}=0$ | $x+\bar{x}=1$ |
| Commutative Law | $x y=y x$ | $x+y=y+x$ |
| Associative Law | $(x y) z=x(y z)$ | $(x+y)+z=x+(y+z)$ |
| Distributive Law | $x+y z=(x+y)(x+z)$ | $x(y+z)=x y+x z$ |
| Absorption Law | $x(x+y)=x$ | $x+x y=x$ |
| DeMorgan's Law | $(\overline{x y})=\bar{x}+\bar{y}$ | $(\overline{x+y})=\bar{x} \bar{y}$ |
| Double Complement Law | $\overline{\bar{x}}=x$ |  |

2) Using Boolean Algebra simplify $\mathrm{F}=\bar{A} \bar{B} C+A \bar{B} \bar{C}+\bar{A} B C+A \bar{B} C$.
3) Draw the simplified logic circuit using ANDs, ORs, and NOT. What is the complexity (sum of \# inputs and \# gates)?
4) Simplify the following using K-maps:
a) $F_{1}=\bar{A} \bar{B} \bar{C}+\bar{A} B+\bar{A} B C+A \bar{B} C+A \bar{B} \bar{C}$
b) $F_{2}=\bar{A} \bar{B} \bar{D}+A C \bar{D}+A B \bar{C} \bar{D}+A \bar{B} \bar{C} \bar{D}+A B C D$

Team \#:
Name: $\qquad$
Absent:
5) For the BCD to seven-segment display, what would the simplified SOP expression for the "c" segment? (Use "d" for don't cares)

| Decimal <br> Value | $\mathbf{x}_{\mathbf{1}}$ | $\mathbf{x}_{\mathbf{2}}$ | $\mathbf{x}_{\mathbf{3}}$ | $\mathbf{x}_{\mathbf{4}}$ | $\mathbf{a}$ | $\mathbf{b}$ | $\mathbf{c}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 1 | 1 |  |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 |  |
| 2 | 0 | 0 | 1 | 0 | 1 | 1 |  |
| 3 | 0 | 0 | 1 | 1 | 1 | 1 |  |
| 4 | 0 | 1 | 0 | 0 | 0 | 1 |  |
| 5 | 0 | 1 | 0 | 1 | 1 | 0 |  |
| 6 | 0 | 1 | 1 | 0 | 1 | 0 |  |
| 7 | 0 | 1 | 1 | 1 | 1 | 1 |  |
| 8 | 1 | 0 | 0 | 0 | 1 | 1 |  |
| 9 | 1 | 0 | 0 | 1 | 1 | 1 |  |
| 10 | 1 | 0 | 1 | 0 | d | d |  |
| 11 | 1 | 0 | 1 | 1 | d | d |  |
| 12 | 1 | 1 | 0 | 0 | d | d |  |
| 13 | 1 | 1 | 0 | 1 | d | d |  |
| 14 | 1 | 1 | 1 | 0 | d | d |  |
| 15 | 1 | 1 | 1 | 1 | d | d |  |


6) Since there are so many 1's in function c above, consider implementing $\overline{\mathrm{c}}$ and then negating it.

