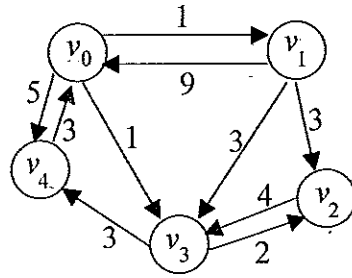


1. Consider the following directed graph (digraph)  $G = (V, E)$ :



- a) What is the set of vertices?  $V = \{v_0, v_1, v_2, v_3, v_4\}$
- b) An edge can be represented by a tuple (from vertex, to vertex [, weight]). What is the set of edges?  
 $E = \{(v_0, v_1, 1), (v_0, v_3, 1), (v_0, v_4, 5), (v_1, v_0, 9), (v_1, v_2, 3), (v_1, v_3, 3), (v_2, v_3, 4), (v_3, v_2, 2), (v_0, v_4, 5), (v_4, v_0, 3), (v_3, v_4, 3)\}$
- c) A path is a sequence of vertices that are connected by edges. In the graph G above, list two different paths from  $v_0$  to  $v_3$ .  $v_0, v_3$  or  $v_0, v_1, v_2, v_3$
- d) A cycle in a directed graph is a path that starts and ends at the same vertex. Find a cycle in the above graph.  $v_0, v_3, v_4, v_0$  or  $v_0, v_1, v_0$

2. Like most data structures, a graph can be represented using an array, or as a linked list of nodes. The array representation is a two-dimensional array (called an *adjacency matrix*) whose elements contain information about the edges and the vertices corresponding to the indices. (Python could use a list-of-lists)

a) Complete the following adjacency matrix for the above graph. (Here a missing edge is represented by  $\infty$ .)

Storage  $O(V^2)$  (from vertex)

Lookup edge existence or weight  $O(1)$

adj Matrix [#row] [ ] (to vertex)

	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$
$v_0$	0	1	$\infty$	1	5
$v_1$	9	0	3	3	$\infty$
$v_2$	$\infty$	$\infty$	0	4	$\infty$
$v_3$	$\infty$	$\infty$	2	0	3
$v_4$	3	$\infty$	$\infty$	$\infty$	0

b) The linked representation maintains a linked-list (or Python dictionary) of vertices with each vertex maintaining a linked list of other vertices that it connects to. Complete the adjacency list representation below:

Storage  $O(V+E)$

Lookup edge or weight  $O(1)$

Python dictionary  $O(1)$

dictionary edge info  $O(1)$

Vertex List:  $v_0, v_1, v_2, v_3, v_4$

Edge Lists for each vertex containing connected-to and edge-weight information:

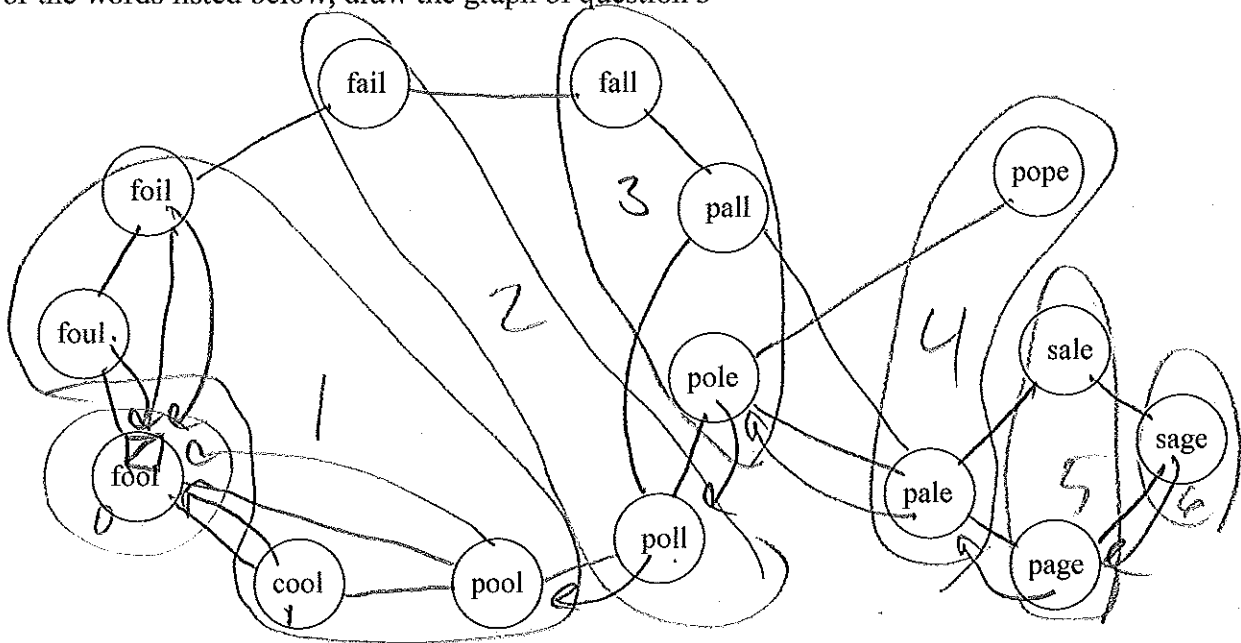
- $v_0$ :  $v_1$  (1),  $v_3$  (1),  $v_4$  (5)
- $v_1$ :  $v_0$  (9),  $v_2$  (3),  $v_3$  (3)
- $v_2$ :  $v_3$  (4)
- $v_3$ :  $v_2$  (2),  $v_4$  (3)
- $v_4$ :  $v_0$  (3)

3. Graphs can be used to solve many problems by modeling the problem as a graph and using "known" graph algorithm(s). For example, consider the *word-ladder puzzle* where you transform one word into another by changing one letter at a time, e.g., transform FOOL into SAGE by FOOL → FOIL → FAIL → FALL → PALL → PALE → SALE → SAGE.

We can use a graph algorithm to solve this problem by constructing a graph such that

- a word represents a vertex
- an edge represents?
  - a word ladder transformation from one word to another represents?

4. For the words listed below, draw the graph of question 3



a) List a different transformation from FOOL to SAGE

*fool → pool → poll → pole → pale → page → sage*

b) If we wanted to find the shortest transformation from FOOL to SAGE, what does that represent in the graph?

*shortest path from fool to sage*

c) There are two general approaches for traversing a graph from some starting vertex  $s$ :

- Breadth First Search (BFS) where you find all vertices a distance 1 (directly connected) from  $s$ , before finding all vertices a distance 2 from  $s$ , etc.
- Depth First Search (DFS) where you explore as deeply into the graph as possible. If you reach a "dead end," we backtrack to the deepest vertex that allows us to try a different path.

Which of these traversals would be helpful for finding the **shortest** solution to the word-ladder puzzle?

*BFS*

- There are two general approaches for traversing a graph from some starting vertex  $s$ :
  - Depth First Search (DFS) where you explore as deeply into the graph as possible. If you reach a "dead end," we backtrack to the deepest vertex that allows us to try a different path.
  - Breadth First Search (BFS) where you find all vertices a distance 1 (directly connected) from  $s$ , before finding all vertices a distance 2 from  $s$ , etc.

What data structure would be helpful in each type of search? Why?

a) Breadth First Search (BFS):

*Q*

*myQ = empty queue  
 Mark all vertices as white  
 set start vertex's distance 0 to gray, predictor = None  
 myQueue (start vertex)  
 while myQ is not empty do  
   currentVert = myQ.dequeue()  
   for nextVert its connected to do.  
     if nextVert is white then  
       update distance, color, predictor to currentVert  
       enqueue nextVert  
   currentVert color "black"*

b) Depth First Search (DFS):

*stack instead of queue  
 recursion with run-time stack as the stack*

2. On the next page is the textbook's edge, vertex, and graph implementations.

a) How does this graph implementation maintain its set of vertices?

*VertList dictionary  $O(1)$  by key*

b) How does this graph implementation maintain its set of edges?

*Vertex object has connected to dictionary  
 $O(1)$  to key in edge*

3. Assuming a graph  $g$  containing the word-ladder graph from lecture 25, on the diagram trace the bfs algorithm by showing the value of each vertex's color, predecessor, and distance attributes?

```

""" File: vertex.py """
class Vertex:
    def __init__(self, key, color = 'white',
                  dist = 0, pred = None):
        self.id = key
        self.connectedTo = {}
        self.color = color
        self.predecessor = pred
        self.distance = dist
        self.discovery = 0
        self.finish = 0

    def addNeighbor(self, nbr, weight=0):
        self.connectedTo[nbr] = weight

    def __str__(self):
        return str(self.id) + ' connectedTo: '
        + str([x.id for x in self.connectedTo])

    def getConnections(self):
        return self.connectedTo.keys()

    def getId(self):
        return self.id

    def getWeight(self, nbr):
        return self.connectedTo[nbr]

    def getColor(self):
        return self.color

    def setColor(self, newColor):
        self.color = newColor

    def getPred(self):
        return self.predecessor

    def setPred(self, newPred):
        self.predecessor = newPred

    def getDiscovery(self):
        return self.discovery

    def setDiscovery(self, newDiscovery):
        self.discovery = newDiscovery

    def getFinish(self):
        return self.Finish

    def setFinish(self, newFinish):
        self.finish = newFinish

    def getDistance(self):
        return self.distance

    def setDistance(self, newDistance):
        self.distance = newDistance

```

white = non visited  
gray = partially  
black = completely done

```

""" File: graph.py """
from vertex import Vertex

class Graph:
    def __init__(self):
        self.vertList = {}
        self.numVertices = 0

    def addVertex(self, key):
        self.numVertices = self.numVertices + 1
        newVertex = Vertex(key)
        self.vertList[key] = newVertex
        return newVertex

    def getVertex(self, n):
        if n in self.vertList:
            return self.vertList[n]
        else:
            return None

    def __contains__(self, n):
        return n in self.vertList

    def addEdge(self, f, t, cost=0):
        if f not in self.vertList:
            nv = self.addVertex(f)
        if t not in self.vertList:
            nv = self.addVertex(t)
        self.vertList[f].addNeighbor \
            (self.vertList[t], cost)

    def getVertices(self):
        return self.vertList.keys()

    def __iter__(self):
        return iter(self.vertList.values())

```

```

""" File: graph_algorithms.py """

from graph import Graph
from vertex import Vertex
from linked_queue import LinkedQueue

def bfs(g, start):
    start.setDistance(0)
    start.setPred(None)
    vertQueue = LinkedQueue()
    vertQueue.enqueue(start)
    while (vertQueue.size() > 0):
        currentVert = vertQueue.dequeue()
        for nbr in currentVert.getConnections():
            if (nbr.getColor() == 'white'):
                nbr.setColor('gray')
                nbr.setDistance(currentVert.getDistance()+1)
                nbr.setPred(currentVert)
                vertQueue.enqueue(nbr)
        currentVert.setColor('black')

```