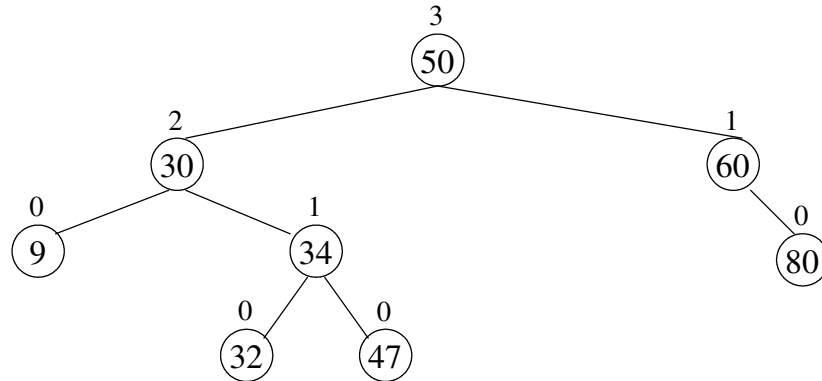1. An *AVL Tree* is a special type of Binary Search Tree (BST) that it is *height balanced*. By height balanced I mean that the height of every nodes left and right subtrees differ by at most one. This is enough to guarantee that a AVL tree with n nodes has a height no worst than $\Theta(\log_2 n)$. Therefore, insertions, deletions, and search are in the worst case $\Theta(\log_2 n)$. An example of an AVL tree with integer keys is shown below. The height of each node is shown.

```
                              3
                            (50)
               2                        1
             (30)                     (60)
      0              1                      0
     (9)           (34)                   (80)
               0         0
              (32)      (47)
```

Each AVL-tree node usually stores a *balance factor* in addition to its key and data. The balance factor keeps track of the relative height difference between its left and right subtrees.
a) Label each node in the above AVL tree with one of the following *balance factors*:
   • 'EQ' if its left and right subtrees are the same height
   • 'TL' if its left subtree is one taller than its right subtree
   • 'TR' if its right subtree is one taller than its left subtree

b) We start an `add` operation by adding the new item into the AVL as leaves just like we did for Binary Search Trees (BSTs).   Add the key 90 to the above tree?
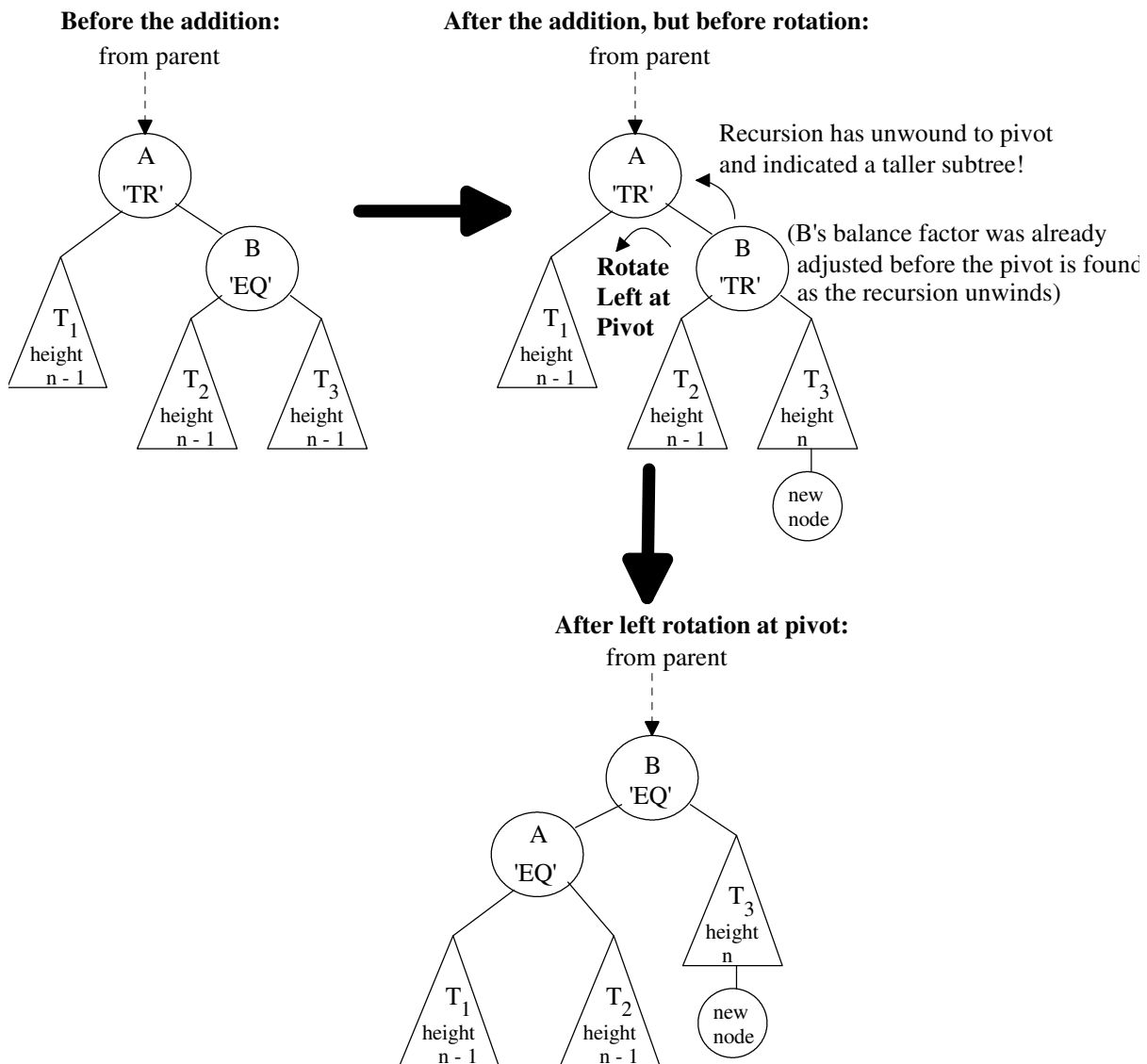
c) Identify the node "closest up the tree" from the inserted node (90) that no longer satisfies the height balanced property of an AVL tree. This node is called the *pivot node.* Label the pivot node above.

d) Consider the subtree whose root is the pivot node. How could we rearrange this subtree to restore the AVL height balanced property of AVL tree? (Draw the rearranged tree below)

2. Typically, the addition of a new key into an AVL requires the following steps:
- compare the new key with the current tree node's key (as we did in the `addHelper` function inside the `add` method in the BST) to determine whether to recursively add the new key into the left or right subtree
- add the new key as a leaf as the base case(s) to the recursion
- as the recursion "unwinds" (i.e., after you return from the recursive call) adjust the balance factors of the nodes on the search path from the new node back up to the root of the tree. To aid in adjusting the balance factors, we'll modify the `addHelper` function so that it returns `True` if the subtree got taller and `False` otherwise.
- as the recursion "unwinds" if we encounter a pivot node (as in question (c) above) we perform one or two "rotations" to restore the AVL tree's height-balanced property.
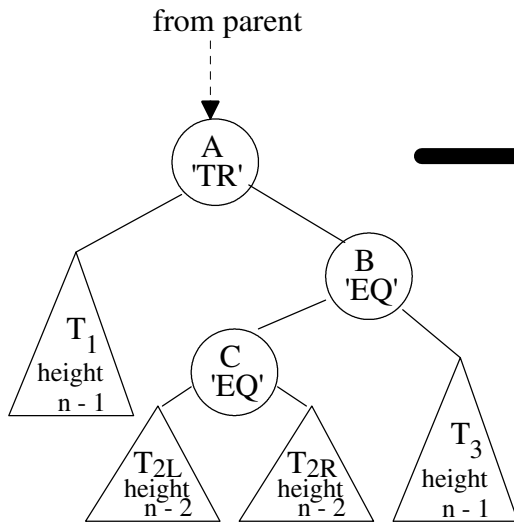
For example, consider the previous example of adding 90 to the AVL tree. Before the addition, the pivot node was already "TR" (tall right - right subtree had a height one greater than its left subtree). After inserting 90, the pivot's right subtree had a height 2 more than its left subtree which violates the AVL tree's height-balance property. This problem is handled with a *left rotation* about the pivot as shown in the following generalized diagram:
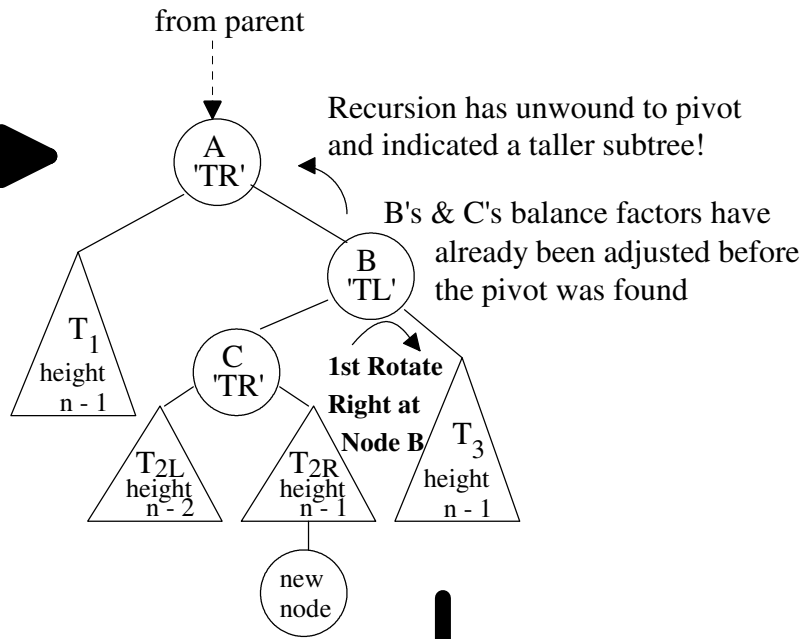

a) Assuming the same initial AVL tree (upper, left-hand of above diagram) if the new node would have increased the height of T_2 (instead of T_3), would a left rotation about the node A have rebalanced the AVL tree?

Before the addition, if the pivot node was already "TR" (tall right) and if the new node is inserted into the left subtree of the pivot node's right child, then we must do two rotations to restore the AVL-tree's height-balance property.
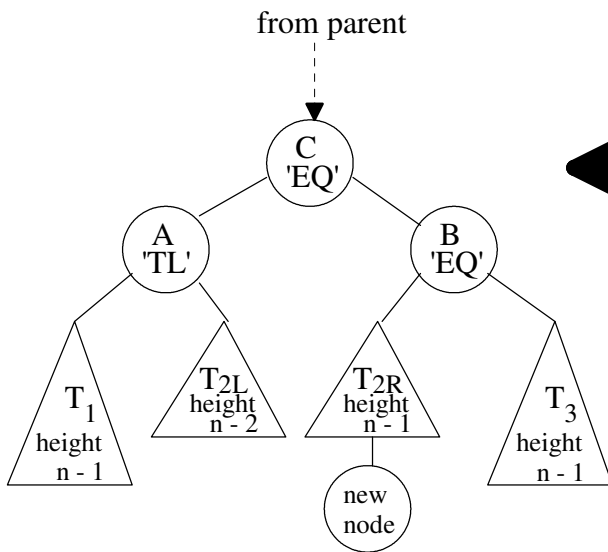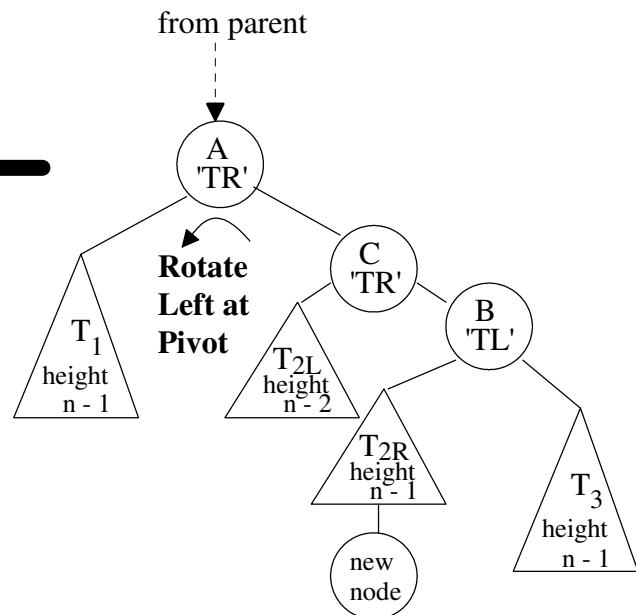
Before the addition:

from parent

After the addition, but before first rotation:

from parent

Recursion has unwound to pivot and indicated a taller subtree!

B's & C's balance factors have already been adjusted before the pivot was found

1st Rotate Right at Node B

After the left rotation at pivot and balance factors adjusted correctly:

from parent

After right rotation at B, but before left rotation at pivot:

from parent

Rotate Left at Pivot

b) Suppose that the new node was added into the right subtree of the pivot's right child, i.e., inserted in T_{2L} instead T_{2R}, then the same two rotations would restore the AVL-tree's height-balance property. However, what should the balance factors of nodes A, B, and C be after the rotations?

Data Structures (CS 1520)   Lecture 19   Name:_____

Consider the BinaryTreeAVL class that inherits and extends the BinaryTree class to include balance factors.

```
from binarytree import *

class BinaryTreeAVL(BinaryTree):
    def __init__(self, item, balance = 'EQ'):
        BinaryTree.__init__(self, item)
        self._balance = balance

    def getBalance(self):
        return self._balance

    def setBalance(self, newBalance):
        self._balance = newBalance

    def __str__(self):
        """Returns a string representation of the tree
        rotated 90 degrees to the left."""
        def strHelper(tree, level):
            result = ""
            if not tree.isEmpty():
                result += strHelper(tree.getRight(), level + 1)
                result += "| " * level
                result += str(tree.getRoot())+ " : " + tree.getBalance() + "\n"
                result += strHelper(tree.getLeft(), level + 1)
            return result
        return strHelper(self, 0)
```

Now let's consider the partial AVL class code that inherits from the BST class:

```
class AVL(BST):
    def __init__(self):
        BST.__init__(self)

    def add(self, newItem):
        """Adds newItem to the tree."""

        # Helper function to search for item's position
        def addHelper(tree):

            def rotateLeft(tree):
                newTree = BinaryTreeAVL(tree.getRoot())
                newTree.setLeft(tree.getLeft())
                newTree.setRight(tree.getRight().getLeft())
                newTree.setBalance(tree.getBalance())
                tree.setRoot(tree.getRight().getRoot())
                tree.setBalance(tree.getRight().getBalance())
                tree.setLeft(newTree)
                tree.setRight(tree.getRight().getRight())

            def rotateRight(tree):
                ### ADD CODE HERE FOR rotateRight ###
```

```
        # start of addHelper code
        currentItem = tree.getRoot()
        left = tree.getLeft()
        right = tree.getRight()

        if newItem < currentItem:
            ### CODE OMITTED HERE TO BE SUPPLIED BY YOU IN LAB

        else: # newItem > currentItem
            if not tree.getRight().isEmpty():
                tallerRightSubtree = addHelper(right)
                if tallerRightSubtree:
                    if tree.getBalance() == 'TL':
                        tree.setBalance('EQ')
                        return False
                    elif tree.getBalance() == 'EQ':
                        tree.setBalance('TR')
                        return True
                    else: # Two too tall on right now
                        if tree.getRight().getBalance() == 'TR': #only rotate left
                            tree.setBalance('EQ')
                            tree.getRight().setBalance('EQ')
                        else: # need to rotate right around right child, then
                              # rotate left on self
                            if tree.getRight().getLeft().getBalance() == 'TR':
                                tree.setBalance('TL')
                                tree.getRight().setBalance('EQ')
                            elif tree.getRight().getLeft().getBalance() == 'TL':
                                tree.setBalance('EQ')
                                tree.getRight().setBalance('TR')
                            else:
                                tree.setBalance('EQ')
                                tree.getRight().setBalance('EQ')

                            tree.getRight().getLeft().setBalance('EQ')
                            rotateRight(tree.getRight())
                        rotateLeft(tree)
                        return False
                else:
                    return False
            else:
                tree.setRight(BinaryTreeAVL(newItem))
                if tree.getBalance() == 'EQ':
                    tree.setBalance('TR')
                    return True
                else:
                    tree.setBalance('EQ')
                    return False
        # End of addHelper

    # Tree is empty, so new item goes at the root
    if self.isEmpty():
        self._tree = BinaryTreeAVL(newItem)
    # Otherwise, search for the item's spot
    else:
        addHelper(self._tree)
    self._size += 1
# end add
```
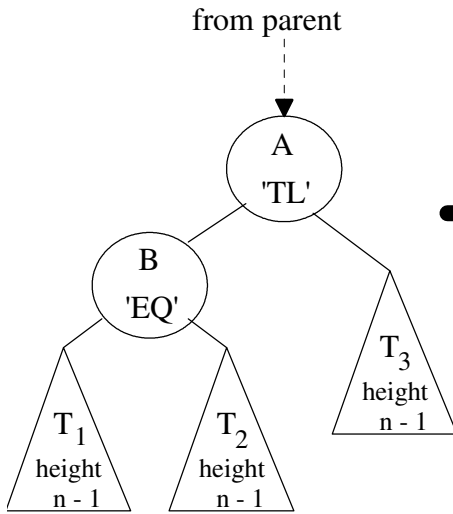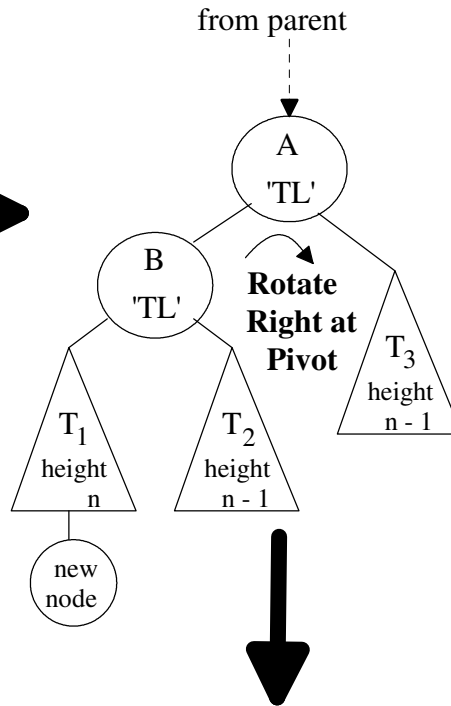
c) Where in the above code do the balance factors get set for your answer to question (b)?

d) Complete the below figure which is a "mirror image" to the figure on page 2, i.e., inserting into the pivot's left child's left subtree. Include correct balance factors after the rotation.

Before the insertion:

from parent

A
'TL'

B
'EQ'

T_1
height
n - 1

T_2
height
n - 1

T_3
height
n - 1

After the insertion, but before rotation:

from parent

A
'TL'

B
'TL'

Rotate Right at Pivot

T_1
height
n

T_2
height
n - 1

T_3
height
n - 1

new node

After right rotation at pivot:

3. a) Complete the below figure which is a "mirror image" to the figure on page 3, i.e., inserting into the pivot's left child's right subtree. Include correct balance factors after the rotation.
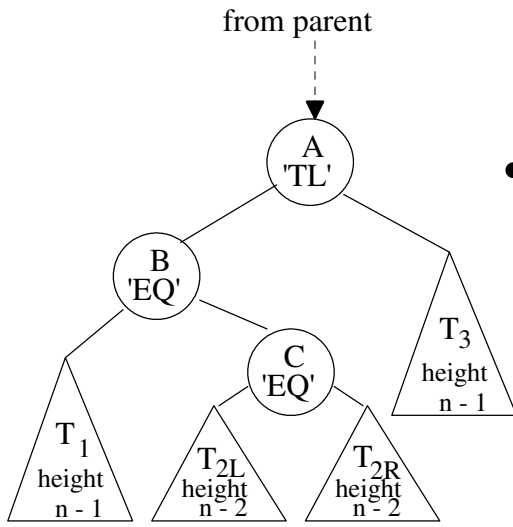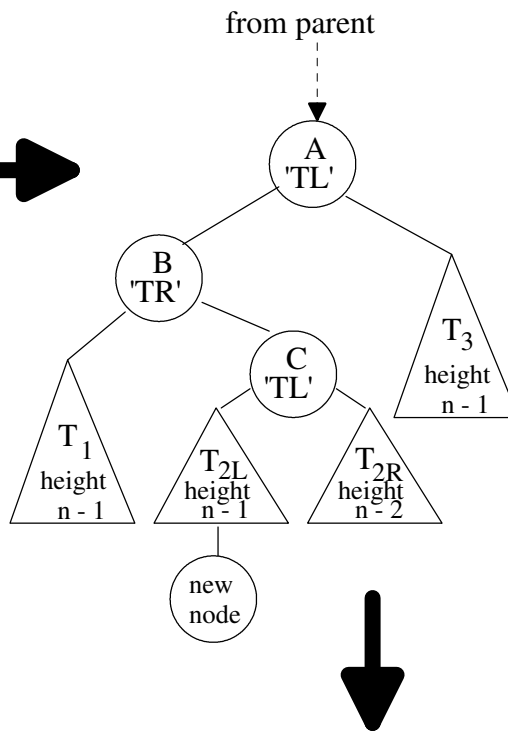
Before the insertion:　　　　　　　**After the insertion, but before first rotation:**

from parent　　　　　　　　　　　　　　from parent

A
'TL'

B
'EQ'

C
'EQ'

T_3
height
n - 1

T_1
height
n - 1

T_{2L}
height
n - 2

T_{2R}
height
n - 2

A
'TL'

B
'TR'

C
'TL'

T_3
height
n - 1

T_1
height
n - 1

T_{2L}
height
n - 1

T_{2R}
height
n - 2

new
node

After the right rotation at pivot and
balance factors adjusted correctly:

After left rotation at B, but
before right rotation at pivot:

In lab you will complete and test the `add` method code including rotateRight.