

Question 1. (4 points) Consider the following Python code.

```

4 for i in range(n):
    j = 1
    while j < n:
        print(i, j)
        j = j + 2
    
```

$O(n^2)$

(+2 for $n \log_2 n$)

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 2. (4 points) Consider the following Python code.

```

4 i = 1
while i < n:
    for j in range(n):
        print(j)
    for k in range(n):
        print(k)
    i = i * 2
    
```

$2n \log_2 n \Rightarrow O(n \log_2 n)$

(+2 for $O(n^2 \log_2 n)$)

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 3. (4 points) Consider the following Python code.

```

4 def main(n):
    for i in range(n):
        doSomething(n)
        doMore(n)
    def doSomething(n):
        for k in range(2**n):
            print(k)
    def doMore(n):
        for k in range(n):
            print(k)
    main(n)
    
```

$O(n 2^n)$

(+3 for $n 2^n$)
(+3 for 2^n)

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 4. (8 points) Suppose a $O(n^4)$ algorithm takes 1 second when $n = 100$. How long would you expect the algorithm to run when $n = 1,000$?

$T(n) = C n^4$

$T(1000) = C 1000^4 = C 10^{12}$

8 $T(100) = C 100^4 = 1 \text{ sec}$

$= \frac{1}{10^8} 10^{12} \text{ sec}$

$C = \frac{1}{100^4} = \frac{1}{10^8} \text{ sec}$

$= 10^4 \text{ sec}$

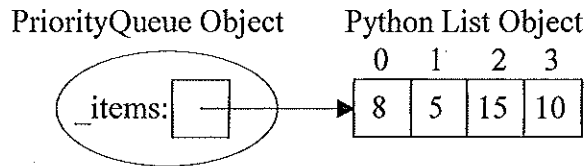
$= 10,000 \text{ sec}$

Question 5. (5 points) In lab 2 (and on the Python Summary) the AdvancedDie class inherited from the Die class. How does inheritance aid a programmer in writing code?

5 The subclass inherits working/correct code which it does not need to duplicate

Question 6. A **priority queue** has the same operations as a regular queue, except the items are NOT returned in the FIFO (first-in, first-out) order. Instead, each item has a priority that determines the order they are removed. One possible implementation of a priority queue would be to use a built-in Python list to store the items such that

- items in the Python list are **unordered** by their priorities,
- lowest number indicates the highest priority (i.e., dequeuing from the below priority queue would return 5)



a) (5 points) Complete the big-oh $O()$, for each PriorityQueue operation, assuming the above implementation. Let n be the number of items in the PriorityQueue.

	isEmpty	enqueue(item)	dequeue	<u>str</u>	size
5	$O(1)$	$O(1)$	$O(n)$	$O(n)$	$O(1)$

b) (15 points) Complete the method for the dequeue operation including the precondition check.

```

class PriorityQueue(object):
    def __init__(self):
        self._items = []

    def dequeue(self):
        """Removes and returns the highest priority (lowest value) item in the
        PriorityQueue

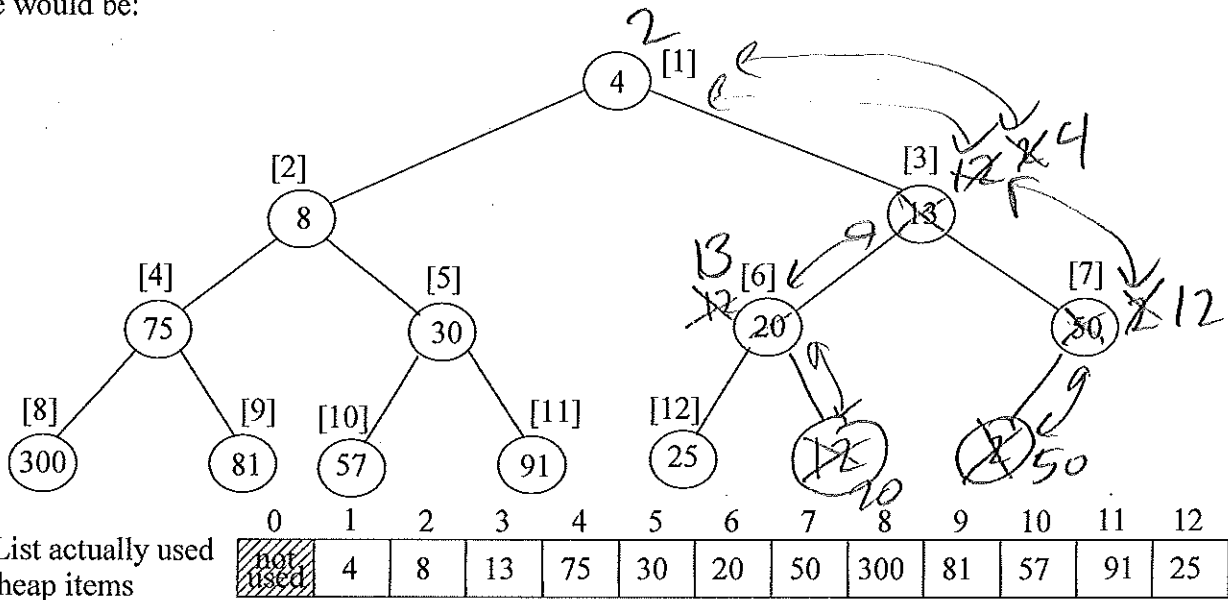
        Precondition: the PriorityQueue is not empty.
        Postcondition: the highest priority (lowest value) item in the PriorityQueue is
        removed and returned"""

        5 { if len(self._items) == 0:
            raise ValueError("Cannot dequeue from empty priority queue")
        }
        5 { minIndex = 0
            for test in range(1, len(self._items)):
                if self._items[test] < self._items[minIndex]:
                    minIndex = test
            }
        5 { return self._items.pop(minIndex)
    }
    
```

c) (5 points) Suggest an alternate PriorityQueue implementation to speed up some of its operations.

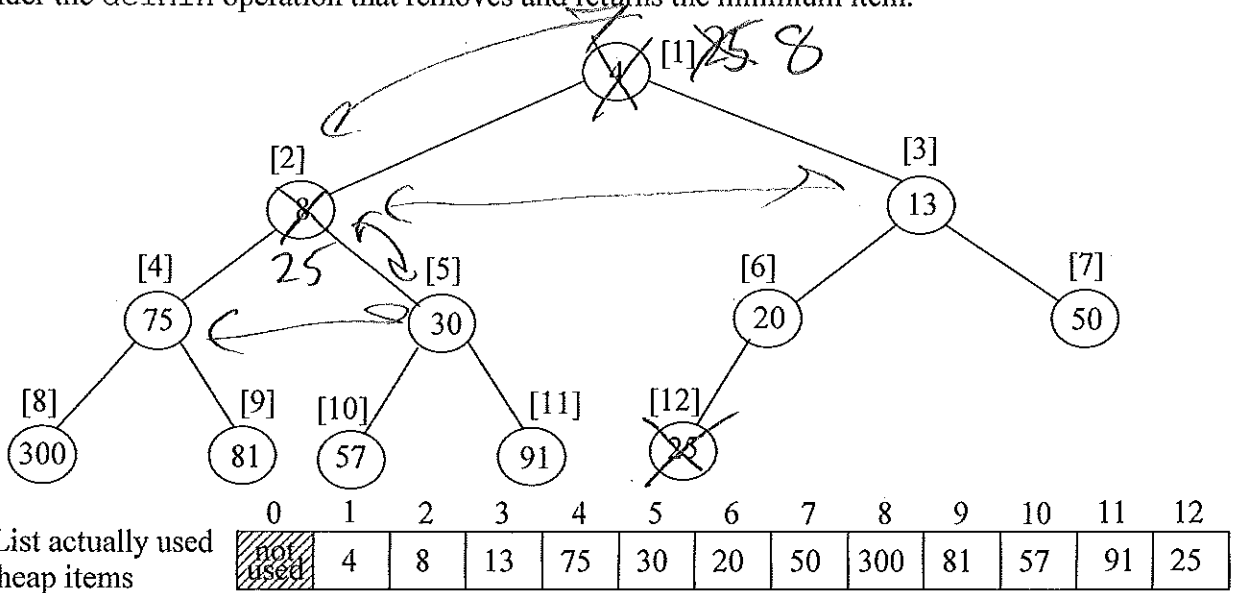
5 Use a binary heap

Question 7. Consider the binary heap approach to implement a priority queue. A Python list is used to store a *complete binary tree* (a full tree with any additional leaves as far left as possible) with the items being arranged by *heap-order property*, i.e., each node is \leq either of its children. An example of a *min heap* "viewed" as a complete binary tree would be:



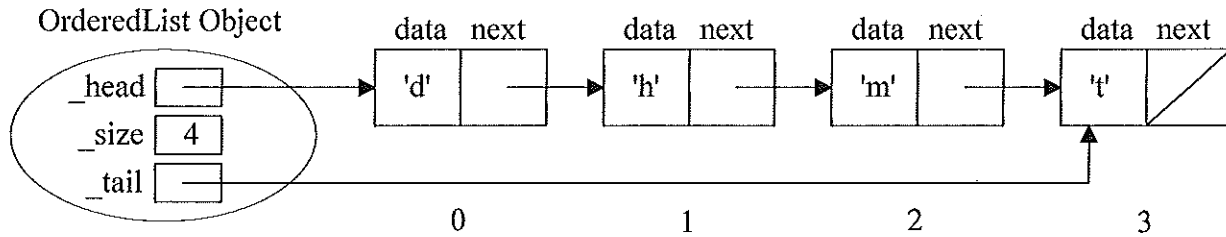
- a) (3 points) For the above heap, the list indexes are indicated in []'s. For a node at index i , what is the index of:
 - its left child if it exists: $i \times 2$
 - its right child if it exists: $i \times 2 + 1$
 - its parent if it exists: $i // 2$
- b) (7 points) What would the above heap look like after inserting 12 and then 2 (show the changes on above tree)
- c) (3 points) What is the big-oh notation for inserting a new item in the heap? $O(\log_2 n)$

Now consider the `delMin` operation that removes and returns the minimum item.



- d) (2 point) What item would `delMin` remove and return from the above heap? 4
- e) (7 points) What would the above heap look like after `delMin`? (show the changes on above tree)
- f) (3 points) Why does a `delMin` operation typically take longer than an `insert` operation?
 - (1) grabbing last item in list to put at root is typically larger and needs to percolate down further than inserted item goes up.
 - (2) to move down one level requires 2 comparisons v.s. 1 compare to move up

Question 8. The textbook's **Ordered list** ADT uses a singly-linked list implementation. I added the `_size` and `_tail` attributes:



a) (15 points) The `index(item)` method returns the position of the `item` in the list (e.g., 'm' is at position 2). Recall that the textbook's implementation, assumes the `item` is in the list!!! Thus, the precondition is that `item` is in the list. Complete the `index(item)` method code including the precondition check.

```

class OrderedList(object):
    def __init__(self):
        self._head = None
        self._size = 0
        self._tail = None
    def index(self, item):
        position = 0
        current = self._head
        while True:
            if current == None or current.getData() > item:
                raise ValueError("item not in list has not index")
            elif current.getData() == item:
                return position
            else:
                current = current.getNext()
                position += 1

```

```

class Node:
    def __init__(self, initdata):
        self.data = initdata
        self.next = None
    def getData(self):
        return self.data
    def getNext(self):
        return self.next
    def setData(self, newdata):
        self.data = newdata
    def setNext(self, newnext):
        self.next = newnext

```

b) (10 points) Assuming the ordered list ADT described above **does not allow duplicate items**. Complete the big-oh $O()$ for each operation. Let n be the number of items in the list.

<code>add(item)</code> adds the item into the list	<code>pop()</code> removes and returns tail item	<code>length()</code> returns number of items in the list	<code>remove(item)</code> removes the item from the list	<code>index(item)</code> returns the position of item in the list
$O(n)$	$O(n)$	$O(1)$	$O(n)$	$O(n)$