

Data Structures - Test 1

Question 1. (4 points) Consider the following Python code.

4
`for i in range(n*n): $-n^2$`
 `j = 1`
 `while j < n: $-\log_2 n$`
 `print (i, j)`
 `j = j * 2`

$$O(n^2 \log_2 n)$$

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 2. (4 points) Consider the following Python code.

4
`i = 1`
`while i < n: $-\log n$`
 `for j in range(n):`
 `print(j)`
 `for k in range(n):`
 `print(k)`
 `i = i * 2`

$$O(n^2 \log n)$$

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 3. (4 points) Consider the following Python code.

4
`def main(n):`
 `for i in range(n):`
 `doSomething(n)`
`def doSomething(n):`
 `for k in range(n):`
 `doMore(n)`
`def doMore(n):`
 `for k in range(n):`
 `print(k)`
`main(n)`

$$O(n^3)$$

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 4. (8 points) Suppose a $O(n^5)$ algorithm takes 10 seconds when $n = 100$. How long would you expect the algorithm to run when $n = 1,000$?

$$T(n) = c n^5$$

$$8 \quad T(100) = c 100^5 = 10 \text{ sec}$$

$$c = \frac{10}{100^5} \text{ sec} = 10^{-9} \text{ sec}$$

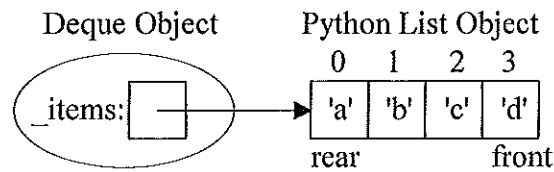
$$\begin{aligned} T(1000) &= c 1000^5 \\ &= 10^{-9} \text{ sec } 10^{15} \\ &= 10^6 \text{ sec} \end{aligned}$$

Question 5. (5 points) Why should you design a program instead of "jumping in" and start by writing code?

5 You will make fewer mistakes working from a design.

Question 6. A Deque (pronounced "Deck") is a linear data structure which behaves like a double-ended queue, i.e., it allows adding or removing items from either the front or the rear of the Deque. One possible implementation of a Deque would be to use a built-in Python list to store the Deque items such that

- the rear item is always stored at index 0,
- the front item is always at index $\text{len}(\text{self}._\text{items})-1$ or -1



a) (6 points) Complete the big-oh $O()$, for each Deque operation, assuming the above implementation. Let n be the number of items in the Deque.

| isEmpty | addRear | removeRear | addFront | removeFront | size |
|---------|---------|------------|----------|-------------|--------|
| $O(1)$ | $O(n)$ | $O(n)$ | $O(1)$ | $O(1)$ | $O(1)$ |

b) (9 points) Complete the method for the removeRear operation including the precondition check.

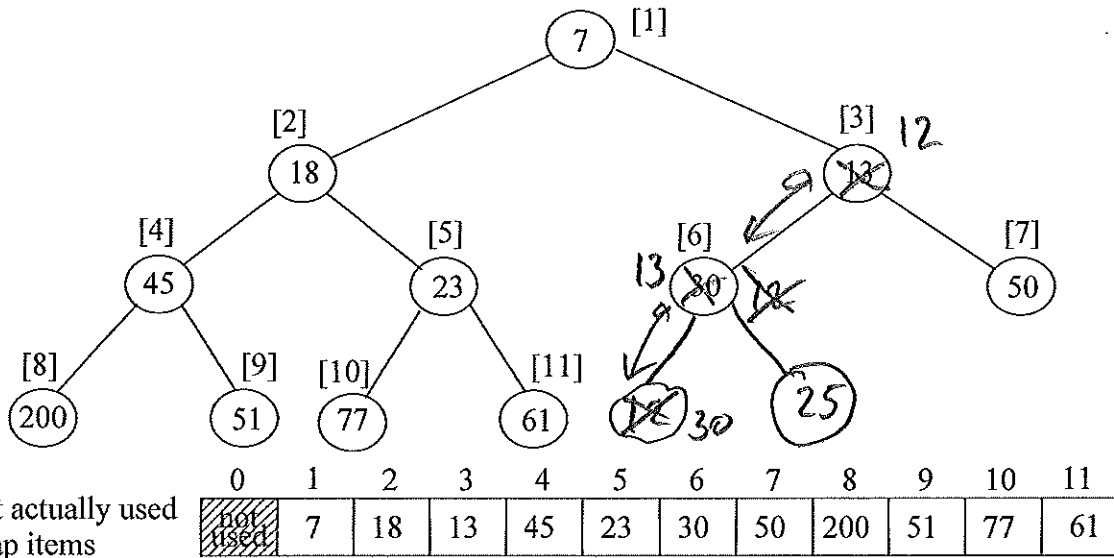
```
def removeRear(self):
    """Removes and returns the rear item of the Deque
    Precondition: the Deque is not empty.
    Postcondition: Rear item is removed from the Deque and returned"""
```

```
    if len(self._items) == 0:
        raise ValueError("Cannot remove from empty Deque")
    return self._items.pop(0)
```

c) (5 points) Suggest an alternate Deque implementation to speed up some of its operations.

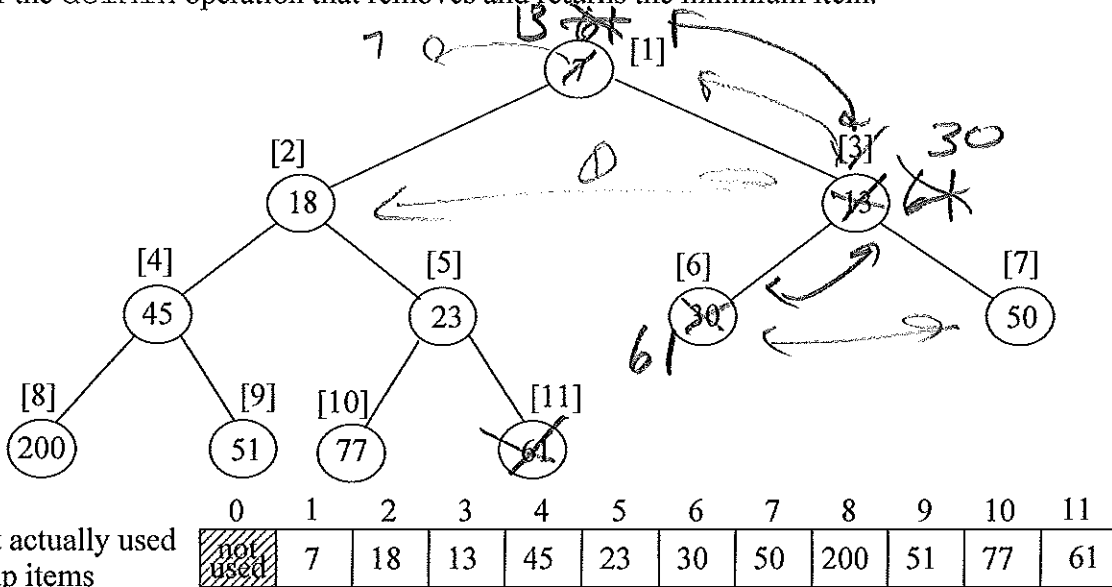
A doubly-linked list implementation.

Question 7. Consider the binary heap approach to implement a priority queue. A Python list is used to store a *complete binary tree* (a full tree with any additional leaves as far left as possible) with the items being arranged by *heap-order property*, i.e., each node is \leq either of its children. An example of a *min* heap “viewed” as a complete binary tree would be:



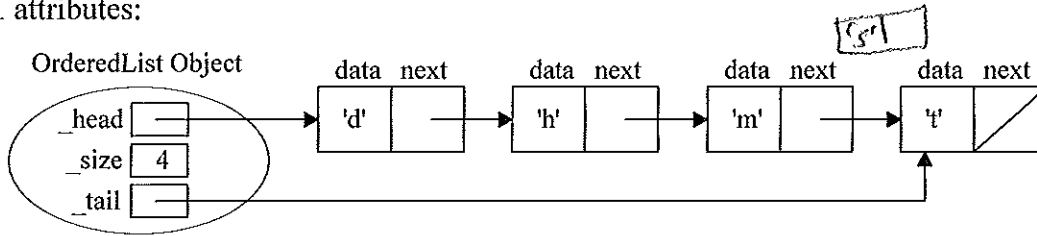
- a) (3 points) For the above heap, the list indexes are indicated in []'s. For a node at index i , what is the index of:
- its left child if it exists: $i * 2$
 - its right child if it exists: $i * 2 + 1$
 - its parent if it exists: $i // 2$
- b) (7 points) What would the above heap look like after inserting 12 and then 25 (show the changes on above tree)
- c) (3 points) What is the big-oh notation for inserting a new item in the heap? $O(\log_2 n)$

Now consider the `delMin` operation that removes and returns the minimum item.



- d) (2 point) What item would `delMin` remove and return from the above heap? 7
- e) (7 points) What would the above heap look like after `delMin`? (show the changes on above tree)
- f) (3 points) What is the big-oh notation for `delMin`? $O(\log_2 n)$

Question 8. The textbook's **Ordered list** ADT uses a singly-linked list implementation. I added the `_size` and `_tail` attributes:



a) (15 points) The `add(item)` method adds the add to the list. Recall that the textbook's implementation, cannot contain duplicate items!!! Thus, the precondition is that `item` is a not already in the list. Complete the `add(item)` method code including the precondition check.

```

class OrderedList(object):
    def __init__(self):
        self._head = None
        self._size = 0
        self._tail = None

    def add(self, item):
        temp = Node(item)
        prev = None
        current = self._head
        while current != None:
            if current.getData() == item:
                raise ValueError("cannot add duplicates")
            if current.getData() > item:
                break
            prev = current
            current = current.getNext()
        if prev == None:
            temp.setNext(self._head)
            self._head = temp
        else:
            temp.setNext(current)
            prev.setNext(temp)
            if current == None:
                self._tail = temp
            self._size += 1

class Node:
    def __init__(self, initdata):
        self.data = initdata
        self.next = None

    def getData(self):
        return self.data

    def getNext(self):
        return self.next

    def setData(self, newdata):
        self.data = newdata

    def setNext(self, newnext):
        self.next = newnext
    
```

b) (10 points) Assuming the ordered list ADT described above **does not allow duplicate items**. Complete the big-oh $O()$ for each operation. Let n be the number of items in the list.

| <code>add(item)</code> | <code>pop()</code> removes and returns tail item | <code>length()</code> returns number of items in the list | <code>remove(item)</code> removes the item from the list | <code>index(item)</code> returns the position of item in the list |
|------------------------|--|---|--|---|
| $O(n)$ | $O(n)$ | $O(1)$ | $O(n)$ | $O(n)$ |