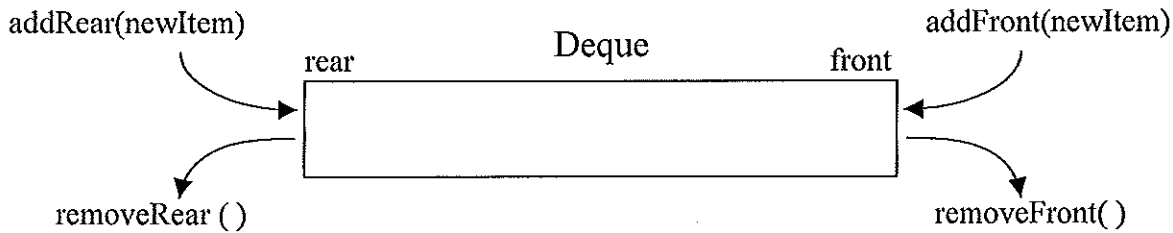
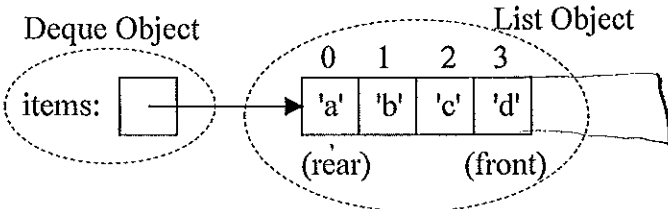


A Deque (pronounced "Deck") is a linear data structure which behaves like a double-ended queue, i.e., it allows adding or removing items from either the front or the rear of the Deque.



- One possible implementation of a Deque would be to use a Python list to store the Deque items such that
 - the rear item is always stored at index 0,
 - the front item is always stored at the highest index (or -1)



```
class Deque(object):
    def __init__(self):
        self.items = []
```

a) Complete the `__init__` method and determine the big-oh, $O()$, for each Deque operation, assuming the above implementation. Let n be the number of items in the Deque.

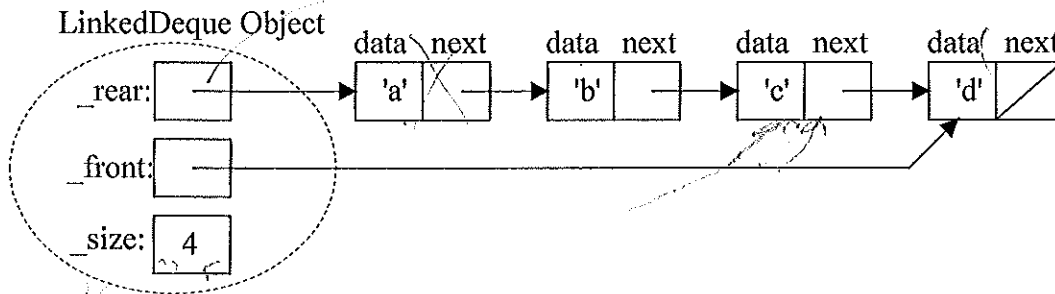
| isEmpty | addFront | removeFront | addRear | removeRear | size |
|---------|----------|-------------|---------|------------|--------|
| $O(1)$ | $O(1)$ | $O(1)$ | $O(n)$ | $O(n)$ | $O(1)$ |

b) Write the methods for the `addRear` and `removeRear` operation.

```
def addRear(self, newItem):
    self.items.insert(0, newItem)
```

```
def removeRear(self):
    if len(self.items) == 0:
        raise Exception("cannot...")
    return self.items.pop(0)
```

2. An alternative implementation of a Deque would be a linked implementation as in:

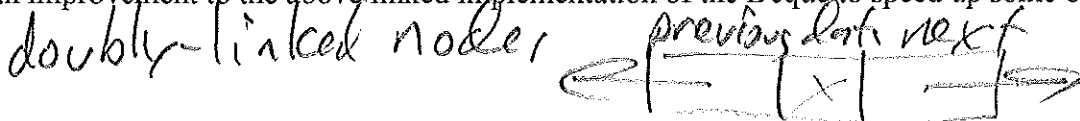


```
class LinkedDeque(object):
    def __init__(self):
        self._rear = None
        self._front = None
        self._size = 0
```

a) Complete the `__init__` method and determine the big-oh, $O()$, for each Deque operation assuming the above linked implementation. Let n be the number of items in the Deque.

| isEmpty | addFront | removeFront | addRear | removeRear | size |
|---------|----------|-------------|---------|------------|--------|
| $O(1)$ | $O(1)$ | $O(n)$ | $O(1)$ | $O(1)$ | $O(1)$ |

b) Suggest an improvement to the above linked implementation of the Deque to speed up some of its operations.



```

from node import Node

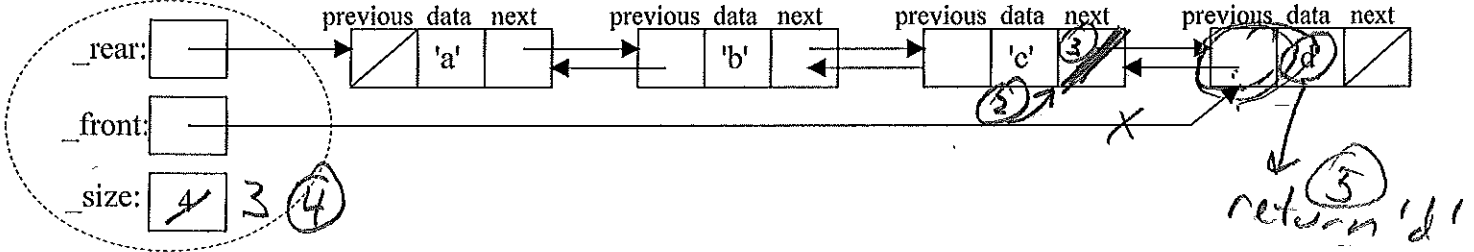
class Node2Way(Node):
    def __init__(self, initdata):
        Node.__init__(self, initdata)
        self.previous = None

    def getPrevious(self):
        return self.previous

    def setPrevious(self, newprevious):
        self.previous = newprevious
    
```

3. An alternative implementation of a Deque would be a doubly-linked implementation as in:

DoublyLinkedList Object

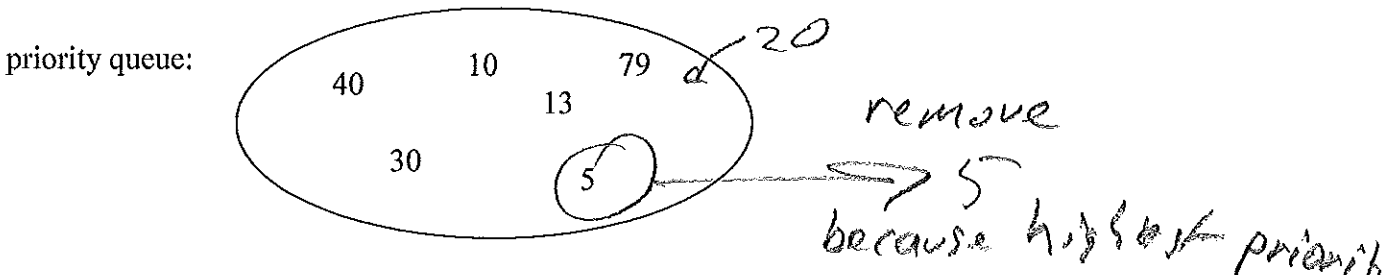


a) Determine the big-oh, $O()$, for each Deque operation assuming the above doubly-linked implementation. Let n be the number of items in the Deque.

| isEmpty | addFront | removeFront | addRear | removeRear | size |
|---------|----------|-------------|---------|------------|--------|
| $O(1)$ | $O(1)$ | $O(1)$ | $O(1)$ | $O(1)$ | $O(1)$ |

4. A *priority queue* has the same operations as a regular queue, except the items are NOT returned in the FIFO (first-in, first-out) order. Instead, each item has a priority that determines the order they are removed. A hospital emergency room operates like a priority queue -- the person with the most serious injure has highest priority even if they just arrived.

a) Suppose that we have a priority queue with integer priorities such that the smallest integer corresponds to the highest priority. For the following priority queue, which item would be dequeued next?



b) To implement a priority queue, we could use an **unordered Python list**. If we did, what would be the big-oh notation for each of the following methods: (justify your answer)

- enqueue: $O(1)$ - append to right end of list
- dequeue: $O(n)$

c) To implement a priority queue, we could use a **Python list order-by priorities in decending order**. If we did, what would be the big-oh notation for each of the following methods: (justify your answer)

- enqueue: $O(n)$
 - dequeue: $O(1)$
- Handwritten notes: $O(n)$ search, $O(n)$ enqueue, overall $O(n)$, $O(n^2)$ insert.

Normal-case remove Front Deque if self._size == 0:
raise Exception('')

① temp = self._front

② self._front = self._front.getPrevious()

③ self._front.setNext(None) if self._size == 1:
self._rear = None
else:

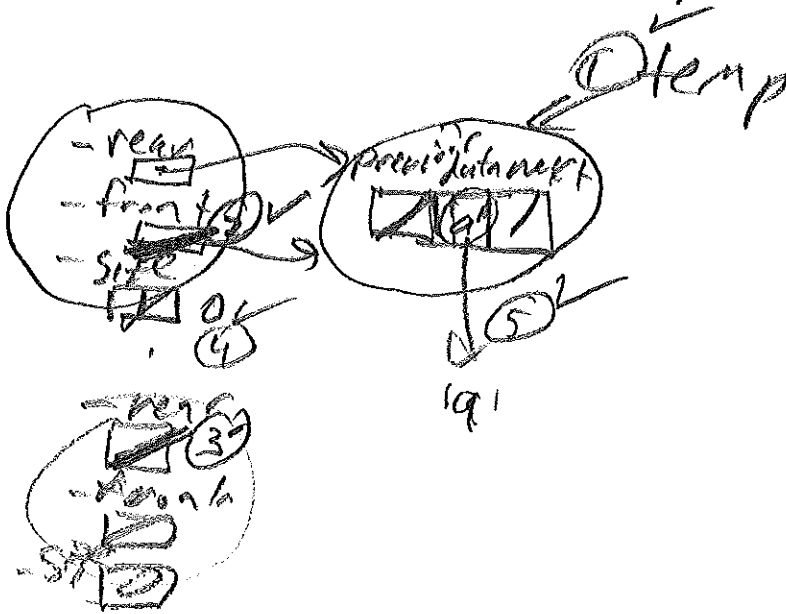
④ self._size -= 1

⑤ return temp.getData()

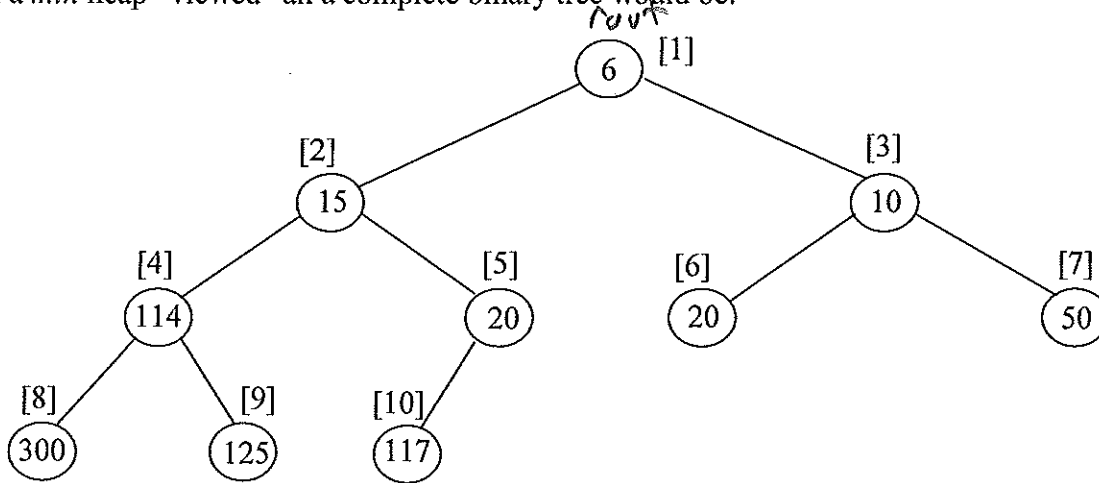
Special cases

(1) removing Front of only / last node

(2) precondition - empty Deque



1. Section 6.6 discusses a very “non-intuitive”, but powerful list/array-based approach to implement a priority queue, call a binary heap. The list/array is used to store a *complete binary tree* (a full tree with any additional leaves as far left as possible) with the items being arranged by *heap-order property*, i.e., each node is \leq either of its children. An example of a *min heap* “viewed” as a complete binary tree would be:



Python List actually used to store heap items

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|----------|---|----|----|-----|----|----|----|-----|-----|-----|
| not used | 6 | 15 | 10 | 114 | 20 | 20 | 50 | 300 | 125 | 117 |

a) For the above heap, the list/array indexes are indicated in []'s. For a node at index i , what is the index of:

- its left child if it exists: $2 \times i$
- its right child if it exists: $2 \times i + 1$
- its parent if it exists: $i // 2$

b) What would the above heap look like after inserting 13 and then 3? (show the changes on above tree)

General Idea of `insert(newItem)`:

- append `newItem` to the end of the list (easy to do, but violates heap-order property)
- restore the heap-order property by repeatedly swapping the `newItem` with its parent until it *percolates* to correct spot

c) What is the big-oh notation for inserting a new item in the heap?

d) Complete the code for the `percUp` method used by `insert`.

```
class BinHeap:
    def __init__(self):
        self.heapList = [0]
        self.currentSize = 0

    def percUp(self, currentIndex):
        parentIndex =
        while

    def insert(self, k):
        self.heapList.append(k)
        self.currentSize = self.currentSize + 1
        self.percUp(self.currentSize)
```