

Data Structures - Test 1

Question 1. (5 points) Consider the following Python code.

```

5 for i in range(n): -n
    j = 1
    while j < n: -log2 n
        for k in range(n): ← n
            print (i, j, k)
        j = j * 2
    
```

$O(n^2 \log_2 n)$

(+3 $n \log_2 n$)

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 2. (5 points) Consider the following Python code.

```

5 i = 2**n # this is 2^n
  while i > 1: ← i's values: 2^n, 2^{n-1}, 2^{n-2}, ..., 2^2, 2^1 ← loops n times
    for j in range(n): ← n
        print(j)
    i = i // 2
  
```

$O(n^2)$
(+3 $n \log_2 n$)

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 3. (5 points) Consider the following Python code.

```

5 def main(n):
    for i in range(n): -n
        doSomething(n)
  def doSomething(n):
    for k in range(n): -n
        print(k)
  main(n)
  
```

$O(n^2)$

What is the big-oh notation $O()$ for this code segment in terms of n ?

Question 4. (10 points) Suppose a $O(n^4)$ algorithm takes 10 seconds when $n = 1,000$. How long would you expect the algorithm to run when $n = 10,000$?

$T(n) = c n^4$ $T(1000) = c 1000^4 = 10 \text{ sec}$ $c = \frac{10 \text{ sec}}{1000^4} = 10^{-11} \text{ sec}$

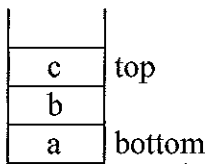
10 $T(10,000) = c 10000^4 = 10^{-11} \text{ sec} 10000^4 = 10^{-11} \text{ sec} \times 10^{16} = 10^5 \text{ sec}$
 $= 100,000 \text{ sec}$
 $= 27.78 \text{ hrs}$
 (+2 100 sec)

Question 5. (10 points) Why should you design a program instead of "jumping in" and start writing code?

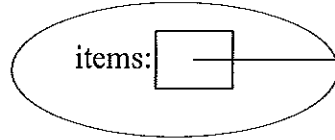
10 By designing a program you can split the work into smaller, managable pieces. It is easier to implement/code, and debug/test the smaller pieces.

Question 6. Consider the following Stack implementation utilizing a Python list:

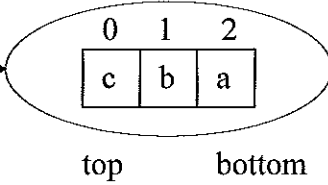
"Abstract"
Stack



Stack Object



Python list Object



a) Complete the big-oh notation for the ~~alternate~~ Stack methods: ("n" is the # items)

	push(item)	pop()	peek()	size()	isEmpty()	__init__
Big-oh	$O(n)$	$O(n)$	$O(1)$	$O(1)$	$O(1)$	$O(1)$

b) (9 points) Complete the method for the pop operation including the precondition check.

class Stack:

```
def __init__(self):
    self._items = []
```

def pop(self):

```
    """Removes and returns the top item of the stack
    Precondition: the stack is not empty.
    Postcondition: the top item is removed from the stack and returned"""
```

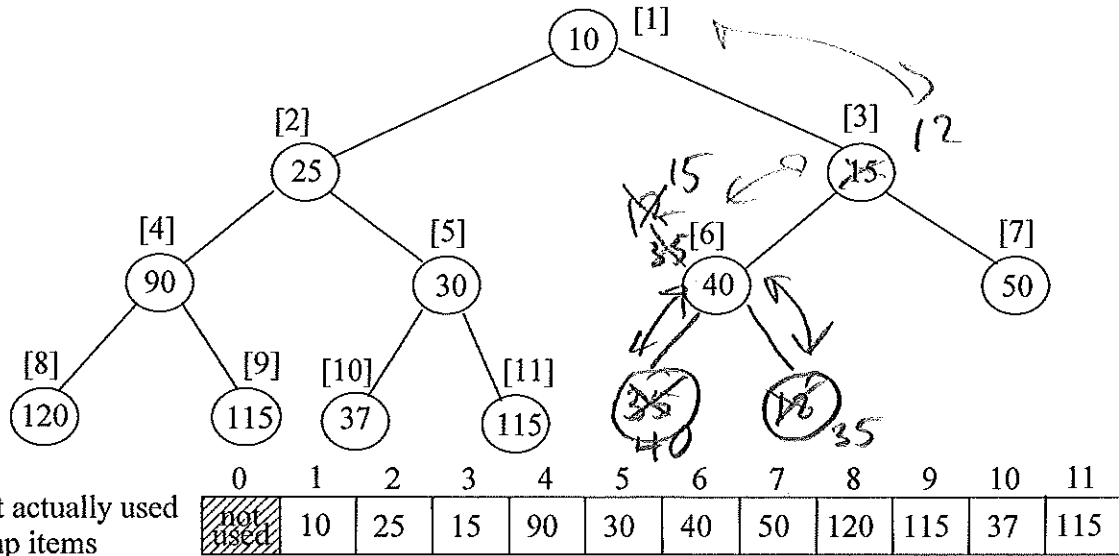
```
    if len(self.items) == 0:
        raise ValueError("cannot pop from empty stack")
    return self.items.pop(0)
```

c) (5 points) Suggest an alternate Stack implementation to speed up some of its operations.

Reverse the top and bottom, i.e., let index 0 hold the bottom item of the stack

Question 7. Consider the binary heap approach to implement a priority queue. A Python list is used to store a *complete binary tree* (a full tree with any additional leaves as far left as possible) with the items being arranged by *heap-order property*, i.e., each node is \leq either of its children. An example of a *min heap* "viewed" as a complete binary tree would be:

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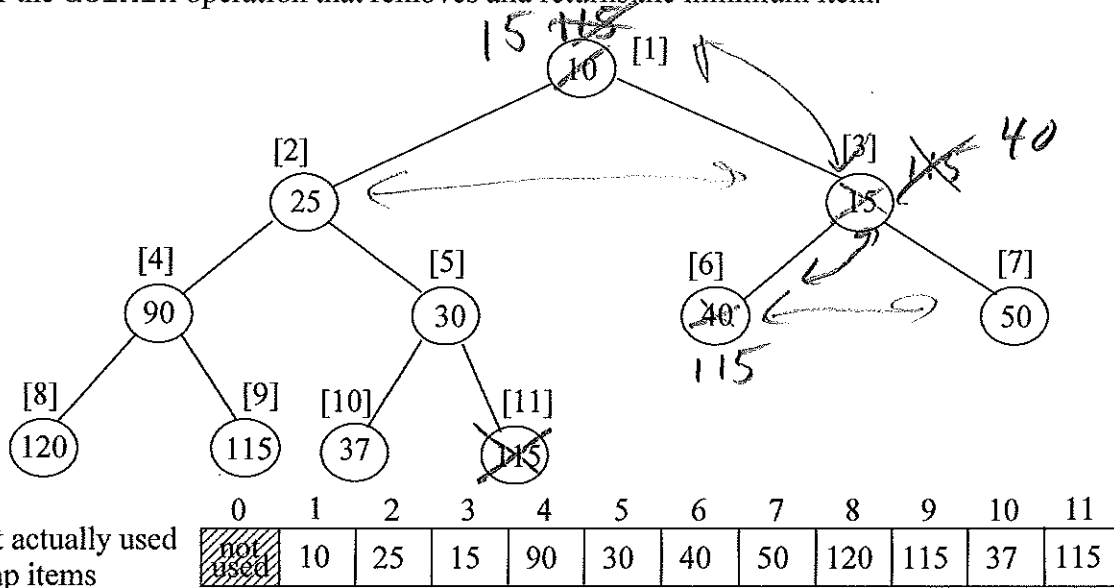
a) (3 points) For the above heap, the list indexes are indicated in []'s. For a node at index i , what is the index of:

- its left child if it exists: $i * 2$
- its right child if it exists: $i * 2 + 1$
- its parent if it exists: $i // 2$

b) (6 points) What would the above heap look like after inserting 35 and then 12 (show the changes on above tree)

c) (2 points) What is the big-oh notation for inserting a new item in the heap? $O(\log n)$

Now consider the `delMin` operation that removes and returns the minimum item.

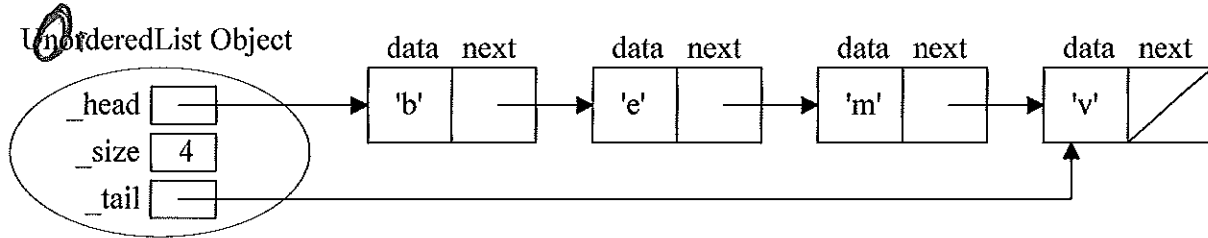


d) (1 point) What item would `delMin` remove and return from the above heap? 10

e) (6 points) What would the above heap look like after `delMin`? (show the changes on above tree)

f) (2 points) What is the big-oh notation for `delMin`? $O(\log n)$

Question 8. The textbook's ordered list ADT uses a singly-linked list implementation. I added the `_size` and `_tail` attributes:



a) (15 points) The `pop(position)` method removes and returns the item at the specified position. The precondition is that `position` is a nonnegative integer corresponding to an actual list item (e.g., for the above list $0 \leq \text{position} \leq 3$). Complete the `pop(position)` method code including the precondition check.

```
class OrderedList:
    def __init__(self):
        self._head = None
        self._size = 0
        self._tail = None

    def pop(self, position):
        if not isinstance(position, int):
            raise TypeError("Pop position must be an integer")
        if position < 0 or position >= self._size:
            raise IndexError("Pop position outside value range")
        previous = None
        current = self._head
        for count in range(position):
            previous = current
            current = current.getNext()
        if previous == None:
            self._head = self._head.getNext()
        else:
            previous.setNext(current.getNext())
        if current == self._tail:
            self._tail = previous
        self._size -= 1
        return current.getData()
```

b) (10 points) Assuming the ordered list ADT described above. Complete the big-oh $O()$ for each operation. Let n be the number of items in the list.

<code>pop(position)</code> removes and returns the item at the specified position	<code>pop()</code> removes and returns tail item	<code>length()</code> returns number of items in the list	<code>index(item)</code> returns the position of item in the list	<code>add(item)</code> adds item to its sorted spot in the list
$O(n)$	$O(n)$	$O(1)$	$O(n)$	$O(n)$

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