Data Structures (CS 1520) Spring 2019

Time and Place: 8 - 9:15 AM Tuesday and Thursday in ITT 328 and either:
• Section 01: 8 – 9:50 AM Wednesday in Wright 112, or
• Section 02: 10 – 11:50 AM Wednesday in Wright 112.

Web-site: http://www.cs.uni.edu/~fienup/cs1520s19/ (Lecture videos posted after class)

Class Email List: Send messages to Google group for the course at CS-1520-01-spring@uni.edu or CS-1520-02-spring@uni.edu

Instructor: Mark Fienup (fienup@cs.uni.edu)
Office: ITT 313
Phone: 319-273-5918 (Home 319-266-5379)
Office Hours: M: 8-11:45, 1:10-2, T: 9:30-11:45, 1:10-2, W: 1:10-2, Th: 9:30-11:45, 1:10-2, F: 8-10

Prerequisite: Intro. to Computing (CS 1510), and pre- or corequisite Discrete Structures (CS 1800)

Student Learning Outcomes: After this course you should be able to:
1. implement efficient data structures include: stacks, queues, lists, heaps, hash tables, trees, and graphs
2. design and implement "medium" sized programs using functional decomposition and be able to select appropriate data structures.


Assignments: Assignments will consist of weekly laboratory exercises along with concurrent weekly or bi-weekly programming assignments.

Pedagogic Approach: In class, I'll tend to break up the lecture with active (and group) learning exercises to aid learning. While this is not formally graded, part (5%) of your grade will be based on your participation in (and attendance for) these in-class activities. Students benefit by (1) increased depth of understanding, (2) increased comfort and confidence, (3) increased motivation, and (4) being better prepared to work in groups on the job. This might sound great, but it will require you (and me) to work differently to prepare for class. Before the class, you must read the assigned reading, thought about what I asked you to think about, etc.; otherwise you won't be able to effectively participate during class.

Grading policy: There will be three tests (including the final). Tentative test dates and weighting of course components are:

- Attendance and In-class Work: 5%
- Labs: 15%
- Programming Assignments: 20%
- In-class Test 1: 20% (Thursday, February 21)
- In-class Test 2: 20% (Thursday, April 4)
- Final: 20% (Tuesday, May 7 from 8 - 9:50 AM in ITT 328)

Grades will be assigned based on straight percentages off the top student score. If the top student's score is 92%, then the grading scale will be, i.e., 100-82 A, 81.9-72 B, 71.9-62 C, 61.9-52 D, and below 52 F. Plus and minus grades will be assigned for students near cutoff points.

Scholastic Conduct: You are responsible for being familiar with the University's Academic Ethics Policies (http://www.uni.edu/pres/policies/301.shtml). Copying from other students is expressly forbidden. Doing so on exams or assignments will be penalized every time it is discovered. The penalty can vary from zero credit for the copied items (first offense) up to a failing grade for the course. If an assignment makes you realize you don't understand the material, ask questions designed to improve your understanding, not ones designed to discover how another student
solved the assignment. The solutions to assignments should be **individual, original** work unless otherwise specified. Remember: discussing assignments is good. Copying code or test-question answers is cheating.

Any substantive contribution to your assignment solution by another person or taken from a publication (or the web) should be properly acknowledged in writing. Failure to do so is plagiarism and will necessitate disciplinary action. In addition to the activities we can all agree are cheating (plagiarism, bringing notes to a closed book exam, texting during an exam, etc.), assisting or collaborating on cheating is cheating (e.g., supplying code for another student). Cheating can result in failing the course and/or more severe disciplinary actions.

**Student Accessibility Services Statement:**
The University of Northern Iowa (UNI) complies with the Americans with Disabilities Act Amendments Act of 2008 (ADAAA), Section 504 of the Rehabilitation Act of 1973, the Fair Housing Act, and other applicable federal and state laws and regulations that prohibit discrimination on the basis of disability. To request accommodations, it is the policy of the University for students with disabilities to register with Student Accessibility Services (SAS). UNI faculty are not obligated to provide accommodations for students with disabilities without proper notification from SAS and the student. Students may initiate the accommodation process at any time. However, accommodations are not retroactive, and the registration process takes time. Therefore, SAS staff always recommends that students initiate the process as soon as possible rather than wait for academic and social-emotional responsibilities to become overly stressful and/or overwhelming. Please contact SAS, located at ITTC 007, for more information either at (319) 273-2677 or accessibilityservices@uni.edu

### Data Structures Schedule Spring 2019

<table>
<thead>
<tr>
<th>Lect #</th>
<th>Tuesday</th>
<th>Thursday</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1/15 Ch. 1: Python Review; Classes</td>
<td>1/17 Ch. 2: Algorithm Analysis; Big-o; timing</td>
</tr>
<tr>
<td>3</td>
<td>1/22 Ch. 2: Performance of Python Built-in data structures</td>
<td>1/24 Ch. 3: Linear data structures; stack implementations</td>
</tr>
<tr>
<td>5</td>
<td>1/29 Stack applications; Queue implementations</td>
<td>1/31 Queue applications; Deque</td>
</tr>
<tr>
<td>7</td>
<td>2/5 Ch. 6.6: Priority Queue: binary heap implementation</td>
<td>2/7 Unordered Lists implementations</td>
</tr>
<tr>
<td>9</td>
<td>2/12 Ordered List implementation Ch. 4: Recursion; Run-time stack</td>
<td>2/14 Recursion examples: Coin-change problem</td>
</tr>
<tr>
<td>11</td>
<td>2/19 Review for Test 1</td>
<td>2/21 Test 1</td>
</tr>
<tr>
<td>13</td>
<td>2/26 Backtracking v.s. dynamic programming coin-change problem</td>
<td>2/28 Ch. 5: Searching: linear and binary search Hashing: Open-address with rehashing</td>
</tr>
<tr>
<td>15</td>
<td>3/5 Chaining/closed-address hashing Dictionary implementations</td>
<td>3/7 Open-address implementation; Simple sorts: bubble, selecton, and insertion sorts</td>
</tr>
<tr>
<td>17</td>
<td>3/12 Advanced sorts: heap, merge and quick sorts</td>
<td>3/14 Ch. 6: tree terminology, binary tree implementation, parse tree application, and tree traversals</td>
</tr>
<tr>
<td>19</td>
<td>3/26 Binary Search Tree implementation</td>
<td>3/28 Binary Search Tree delete method</td>
</tr>
<tr>
<td>21</td>
<td>4/2 Review for Test 2</td>
<td>4/4 Test 2</td>
</tr>
<tr>
<td>23</td>
<td>4/9 AVL height-balanced trees</td>
<td>4/11 File structures vs. In-memory Data Structures B+ trees</td>
</tr>
<tr>
<td>25</td>
<td>4/16 Ch. 7: Graph terminology, traversals (BFS &amp; DFS) Graph implementations</td>
<td>4/18 Select Graph algorithms: topological sort and Dijkstra’s algorithm</td>
</tr>
<tr>
<td>27</td>
<td>4/23 Select Graph algorithms: Prim’s algorithm; Modifications to BinHeap</td>
<td>4/25 Select Graph algorithms: TSP Backtracking and Best-first search Branch-and-bound</td>
</tr>
<tr>
<td>29</td>
<td>4/30 TSP: Approximation algorithm</td>
<td>5/2 Review for Final Exam</td>
</tr>
</tbody>
</table>

**Final: 8:00 - 9:50 AM Tuesday, May 7 in ITTC 328**
Python Summary

Print Function: the print() function takes a list of values to print and writes them to the output, e.g.,
```
print('cat', 5, 'dog')
print() # blank line
print('pi is about', 3.14)
```
```
cat 5 dog
pi is about 3.14
```

Optional keyword arguments can be used to replace the defaults: space-character (' ') as a separator, the new-line-character (\n') as the ending character, and output file of the console (sys.stdout). The syntax with default parameters explicitly shown is:
```
print(value,...,sep=' ',end='\n', file=sys.stdout)
```

<table>
<thead>
<tr>
<th>Print Function</th>
<th>Expected Output</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>print('cat', 5, end='')</code></td>
<td>cat 5 horse</td>
</tr>
<tr>
<td><code>print('horse')</code></td>
<td>cow</td>
</tr>
<tr>
<td><code>print('cow')</code></td>
<td></td>
</tr>
<tr>
<td><code>print('cat', 5, 'dog', sep='23', end='#')</code></td>
<td>cat23523dog#</td>
</tr>
<tr>
<td><code>print('&gt;', end='')</code></td>
<td>&lt;cat&gt;, &lt;5&gt;, &lt;dog&gt;</td>
</tr>
<tr>
<td><code>print('cat', 5, 'dog', sep='&gt;', end='\n')</code></td>
<td></td>
</tr>
<tr>
<td><code>print('cat', 5, 'dog', sep='23', 'horse')</code></td>
<td>error since keyword arguments must be at the end of the parameter list</td>
</tr>
</tbody>
</table>

String Formatting: inside a string we use formatting placeholders (e.g., %s, %d, %f), follow the string with the format operator (%), and a tuple supplying values for corresponding placeholders. For example:
```
print("Name: %s Age: %7d GPA: %2.2f" % ("Bob", 20, 3.138))
```
```
Name: Bob Age: 20 GPA: 3.14
```

Assignment Statement: the assignment statement creates a variable in memory and sets its value. The syntax is:
```
<variable identifier> = <constant or expression value>, where identifiers must start with a letter or underscore ('_'), and then can be followed by letters, underscores, or digits.
```

<table>
<thead>
<tr>
<th>Assignment Statements and Print Functions</th>
<th>Expected Output</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>a = 123</code></td>
<td>a is 124</td>
</tr>
<tr>
<td><code>b = a</code></td>
<td>b is 123</td>
</tr>
<tr>
<td><code>a = a + 1</code></td>
<td></td>
</tr>
<tr>
<td><code>print ('a is', a)</code></td>
<td></td>
</tr>
<tr>
<td><code>print ('b is', b)</code></td>
<td></td>
</tr>
<tr>
<td><code>c = 'cat'</code></td>
<td>c is catfish</td>
</tr>
<tr>
<td><code>d = c</code></td>
<td>d is cat</td>
</tr>
<tr>
<td><code>c = c + 'fish'</code></td>
<td></td>
</tr>
<tr>
<td><code>print ('c is', c)</code></td>
<td></td>
</tr>
<tr>
<td><code>print ('d is', d)</code></td>
<td></td>
</tr>
<tr>
<td><code>e = ['cat', 'dog']</code></td>
<td>e = ['cat', 'dog', 'cow']</td>
</tr>
<tr>
<td><code>f = e</code></td>
<td>f = ['cat', 'dog', 'cow']</td>
</tr>
<tr>
<td><code>e.append('cow')</code></td>
<td></td>
</tr>
<tr>
<td><code>print('e is', e)</code></td>
<td></td>
</tr>
<tr>
<td><code>print('f is', f)</code></td>
<td></td>
</tr>
</tbody>
</table>

NOTE: This last example deals with assigning lists. In e = ['cat', 'dog'] the variable e is assigned a reference/pointer to the list, so f = e assigns f a reference to the same list. There is only a single list! Thus, when we append 'cow' to the list using e's reference, it also printed in f's list because it's the same list.

The first two examples deal with integers and strings which are immutable (i.e., unchangable). New immutable values are created with new references being assigned. After b = a both variables reference 123. When a = a + 1 executes, a new integer constant of 124 is created and its reference is assign to variable a. Variable b still references 123.
**Input Function:** the `input()` function reads a line from the keyboard (`sys.stdin`) and returns it as a string with the trailing new-line stripped. To input numeric values, the string needs to be explicitly cast (e.g. `eval(input())`). For example, we can input and echo the user's name and age.

```python
name = input("Enter your name: ")
age = eval(input("Enter your age: "))
print('Hi', name, end='!')
print(' Your age is', age)
Enter your name: Bob
Enter your age: 10
Hi Bob! Your age is 10
```

**Control Statements:** the body of control statements are indented and there is NO other "end" (\"\") marker

**if statements:** An if statement allows code to be executed or not based on the result of a comparison. If the condition evaluates to True, then the statements of the indented body is executed. If the condition is False, then the body is skipped. The syntax of if statements is:

```python
if <condition>:
    statement_1
    statement_2
    statement_3
else:
    statement_1
    statement_2
```

Typically, the condition involves comparing “stuff” using relational operators (`<`, `>`, `==`, `<=`, `>=`, `!=`).

Complex conditions might involve several comparisons combined using Boolean operators: `not`, `or`, and. For example, we might want to print “Your grade is B.” if the variable score is less than 90, but greater than or equal to 80.

```python
if score < 90 and score >= 80:
    print("Your grade is B.")
```

The precedence for mathematical operators, Boolean operators, and comparisons are given in the table.

<table>
<thead>
<tr>
<th>Operator(s)</th>
<th>** (exponential)</th>
<th>+, - (add, sub)</th>
<th>*, / % (rem), // (integer div)</th>
</tr>
</thead>
<tbody>
<tr>
<td>highest</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>lowest</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

**for loop:** the for loop iterates once for each item in some sequence type (i.e, list, tuple, string).

```python
for value in [1, 3, 9, 7]:
    print(value)
for character in 'house':
    print(character)
```

The for loop iterates over an iterable data-structure object (list, string, dictionary) or a range object created by the built-in `range` function which generate each value one at a time for each iteration of the loop. The syntax of: `range([start,] end, [, step])`, where [] are used to denote optional parameters. Examples:
- `range(5)` generates the sequence of values: 0, 1, 2, 3, 4
- `range(2, 7)` generates the sequence of values: 2, 3, 4, 5, 6
- `range(10, -1)` generates the sequence of values: 10, 9, 8, 7, 6, 5, 4, 3

For example:

```python
for count in range(1,6):
    print(count, end=" ")
print("\nDone")
```
**while loop:** A while statement allows code to be executed repeated (zero or more times) as long as the condition evaluates to True. The syntax of a while statement is:

```
while <condition>:
    statement_1
    statement_2
    statement_3
```

An infinite loop is one that would loop forever. (FYI, in a Python shell ctrl-c (^c) can be used to kill the running program.) Most infinite loops are caused by programmer error, but sometimes they are intentional. The following “sentinel-controlled” code uses an infinite loop and a `break` statement that immediately causes control to exit the loop.

```python
total = 0
counter = 0
while True:  # an infinite loop
    score = eval(input("Enter a score (or negative value to exit): "))
    if score < 0:
        break
    total += score
    counter += 1
print("Average is", total/counter)
```

**Strings:** Strings in Python are sequential collections of only characters. Strings are immutable (i.e., cannot be changed), so new strings are generated by string operations. Operations on strings (or any sequence collection) include:

<table>
<thead>
<tr>
<th>Operation</th>
<th>Operator</th>
<th>Explanation</th>
<th>Example</th>
<th>Result of Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>Indexing</td>
<td>[ &lt;index&gt; ]</td>
<td>Access the element specified by the index</td>
<td><code>myString[1]</code></td>
<td>‘e’</td>
</tr>
<tr>
<td>Slicing</td>
<td>[ : ]</td>
<td>Extract a part of the string</td>
<td><code>myString[1:5]</code></td>
<td>‘ello’</td>
</tr>
<tr>
<td>Concatenation</td>
<td>+</td>
<td>Combine strings together</td>
<td><code>myString + aString</code></td>
<td>‘Hello!!!cat’</td>
</tr>
<tr>
<td>Repetition</td>
<td>*</td>
<td>Concatenate a repeated number of times</td>
<td><code>aString * 3</code></td>
<td>‘catcatcat’</td>
</tr>
<tr>
<td>Membership</td>
<td>in</td>
<td>Ask whether a substring is in a string</td>
<td>‘ell‘ in <code>myString</code></td>
<td>True</td>
</tr>
<tr>
<td>Length</td>
<td><code>len(string)</code></td>
<td>How many items are in the string?</td>
<td><code>len( </code>myString<code> )</code></td>
<td>8</td>
</tr>
</tbody>
</table>

Indexing of strings starts with 0 on the left end, and -1 on the right end:

```
1111
01234567890123
cheer = ‘GO Panthers!!!’
```

```
-4-3-2-1
```

Omitted indexes in a slice means “from the end.” For example, `cheer[4]` generates ‘GO P’.

Omitted indexes in a slice means “from the end.” For example, `cheer[-4:]` generates ‘s!!!’. 
String objects also have the following methods: (the `string` module can be imported to provide more operations.)

<table>
<thead>
<tr>
<th>Method</th>
<th>Usage</th>
<th>Explanation</th>
</tr>
</thead>
<tbody>
<tr>
<td>center</td>
<td><code>myString.center(w)</code></td>
<td>Returns a string with <code>myString</code> centered in a field of size <code>w</code></td>
</tr>
<tr>
<td>ljust</td>
<td><code>myString.ljust(w)</code></td>
<td>Returns a string with <code>myString</code> left-justified in a field of size <code>w</code></td>
</tr>
<tr>
<td>rjust</td>
<td><code>myString.rjust(w)</code></td>
<td>Returns a string with <code>myString</code> right-justified in a field of size <code>w</code></td>
</tr>
<tr>
<td>upper</td>
<td><code>myString.upper()</code></td>
<td>Returns a string with <code>myString</code> in all upper-case characters</td>
</tr>
<tr>
<td>lower</td>
<td><code>myString.lower()</code></td>
<td>Returns a string with <code>myString</code> in all lower-case characters</td>
</tr>
<tr>
<td>strip</td>
<td><code>myString.strip()</code></td>
<td>Returns a string with leading and trailing whitespace (space, tab, new-line) chars. removed. An optional string parameter can be used to supply characters to strip instead of whitespace.</td>
</tr>
<tr>
<td>count</td>
<td><code>myString.count(sub)</code></td>
<td>Returns number of occurrences of <code>sub</code> in <code>myString</code> (Optional parameters: <code>myString.count(sub [, start [, end] ]</code>)</td>
</tr>
<tr>
<td>endswith</td>
<td><code>myString.endswith(sub)</code></td>
<td>Returns True if <code>myString</code> ends with the substring <code>sub</code>; otherwise it returns False</td>
</tr>
<tr>
<td>startswith</td>
<td><code>myString.startswith(sub)</code></td>
<td>Returns True if <code>myString</code> starts with the substring <code>sub</code>; otherwise it returns False</td>
</tr>
<tr>
<td>isdigit</td>
<td><code>myString.isdigit()</code></td>
<td>Returns True if <code>myString</code> contains only digits; otherwise it returns False</td>
</tr>
<tr>
<td>isalpha</td>
<td><code>myString.isalpha()</code></td>
<td>Returns True if <code>myString</code> contains only letters; otherwise it returns False</td>
</tr>
<tr>
<td>split</td>
<td><code>myString.split()</code></td>
<td>Returns a list of substrings of <code>myString</code> splits at whitespace characters. An optional string parameter can supply characters to split on.</td>
</tr>
<tr>
<td>find</td>
<td><code>myString.find(sub)</code></td>
<td>Returns the starting index of the first occurrence of <code>sub</code>. (Optional parameters: <code>myString.find(sub [, start [, end] ]</code>)</td>
</tr>
<tr>
<td>replace</td>
<td><code>myString.replace(old,new)</code></td>
<td>Returns a string with all occurrences of substring <code>old</code> replaced by substring <code>new</code>. An additional integer parameter can specify the number of replacements to perform, e.g., <code>myString.replace(old,new,3)</code></td>
</tr>
</tbody>
</table>

**Lists:** A Python list is also a sequence collection, but a list can contain items of any type (e.g., character, strings, integers, floats, other lists, etc.), and lists are mutable. Lists are represented by comma-separated values enclosed in square brackets ("[", "]"). Operations on lists (or any sequence collection, e.g., strings) include:

<table>
<thead>
<tr>
<th>Operation</th>
<th>Operator</th>
<th>Explanation</th>
<th>Example</th>
<th>Result of Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>Slicing</td>
<td><code>[ : ]</code></td>
<td>Extract a part of the list</td>
<td><code>myList[ 1:3 ]</code></td>
<td><code>[6, 7]</code></td>
</tr>
<tr>
<td>Concatenation</td>
<td><code>+</code></td>
<td>Combine lists together</td>
<td><code>myList + ListB</code></td>
<td><code>[5, 6, 7, 8, 8, 9]</code></td>
</tr>
<tr>
<td>Repetition</td>
<td><code>*</code></td>
<td>Concatenate a repeated number of times</td>
<td><code>ListB * 3</code></td>
<td><code>[8, 9, 8, 9, 8, 9]</code></td>
</tr>
<tr>
<td>Membership</td>
<td><code>in</code></td>
<td>Ask whether an item is in a list</td>
<td><code>3 in myList</code></td>
<td><code>False</code></td>
</tr>
<tr>
<td>Length</td>
<td><code>len(list)</code></td>
<td>How many items are in the list?</td>
<td><code>len( myList )</code></td>
<td><code>4</code></td>
</tr>
</tbody>
</table>
The following list methods are provided by Python:

<table>
<thead>
<tr>
<th>Method</th>
<th>Usage</th>
<th>Explanation</th>
</tr>
</thead>
<tbody>
<tr>
<td>append</td>
<td>myList.append(item)</td>
<td>Adds item to the end of myList</td>
</tr>
<tr>
<td>extend</td>
<td>myList.extend(otherList)</td>
<td>Extends myList by adding all items in otherList to myList's end</td>
</tr>
<tr>
<td>insert</td>
<td>myList.insert(i, item)</td>
<td>Insert item in myList at index i</td>
</tr>
<tr>
<td>pop</td>
<td>myList.pop( )</td>
<td>Remove and return the last item in myList</td>
</tr>
<tr>
<td>pop(i)</td>
<td>myList.pop(i)</td>
<td>Remove and return the ith item in myList *</td>
</tr>
<tr>
<td>del</td>
<td>del myList[i]</td>
<td>Deletes the item in the ith position of myList *</td>
</tr>
<tr>
<td>remove</td>
<td>myList.remove(item)</td>
<td>Removes the first occurrence of item in myList **</td>
</tr>
<tr>
<td>index</td>
<td>myList.index(item)</td>
<td>Returns the index of the first occurrence of item in myList **</td>
</tr>
<tr>
<td>count</td>
<td>myList.count(item)</td>
<td>Returns the number of occurrences of item in myList</td>
</tr>
<tr>
<td>sort</td>
<td>myList.sort( )</td>
<td>Modifies myList to be sorted</td>
</tr>
<tr>
<td>reverse</td>
<td>myList.reverse( )</td>
<td>Modifies myList to be in reverse order</td>
</tr>
</tbody>
</table>

* Note: raises an IndexError if the index i is not in the list
** Note: raises a ValueError if the item is not in the list

**Tuples:** A tuple is another sequence data type, so the sequence operations of indexing, slicing, concatenation, repetition, membership (in), and len() work on tuples too. Tuples are very similar to lists, i.e., comma-separated items enclosed in parentheses. The main difference is that **tuples are immutable** (cannot be modified).

Create two tuples as:

```python
student1 = ('Bob', 123456, 'Jr', 3.12)
student2 = 'Sally', 654321, 'Fr', 0.0
```

In addition to indexing, "fields" of a tuple can be **unpacked** using a single assignment statement as:

```python
name, idnum, rank, gpa = student1
```

(NOTE: This allows multiple values to be returned from a function)

**Dictionaries:** A dictionary is an unordered set of key-value pairs (written as key:value). Keys must be unique and immutable (e.g., numerics, strings, tuples of immutable objects). Dictionaries are typically used to lookup the value corresponding to a specified key. Dictionaries can be written as comma-separated key:value pairs enclosed in curly braces. For example,

```python
phoneNumbers = {'fienup':35918, 'gray':35917, 'east':32939, 'drake':35811, 'schafer':32187}
```

Access to individual key:value pairs looks syntactically like a sequence lookup using a key instead of an index. For example, `phoneNumbers['east']` returns 32939, and a new key:value pair can be added by `phoneNumbers['wallingford'] = 35919`. Additional, methods on dictionaries are:

<table>
<thead>
<tr>
<th>Method</th>
<th>Usage</th>
<th>Explanation</th>
</tr>
</thead>
<tbody>
<tr>
<td>keys</td>
<td>myDictionary.keys()</td>
<td>Returns keys in an iterable dict_keys object</td>
</tr>
<tr>
<td>values</td>
<td>myDictionary.values()</td>
<td>Returns values in an iterable dict_values object</td>
</tr>
<tr>
<td>items</td>
<td>myDictionary.items()</td>
<td>Returns key:value tuples in an iterable dict_items object</td>
</tr>
<tr>
<td>get</td>
<td>value = myDictionary[myKey]</td>
<td>Returns the value associated with myKey; otherwise raises a KeyError if myKey is not in the dictionary</td>
</tr>
<tr>
<td>get</td>
<td>myDictionary.get(myKey)</td>
<td>Returns the value associated with myKey; otherwise None</td>
</tr>
<tr>
<td>get</td>
<td>myDictionary.get(myKey, alt)</td>
<td>Returns the value associated with myKey; otherwise alt</td>
</tr>
<tr>
<td>in</td>
<td>myKey in myDictionary</td>
<td>Returns True if myKey is in myDictionary; otherwise False</td>
</tr>
<tr>
<td>del</td>
<td>del myDictionary[myKey]</td>
<td>Deletes the key:value pair whose key is myKey</td>
</tr>
</tbody>
</table>
## Python Summary

### Text Files:
Below is a summary of the important text-file operations in Python.

<table>
<thead>
<tr>
<th>General Syntax</th>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>open(filename)</td>
<td>f = open('data.txt', 'w')</td>
<td>Modes: 'r' read only; 'w' write only; 'a' append; 'r+' both reading and writing. Default mode is 'r'.</td>
</tr>
<tr>
<td>open(filename, mode)</td>
<td>f.close()</td>
<td>Close the file to free up system resources.</td>
</tr>
<tr>
<td>f.close()</td>
<td>for line in f:</td>
<td>Memory efficient, fast and simple code to loop over each line in the file.</td>
</tr>
<tr>
<td>loop over the file</td>
<td>print (line)</td>
<td></td>
</tr>
<tr>
<td>object</td>
<td>f.readline()</td>
<td>Returns the next line from the file. The newline ('\n') character is left at the end of the string.</td>
</tr>
<tr>
<td>f.write(string)</td>
<td>f.write('cats and dogs\n')</td>
<td>Writes the string to the file.</td>
</tr>
<tr>
<td>f.read()</td>
<td>all = f.read()</td>
<td>Returns the whole file as a single string.</td>
</tr>
<tr>
<td>f.read(size)</td>
<td>chunk = f.read(100)</td>
<td>Returns a string of at most 100 (size) bytes. If the file has been completely read, an empty string is returned.</td>
</tr>
<tr>
<td>f.readlines()</td>
<td>allLines = f.readlines()</td>
<td>Returns a list containing all the lines of the file.</td>
</tr>
<tr>
<td>f.readlines(size)</td>
<td>somelines = f.readlines(5000)</td>
<td>Returns the next 5000 bytes of line. Only complete lines will be returned.</td>
</tr>
</tbody>
</table>

Below is a summary of the important file-system functions from the `os` and `os.path` modules in Python.

### os Module File-system Functions

<table>
<thead>
<tr>
<th>General Syntax</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>getcwd()</td>
<td>Returns the complete path of the current working directory</td>
</tr>
<tr>
<td>chdir(path)</td>
<td>Changes the current working directory to path</td>
</tr>
<tr>
<td>listdir(path)</td>
<td>Returns a list of the names in directory named path</td>
</tr>
<tr>
<td>mkdir(path)</td>
<td>Creates a new directory named path and places it in the current working directory</td>
</tr>
<tr>
<td>rmdir(path)</td>
<td>Removes the directory named path from the current working directory</td>
</tr>
<tr>
<td>remove(path)</td>
<td>Removes the file named path from the current working directory</td>
</tr>
<tr>
<td>rename(old, new)</td>
<td>Renames the file or directory named old to new</td>
</tr>
</tbody>
</table>

### os.path Module File-system Functions

<table>
<thead>
<tr>
<th>General Syntax</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>exists(path)</td>
<td>Returns True if path exists and False otherwise</td>
</tr>
<tr>
<td>isdir(path)</td>
<td>Returns True if path is a directory and False otherwise</td>
</tr>
<tr>
<td>isfile(path)</td>
<td>Returns True if path is a file and False otherwise</td>
</tr>
<tr>
<td>getsize(path)</td>
<td>Returns the size in bytes of the object named path</td>
</tr>
</tbody>
</table>

**NOTE:** The initial "current working directory" is the directory where the program is located. Typically, it is useful to access files relative to the "current working directory" instead of specifying an absolute (complete) path. You can use the strings:

- `'.' to specify the current working directory, e.g., `currentDirectoryList = os.listdir('.'))`
- `'' ..' to specify the parent of current working directory, e.g., `os.chdir('..')` which changes the current working directory to the parent directory
Functions:
A function is a procedural abstract, i.e., a named body of code that performs some task when it is called/invoked. Often a function will have one or more parameter that allows it to perform a more general (variable) task. For example, the cube function below can be called with any numeric value with the corresponding cube of that number being returned.

```python
# Function to calculate the cube of a number
def cube(num):
    num_squared = num * num
    return num_squared * num

# call the function
value = 2
print('The value', value, 'raised to the power 3 is', cube(value))
print('The value 3 raised to the power 3 is', cube(3))
```

Terminology:
- a formal parameter is the name of the variable used in the function definition. It receives a value when the function is called. In the function cube, num is the formal parameter. Formal parameters are only known inside of the function definition. The section of a program where a variable is known is called its scope, so the scope of a formal parameter (and other local variable defined in the function such as num_squared) is limited to the function in which it is defined.
- an actual parameter/argument is the value used in the function call that is sent to the function. In the call to function cube, the variable value supplies the actual parameter value of 2.
- a global variable is created outside all functions and is known throughout the whole program file, e.g. value.

It is helpful to understand the “rules of the game” when a function is called. Memory is used to store the current program and the data associated with it. The memory used to store the data is divided as shown below.
- Global memory is used to store the global variables (and constants).
- The heap is used to store dynamically allocated objects as the program runs, e.g. lists, strings, ints, objects.
- The run-time stack is used to store call-frames (or activation records) that get pushed on the stack when a function is called, and popped off the stack when a function returns.

When a function is called the section of code doing the calling is temporarily suspended, and a new call-frame gets pushed on top of the stack before execution of the function body. The call-frame contains the following information about the function being called:
- the return address -- the spot in code where the call to the function occurred. This is needed so execution (control) can return there when the end of the function is reached or a return statement executes.
- room to store the formal parameters used by the function. In Python, parameters are passed-by-value which means that the value of each actual parameter in the function call is assigned to the corresponding formal parameter in the function definition before the function starts executing. However, the memory location for actual parameters for strings, lists, dictionaries, tuples, atomic objects contain only references to the heap.
- room to store the local variables defined in the function (these are probably references to objects in the heap).

When a function returns, execution resumes at the function call (which is specified by the return address). A function typically sends back a value to the call by specifying an expression after return in the return statement. In Python if no expression is specified returned, then the special object None is returned.
def play(myInt, myLongInt, myList, myString):
    print('START OF play Function')
    print('myInt=', myInt, 'myLongInt=', myLongInt)
    print('myList=', myList, 'myString=', myString)
    myInt += 1
    myLongInt += 1
    myList.append(1)
    myString += 'a'
    print('END OF play Function')
    print('myInt=', myInt, 'myLongInt=', myLongInt)
    print('myList=', myList, 'myString=', myString)
    return

anInt = 10  # 1st STATEMENT EXECUTED, ABOVE IS SKIPPED
aLongInt = 123456789012345678901234567890L
aList = range(5)
aString = 'hello'
print('BEFORE CALL')
print('anInt=', anInt, 'aLongInt=', aLongInt)
print('aList=', aList, 'aString=', aString)
play(anInt, aLongInt, aList, aString)
print('AFTER CALL')
print('anInt=', anInt, 'aLongInt=', aLongInt)
print('aList=', aList, 'aString=', aString)

Output of complete program:

```python
>>> BEFORE CALL
anInt= 10 aLongInt= 123456789012345678901234567890
aList= [0, 1, 2, 3, 4] aString= hello
START OF play Function
myInt= 10 myLongInt= 123456789012345678901234567890
myList= [0, 1, 2, 3, 4] myString= hello
END OF play Function
myInt= 11 myLongInt= 123456789012345678901234567891
myList= [0, 1, 2, 3, 4, 1] myString= helloa
AFTER CALL
anInt= 10 aLongInt= 123456789012345678901234567890
aList= [0, 1, 2, 3, 4, 1] aString= hello
``` 

NOTES:
- `aList` passes a copy of its reference to the list as the initial value of the formal parameter `myList`, so both refer to the single list. In `play` when `myList.append(1)` executes, the single list is changed. Thus, when play terminates, `aList` still reflect this change.
- `aString` passes a copy of its reference to the string ('hello') as the initial value of the formal parameter `myString`. Since strings are immutable (cannot be changed) in Python, executing `myString += 'a'` cause a whole new string 'helloa' to be recreated in memory with `myString` referencing it. `aString` still refers to the string 'hello'.
- the `int` objects should actually be stored in the heap like the string objects since `int`'s are immutable.
Classes: A class definition is like a blueprint (recipe) for each of the objects of that class. A class specifies a set of data attributes and methods for the objects of that class:
- The values of the data attributes of a given object make up its state
- The behavior of an object depends on its current state and on the methods that manipulate this state
- The set of a class's methods is called its interface

The general syntax of class definition is:

```python
class MyClass ( ( superClass1 [, superClass2 ]* ) ) :
    '''Document comment which becomes the __doc__ attribute for the class'''
    def __init__(self, [param1, param2]*):
        '''Document comment for constructor method with self be referencing to the object itself'''
        #_init_body

    # def's of other class methods and assignments to class attributes
    # end class MyClass
```

File: simple_die.py
Description: This module defines a six-sided Die class.

```python
from random import randint

class Die(object):
    '''This class represents a six-sided die.'''
    def __init__(self):
        '''The initial face of the die.'''
        self.currentRoll = randint(1, 6)
    def roll(self):
        '''Resets the die's value to a random number between 1 and 6.'''
        self.currentRoll = randint(1, 6)
    def getRoll(self):
        '''Returns the face value of the die.'''
        return self.currentRoll
    def __str__(self):
        '''Returns the string representation of the die.'''
        return str(self.currentRoll)
```

Consider the following script to test the Die class and its associated output:

```python
# testDie.py - script to test Die class
from simple_die import Die
die1 = Die()
die2 = Die()
print('die1 = ', die1)  #calls __str__
print('die2 = ', die2)
print()
print('die1.getRoll() = ', die1.getRoll())
print('die2.getRoll() = ', die2.getRoll())
die1.roll()
print('die1.getRoll() = ', die1.getRoll())
print('str(die1): ' + str(die1))
print('die1 + die2: ', die1.getRoll() + die2.getRoll())

>>>
die1 = 2
die2 = 5
die1.getRoll() = 2
die2.getRoll() = 5
die1.getRoll() = 3
str(die1): 3
die1 + die2: 8
```
Classes in Python have the following characteristics:
- all class attributes (data attributes and methods) are public by default, unless your identifier starts with a single underscores, e.g. self._currentRoll
- all data types are objects, so they can be used as inherited base classes
- objects are passed by reference when used as parameters to functions
- all classes have a set of standard methods provided, but may not work properly (_str_, _doc_, etc.)
- most built-in operators (+, -, *, <, >, ==, etc.) can be redefined for a class. This makes programming with objects a lot more intuitive. For example suppose we have two Die objects: die1 & die2, and we want to add up their combined rolls. We could use accessor methods to do this:
  
  diceTotal = die1.getRoll() + die2.getRoll()

Here, the getRoll method returns an integer (type int), so the ‘+’ operator being used above is the one for int. But, it might be nice to ‘overload’ the + operator by defining an _add_ method as part of the Die class, so the programmer could add dice directly as in:

  diceTotal = die1 + die2

The three most important features of Object-Oriented Programming (OOP) to simplify programs and make them maintainable are:

1. **encapsulation** - restricts access to an object's data to access only by its methods
   - helps to prevent indiscriminative changes that might cause an invalid object state (e.g., 6-side die with a of roll 8)

2. **inheritance** - allows one class (the subclass) to pickup data attributes and methods of other class(es) (the parents)
   - helps code reuse since the subclass can extend its parent class(es) by adding addition data attributes and/or methods, or overriding (through polymorphism) a parent's methods

3. **polymorphism** - allows methods in several different classes to have the same names, but be tailored for each class
   - helps reduce the need to learn new names for standard operations (or invent strange names to make them unique)

Consider using inheritance to extend the Die class to a generalized AdvancedDie class that can have any number of sides. The interface for the AdvancedDie class are:

<table>
<thead>
<tr>
<th>Method</th>
<th>Example Usage</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>init</strong></td>
<td>myDie = AdvancedDie(8)</td>
<td>Constructs a die with a specified number of sides and randomly rolls it (Default of 6 sides if no argument supplied)</td>
</tr>
<tr>
<td>getRoll</td>
<td>myDie.getRoll()</td>
<td>Returns the current roll of the die (inherited from Die class)</td>
</tr>
<tr>
<td>getSides</td>
<td>myDie.getSides()</td>
<td>Returns the number of sides on the die (did not exist in Die class)</td>
</tr>
<tr>
<td>roll</td>
<td>myDie.roll()</td>
<td>Rerolls the die randomly (By overriding the roll method of Die, an AdvancedDie can generate a value based on its # of sides)</td>
</tr>
<tr>
<td><strong>eq</strong></td>
<td>if myDie == otherDie:</td>
<td>Allows == operations to work correctly for AdvancedDie objects.</td>
</tr>
<tr>
<td><strong>lt</strong></td>
<td>if myDie &lt; otherDie:</td>
<td>Allows &lt; operations to work correctly for AdvancedDie objects.</td>
</tr>
<tr>
<td><strong>gt</strong></td>
<td>if myDie &gt; otherDie:</td>
<td>Allows &gt; operations to work correctly for AdvancedDie objects.</td>
</tr>
<tr>
<td><strong>add</strong></td>
<td>sum = myDie + otherDie</td>
<td>Allows the direct addition of AdvancedDie objects, and returns the integer sum of their current roll values.</td>
</tr>
<tr>
<td><strong>str</strong></td>
<td>myDie.<em>str</em>() or str(myDie)</td>
<td>Returns a string representation for the AdvancedDie. By overriding the <em>str</em> method of the Die class, so the “print” statement will work correctly with an AdvancedDie.</td>
</tr>
</tbody>
</table>
Consider the following script and associated output:

```python
# testAdvancedDie.py - script to test AdvancedDie class
from advanced_die import AdvancedDie

die1 = AdvancedDie(100)
die2 = AdvancedDie(100)
die3 = AdvancedDie()

print('die1 =', die1)     # calls __str__
print('die2 =', die2)
print('die3 =', die3)

print('die1.getRoll() = ', die1.getRoll())
print('die1.getSides() = ', die1.getSides())
die1.roll()
print('die1.getRoll() = ', die1.getRoll())
print('die2.getRoll() = ', die2.getRoll())
print('die1 == die2:', die1==die2)
print('die1 < die2:', die1<die2)
print('die1 > die2:', die1>die2)
print('die1 != die2:', die1!=die2)
print('str(die1): ' + str(die1))
print('die1 + die2:', die1 + die2)

help(AdvancedDie)
```

```
die1 = Number of Sides=100 Roll=32
die2 = Number of Sides=100 Roll=76
die3 = Number of Sides=6 Roll=5
die1.getRoll() = 32
die1.getSides() = 100

die1.getRoll() = 70
die2.getRoll() = 76
die1 == die2: False
die1 < die2: True
die1 > die2: False
die1 != die2: True
str(die1): Number of Sides=100 Roll=70
die1 + die2: 146
```

Help on class AdvancedDie in module advanced_die:

```
class AdvancedDie(simple_die.Die):
    Advanced die class that allows for any number of sides
    | Method resolution order:
    | AdvancedDie
    | simple_die.Die
    | __builtin__.object
    | Methods defined here:
```
The AdvancedDie class that inherits from the Die superclass.

```python
""
File: advanced_die.py
Description: Provides an AdvancedDie class that allows for any number of sides
Inherits from the parent class Die in module die_simple
""
from simple_die import Die
from random import randint

class AdvancedDie(Die):
    """Advanced die class that allows for any number of sides""
    def __init__(self, sides = 6):
        """Constructor for any sided Die that takes an the number of sides
        as a parameter; if no parameter given then default is 6-sided.""
        Die.__init__(self)    # call Die parent class constructor
        self._numSides = sides
        self._currentRoll = randint(1, self._numSides)

    def roll(self):
        """Causes a die to roll itself -- overrides Die class roll""
        self._currentRoll = randint(1, self._numSides)

    def __eq__(self, rhs_Die):
        """Overrides default '__eq__' operator to allow for deep comparison of Dice""
        return self._currentRoll == rhs_Die._currentRoll

    def __lt__(self, rhs_Die):
        """Overrides default '__lt__' operator to allow for deep comparison of Dice""
        return self._currentRoll < rhs_Die._currentRoll

    def __gt__(self, rhs_Die):
        """Overrides default '__gt__' operator to allow for deep comparison of Dice""
        return self._currentRoll > rhs_Die._currentRoll

    def __str__(self):
        """Returns the string representation of the AdvancedDie.""
        return 'Number of Sides='+str(self._numSides)+' Roll='+str(self._currentRoll)

    def __add__(self, rhs_Die):
        """Returns the sum of two dice rolls""
        return self._currentRoll + rhs_Die._currentRoll

    def getSides(self):
        """Returns the number of sides on the die.""
        return self._numSides
```
1. The print function has optional keyword arguments which can be listed last that modify it behavior. The print function syntax: `print(value,...,sep=' ',end='
', file=sys.stdout)

a) Predict the expected output of each of the following.

<table>
<thead>
<tr>
<th>Program</th>
<th>Expected Output</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>print('cat',5,'dog')</code></td>
<td></td>
</tr>
<tr>
<td><code>print()</code></td>
<td></td>
</tr>
<tr>
<td><code>print('cat',5,end='')</code></td>
<td></td>
</tr>
<tr>
<td><code>print('horse')</code></td>
<td></td>
</tr>
<tr>
<td><code>print('cow')</code></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Program</th>
<th>Expected Output</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>print ('cat',5,'dog',end='#',sep='23')</code></td>
<td></td>
</tr>
<tr>
<td><code>print ('cat',5,'dog',sep='23','horse')</code></td>
<td></td>
</tr>
<tr>
<td><code>print ('cat',5,'dog',sep='&gt;!*3')</code></td>
<td></td>
</tr>
</tbody>
</table>

2. Review of assignment statements. Predict the output of the following programs

```
# Assignment statements
a = 123
b = a
a += 1
print ('a is', a)
print ('b is', b)
```

```
c = ['cat', 'dog']
d = c
c.append('cow')
print('c is', c)
print('d is', d)
```

```
c = 'cat'
d = c
c += 'fish'
print('c is', c)
print('d is', d)
```

---

Most simple programs have a similar functional-decomposition design pattern (IPO - Input, Process, Output):

```
input, initialize variables  main  results params.

Input - welcome user
- initialize variables
- prompt for input

Process - compute answer

subtask 1

Output - display results

subtask 2
```

""" Simple IPO program to sum a list of numbers. """

```python
def main():
    label, values = getInput()
    total = sum(values)
    displayResults(label, total)

def getInput():
    """ Get label and list of values to sum. """
    label = input("What are we summing? ")
    numberOfValues = int(input("How many values are there? "))
    values = []
    for i in range(numberOfValues):
        values.append(eval(input("Enter the next number: ")))
    return label, values
def displayResults(label, total):
    """ Display sum of values. """
    print("The sum of", label, "values is", total)

main()  # starts the main function running
```

What are we summing? money
How many values are there? 4
Enter the next number: 10
Enter the next number: 20
Enter the next number: 30
Enter the next number: 50
The sum of money values is 110
3. Design a program to roll two 6-sided dice 1,000 times to determine the percentage of each outcome (i.e., sum of both dice). Report the outcome(s) with the highest percentage.

Customize the diagram for the dice problem by briefly describing what each function does and what parameters are passed.
Data Structures (CS 1520)  Lab 1  Name:________________________

Objective: To practice writing Python code.

To start the lab: Download and unzip the file lab1.zip from
http://www.cs.uni.edu/~fienup/cs1520s19/labs/lab1.zip

Part A: In the folder lab1, open the diceOutcomes.py program in IDLE. (Right-click on diceOutcomes.py | Edit with IDLE) It contains a partial program we started to discuss in class to solve the problem:

"Write a program to roll two 6-sided dice 1,000 times to determine the percentage of each outcome (i.e., sum of both dice). Report the outcome(s) with the highest percentage."

I decided to functional-decomposition this problem as:

```
main - provides an outline of program by calling top-level functions
welcomeAndInpuntRolls - Displays welcome message for the user. Gets and returns the number of dice rolls from the user.
calculateFrequentRolls - Rolls the dice the correct number of times, tallies the outcomes, and returns a list of outcomes with the highest count and highest count.
rollAndTallyOutcomes - Rolls the dice the correct number of times and tallies the outcomes. Returns a list of tallies with the index being the outcome.
max - built-in function to return the largest item in an iterable data structure like a list.
findOutcomes - Returns a list of outcomes with the highest count.
displayResults - Displays the outcome(s) with the highest percentage.
```

Consider running the program with only 10 dice rolls instead of 1,000. The program output with some extra debugging prints showing the two Python lists used: outcomeCounts and mostFrequentRolls.

```
This program rolls two 6-sided dice many times to determine the outcome(s) with the highest percentage.
How many times would you like to roll the pair of dice? 10
outcomeCounts: [0, 0, 1, 0, 2, 1, 0, 3, 0, 3, 0, 0]
mostFrequentRolls: [7, 10] and highestCount: 3
The highest percentage is 30.0 for outcome(s): 7 10
```

Your task for lab 1 is to complete the code for the rollAndTallyOutcomes and findOutcomes functions.

After you have working code, raise your hand and demonstrate your code.

If you complete all parts of the lab, nothing needs to be turned in for this lab. If you do not get done today, then show me the completed lab in next week’s lab period. When done, remember save your program to a USB drive, email to yourself, Google drive, etc.
EXTRA CREDIT -- Part B: Rewrite the program using a dictionary instead of a list for the outcomeCounts. Your dictionary will have the outcome for the key and its corresponding tally as its value.

After you have working code, raise your hand and demonstrate your code.
Data Structures (CS 1520)  
Lecture 2  
Name: ____________

1. An alternative to functional-decomposition design is to use object-oriented design (OOD). For the following program, what objects would be useful and what methods (operations on the objects) should each support?

   “Write a program to roll two 6-sided dice 1,000 times to determine the percentage of each outcome (i.e., sum boths dice). Report the outcome(s) with the highest percentage.” (You only need consider the program’s OOD)

2. Consider the Die and AdvancedDie classes from the Python Summary handout.
   a) What data attributes of AdvancedDie are inherited from the parent Die class?
   b) What new data attributes are added as part of the subclass AdvancedDie?
   c) Which Die class methods are used directly for an AdvancedDie object?
   d) Which Die class methods are redefined/overridden by the AdvancedDie object?
   e) Which methods are new to the AdvancedDie class and not in the Die class?

   f) If die1 and die2 are AdvancedDie objects, then the statement “if die1 == die2;” invokes the __eq__ method of AdvancedDie with die1 “passed” as self and die2 passed as rhs_Die.

   ```python
   def __eq__(self, rhs_Die):
       """Overides default '__eq__' operator to allow for deep comparison of dice""
       return self._currentRoll == rhs_Die._currentRoll
   ```

   What would the code be for AdvancedDie __le__ method to allow for the “if die1 <= die2;” statement?

   g) Good software engineering practice is to include precondition and postcondition comments on each method/function where the:
   - precondition - indicates what must be true for the method to work correctly. Typically, the precondition describes the valid values of the parameters. If the precondition is not satisfied, the method does not need to work correctly!
   - postcondition - describes the expected state after the method has executed

   Consider the AdvancedDie constructor:

   ```python
   class AdvancedDie(Die):
       """Advanced die class that allows for any number of sides""
       def __init__(self, sides = 6):
           """Constructor for any sided Die that takes an number of sides as a parameter; if no parameter given then default is 6-sided.""
           Die.__init__(self) # call Die parent class constructor
           self.numSides = sides
           self._currentRoll = randint(1, self.numSides)
   ```

   What precondition and postcondition comments should we add?

   h) If a method/function has a precondition that is not met when invoked (e.g., die1 = AdvancedDie("six")), why should the method raise an error?
3. General "Algorithmic-Complexity Analysis" terminology:

problem - question we seek an answer for, e.g., "What is the sum of all the items in a list/array?"
parameters - variables with unspecified values
problem instance - assignment of values to parameters, i.e., the specific input to the problem

\[
\begin{array}{cccccc}
0 & 1 & 2 & 3 & 4 & 5 & 6 \\
\text{myList:} & 5 & 10 & 2 & 15 & 20 & 1 & 11 \\
\text{(number of elements)} & n: & 7 & \\
\text{sum: ?}
\end{array}
\]

algorithm - step-by-step procedure for producing a solution
basic operation - fundamental operation in the algorithm (i.e., operation done the most) Generally, we want to derive a function for the number of times that the basic operation is performed related to the problem size.
problem size - input size. For algorithms involving lists/arrays, the problem size is the number of elements ("n").

Big-oh notation (O()) - As the size of a problem grows (i.e., more data), how will our program's run-time grow.
Consider the following sumList function.

```python
def sumList(myList):
    """Returns the sum of all items in myList""
    total = 0
    for item in myList:
        total = total + item
    return total
```

a) What is the basic operation of sumList (i.e., operation done the most)?

b) What is the problem size of sumList?

c) If n is 10000 and sumList takes 10 seconds, how long would you expect sumList to take for n of 20000?

d) What is the big-oh notation for sumList?

4. Consider the following someLoops function.

```python
def someLoops(n):
    total = 0
    for i in range(n):
        for j in range(n):
            total = total + i + j
    return total
```

a) What is the basic operation of someLoops (i.e., operation done the most)?

b) How many times will the basic operation execute as a function of n?

c) What is the big-oh notation for someLoops?

d) If we input n of 10000 and someLoops takes 10 seconds, how long would you expect someLoops to take for n of 20000?
1. Draw the graph for \texttt{sumList} \((O(n))\) and \texttt{sumLoops} \((O(n^2))\) from the previous lecture.

\begin{figure}[h]
\centering
\includegraphics[width=0.5\textwidth]{graph.png}
\caption{Execution Time vs. \(n\)}
\end{figure}

2. Consider the following \texttt{sumSomeListItems} function.

```python
import time

def main():
    n = eval(input("Enter size of list: "))
    alist = list(range(1, n+1))
    start = time.clock()
    sum = sumSomeListItems(alist)
    end = time.clock()
    print("Time to sum the list was \%.9f seconds" % (end-start))

def sumSomeListItems(myList):
    """Returns the sum of some items in myList""
    total = 0
    index = len(myList) - 1
    while index > 0:
        total = total + myList[index]
        index = index // 2
    return total

main()
```

a) What is the problem size of \texttt{sumSomeListItems}?

b) If we input \(n\) of 10,000 and \texttt{sumSomeListItems} takes 10 seconds, how long would you expect \texttt{sumSomeListItems} to take for \(n\) of 20,000?

(Hint: For \(n\) of 20,000, how many more times would the loop execute than for \(n\) of 10,000?)

c) What is the big-oh notation for \texttt{sumSomeListItems}?

d) Add the execution-time graph for \texttt{sumSomeListItems} to the graph.
3.

```python
i = 1
while i <= n:
    for j in range(n):
        # something of O(1)
    # end for
    i = i * 2
# end while
```
a) Analyze the above algorithm to determine its big-oh notation, $O()$.

b) If $n$ of 10,000, takes 10 seconds, how long would you expect the above code to take for $n$ of 20,000?

c) Add the execution-time graph for the above code to the graph.

4. Most programming languages have a built-in array data structure to store a collection of same-type items. Arrays are implemented in RAM memory as a contiguous block of memory locations. Consider an array $X$ that contains the odd integers:

<table>
<thead>
<tr>
<th>address</th>
<th>Memory</th>
</tr>
</thead>
<tbody>
<tr>
<td>4000</td>
<td>1</td>
</tr>
<tr>
<td>4004</td>
<td>3</td>
</tr>
<tr>
<td>4008</td>
<td>5</td>
</tr>
<tr>
<td>4012</td>
<td>7</td>
</tr>
<tr>
<td>4016</td>
<td>9</td>
</tr>
<tr>
<td>4020</td>
<td>11</td>
</tr>
<tr>
<td>4024</td>
<td>13</td>
</tr>
</tbody>
</table>
| ...     | ...    | ...

a) Any array element can be accessed randomly by calculating its address. For example, address of $X[5] = 4000 + 5 * 4 = 4020$. What is the general formula for calculating the address of the $i$th element in an array?

b) What is the big-oh notation for accessing the $i$th element?

c) A Python list uses an array of references (pointers) to list items in their implementation of a list. For example, a list of strings containing the alphabet:

```
0 1 2 3
\[\downarrow \downarrow \downarrow \downarrow \downarrow \downarrow\]
'a' 'b' 'c' 'd' \[\downarrow \downarrow \downarrow \downarrow \downarrow \downarrow\]

(len()-1)
```

Since a Python list can contain heterogeneous data, how does storing references in the list aid implementation?
5. Arrays in most HLLs are static in size (i.e., cannot grow at run-time), so arrays are constructed to hold the “maximum” number of items. For example, an array with 1,000 slots might only contain 3 items:

| size: 3 | scores: 20 | 10 | 30 | 999 |

a) The physical size of the array is the number of slots in the array. What is the physical size of scores?

b) The logical size of the array is the number of items actually in the array. What is the logical size of scores?

c) The load factor is fraction of the array being used. What is the load factor of scores?

d) What is the $O(\cdot)$ for “appending” a new score to the “right end” of the array?

e) What is the $O(\cdot)$ for adding a new score to the “left end” of the array?

f) What is the average $O(\cdot)$ for adding a new score to the array?

g) During run-time if an array fills up and we want to add another item, the program can usually:
   - Create a bigger array than the one that filled up
   - Copy all the items from the old array to the bigger array
   - Add the new item
   - Delete the smaller array to free up its memory
When creating the bigger array, how much bigger than the old array should it be?

h) What is the $O(\cdot)$ of moving to a larger array?

6. Consider the following list methods in Python:

<table>
<thead>
<tr>
<th>Method</th>
<th>Usage</th>
<th>Average $O(\cdot)$ for myList containing (n) items</th>
</tr>
</thead>
<tbody>
<tr>
<td>index []</td>
<td>itemValue = myList[i] \n myList[i] = newValue</td>
<td></td>
</tr>
<tr>
<td>append</td>
<td>myList.append(item)</td>
<td></td>
</tr>
<tr>
<td>extend</td>
<td>myList.extend(otherList)</td>
<td></td>
</tr>
<tr>
<td>insert</td>
<td>myList.insert(i, item)</td>
<td></td>
</tr>
<tr>
<td>pop</td>
<td>myList.pop()</td>
<td></td>
</tr>
<tr>
<td>pop(i)</td>
<td>myList.pop(i)</td>
<td></td>
</tr>
<tr>
<td>del</td>
<td>del myList[i]</td>
<td></td>
</tr>
<tr>
<td>remove</td>
<td>myList.remove(item)</td>
<td></td>
</tr>
<tr>
<td>index</td>
<td>myList.index(item)</td>
<td></td>
</tr>
<tr>
<td>iteration</td>
<td>for item in myList:</td>
<td></td>
</tr>
<tr>
<td>reverse</td>
<td>myList.reverse( )</td>
<td></td>
</tr>
</tbody>
</table>

Dictionary Operations:

<table>
<thead>
<tr>
<th>Method</th>
<th>Usage</th>
<th>Explanation</th>
<th>Average $O(\cdot)$ for (n) keys</th>
</tr>
</thead>
<tbody>
<tr>
<td>get item</td>
<td>myDictionary.get(myKey) \n value = myDictionary[myKey]</td>
<td>Returns the value associated with myKey; otherwise None</td>
<td>(O(1))</td>
</tr>
<tr>
<td>set item</td>
<td>myDictionary[myKey] = value</td>
<td>Change or add myKey:value pair</td>
<td>(O(1))</td>
</tr>
<tr>
<td>in</td>
<td>myKey in myDictionary</td>
<td>Returns True if myKey is in myDictionary; otherwise False</td>
<td>(O(1))</td>
</tr>
<tr>
<td>del</td>
<td>del myDictionary[myKey]</td>
<td>Deletes the mykey:value pair</td>
<td>(O(1))</td>
</tr>
</tbody>
</table>
Objective: To get a feel for big-oh notation by analyzing algorithms as well as timing them (Part A), and gain some experience writing Python classes.

Informal Big-oh (and Big-Theta) Definition: As the size of a computational problem grows (i.e., more data), we expect our program to run longer, but this run-time growth is not necessarily linear. Big-oh notation gives us an idea how our program’s run-time will grow with respect to its problem size on larger data.

This might seem like a lot of mathematical mumbo-jumbo, but knowing an algorithms big-oh notation can help us predict its run-time on large problem sizes. While running a large size problem, we might want to know if we have time for a quick lunch, a long lunch, a long nap, go home for the day, take a week of vacation, pack-up the desk because the boss will fire you for a slow algorithm, etc.

For example, consider the following algorithm:

```python
class = 0
for r in range(n):
    for c in range(n):
        for d in range(n/2):
            result = result + d
    # end for
# end for
```

Clearly, the body of the inner-most loop (the “result = result + d” statement) will execute n^3/2 times, so this algorithm is “big-oh” of n-cubed, O(n^3). Thus, the execution-time formula, T(n), with respect to n is:

```
T(n) = c n^3 + (slower growing terms).
```

For large values of n, the execution time as a function of n, T(n) \propto c n^3, where c is the constant of proportionality on the fastest growing term (the machine dependent time related to how long it takes to execute the inner-most loop once). If we know that T(10,000) = 1 second = c \times 10,000^3, then we can calculate c and use it to predict what T(1,000,000). First approximate c as c \approx T(n) / n^3 = 1 second / 10,000^3 = 1 second / 10^{12} = 10^{-12} seconds. Since we are running the algorithm on the same machine, c is unchanged for the larger problem. Thus, T(1,000,000) \approx c \times 1,000,000^3 = c \times 10^{18} = 10^{-12} seconds * 10^{18} = 10^6 seconds or about 11.6 days. (A couple weeks of vacation is appropriate!)

To start the lab: Download and unzip the file lab2.zip from [http://www.cs.uni.edu/~jienup/cs1520s19/labs/lab2.zip](http://www.cs.uni.edu/~jienup/cs1520s19/labs/lab2.zip)

Part A: In the folder lab2, open the timeStuff.py program in IDLE. (Right-click on timeStuff.py | Edit with IDLE) Start it running in IDLE by selecting Run | Run Module from the menu. While it is running, answer the following questions about each of the algorithms in timeStuff.py.

a) What is the big-oh of Algorithm 0?

Algorithm 0:

```python
result = 0
for i in range(10000000):
    result = result + i
```

b) What is the big-oh of Algorithm 1?

Algorithm 1:

```python
result = 0
for i in range(n):
    result = result + i
# end for
```
Data Structures (CS 1520)

Lab 2

Name: ______________________

c) What is the big-oh of Algorithm 2?

Algorithm 2:
result = 0
for r in range(n):
    c = n
    while c > 1:
        result = result + c
        c = c // 2
    # end while
# end for

d) What is the big-oh of Algorithm 3?

Algorithm 3:
result = 0
for r in range(n):
    for c in range(n):
        result = result + c
    # end for
# end for

e) What is the big-oh of Algorithm 4?

Algorithm 4:
result = 0
for r in range(n):
    for c in range(n):
        for d in range(n*n*n):
            result = result + d
    # end for
# end for

f) What is the big-oh of Algorithm 5?

Algorithm 5:
result = 0
i = 0
while i < 2**n:
    result = result + i
    i += 1
# end while
g) Complete the following timing table from the output of timeStuff.py.

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>n = 0</th>
<th>n = 10</th>
<th>n = 20</th>
<th>n = 30</th>
<th>n = 40</th>
<th>n = 50</th>
</tr>
</thead>
<tbody>
<tr>
<td>Algorithm 0</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Algorithm 1</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Algorithm 2</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Algorithm 3</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Algorithm 4</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Algorithm 5</td>
<td></td>
<td></td>
<td></td>
<td>(work on parts h - j while waiting)</td>
<td></td>
<td>These take too long!</td>
</tr>
</tbody>
</table>

h) For Algorithm 5, use the timing for n = 20 to compute the constant of proportionality, c, on the fastest growing term.

i) Using the constant of proportionality computed in (h), predict the run-time of Algorithm 5 for n = 30.

j) How does your prediction in (i) compare to the actual time from (g)?

After you have answered the above questions, raise your hand and explain your answers.

(Note: Part B is on the backside of this sheet)
Data Structures (CS 1520)  Lab 2  Name:____________________

**Part B:** The `lab2.zip` file also contains:
- A simple `Die` class (in the `simple_die.py` module) for a six-sided die.
- An `AdvancedDie` class (in the module `advanced_die.py` module) for a die which can be constructed with any number of sides. The `AdvancedDie` class inherits from the `Die` class.
- An `averageOutcome.py` program that computes the average outcome (i.e., average total on the pair of dice) on a pair of 10-side dice over 1,000 rolls. Unfortunately, it uses `randint(1,10)+randint(1,10)`.

a) Modify the `averageOutcome.py` program so that it uses the `AdvancedDie` class by:
- creating two 10-sided `AdvancedDie` objects (remember to "from advanced_die import AdvancedDie")
- rolls the pair of dice 1,000 times to compute the average outcome
(Note: most of the program will remain unchanged)

**Part C:**
a) For testing certain dice games, suppose we want to extend the `AdvancedDie` class to include a new method `setRoll` which takes as a parameter a roll value that is used to set a die’s roll to a specified value. We might invoke this method as:

```python
myDie.setRoll(3)  # sets myDie's current roll to 3
```

**Implement a new subclass** `MoreAdvancedDie` (in a new file `more_advanced_die.py`) which inherits from the `AdvancedDie` class, and includes the new `setRoll` method. When implementing the `MoreAdvancedDie` class:
- Include documentation with the `MoreAdvancedDie` class (comments right below it’s class line)
- Include documentation with the `setRoll` method including preconditions and postconditions
- Enforce the `setRoll` method’s preconditions by raising appropriate exceptions.
(see the `AdvancedDie` class and its methods for examples of raising exceptions)

b) View the programmer-authored documentation for the `MoreAdvancedDie` class by typing `help(MoreAdvancedDie)` at the IDLE shell prompt (">>>") after selecting Run | Run Module in the file `more_advanced_die.py` which contains the `MoreAdvancedDie` class.

After you have implemented AND fully tested your `MoreAdvancedDie` class, raise you hand and demonstrate it.

If you complete all parts of the lab, nothing needs to be turned in for this lab. If you do not get done today, then show me the completed lab in next week’s lab period. When done save your lab 2 files (USB drive, etc.) and remember to log off.

If you have extra time, this would be a good chance to work on Homework #1!
1. An “abstract” view of the stack:

Using an array implementation would look something like:

<table>
<thead>
<tr>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>(max-1)</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>b</td>
<td>c</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

top: 2 max: 100

Complete the big-oh notation for the following stack methods assuming an array implementation: ("n" is the # items)

<table>
<thead>
<tr>
<th>push(item)</th>
<th>pop()</th>
<th>peek()</th>
<th>size()</th>
<th>isEmpty()</th>
<th>isFull()</th>
<th>Constructor</th>
</tr>
</thead>
<tbody>
<tr>
<td>Big-oh</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

2. Since Python does not have a (directly accessible) built-in array, we can use a list.

```python
class Stack:
    def __init__(self):
        self.items = []

    def isEmpty(self):
        return self.items == []

    def push(self, item):
        self.items.append(item)

    def pop(self):
        return self.items.pop()

    def peek(self):
        return self.items[len(self.items)-1]

    def size(self):
        return len(self.items)
```

Since Python uses an array of references (pointers) to list items in their implementation of a list.

"Abstract" Stack Object list Object

```
<p>| | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>c</td>
<td></td>
<td>a</td>
<td></td>
</tr>
<tr>
<td>top</td>
<td></td>
<td>list</td>
<td>0 1 2 3</td>
</tr>
<tr>
<td>b</td>
<td></td>
<td>b</td>
<td></td>
</tr>
<tr>
<td>a</td>
<td></td>
<td>c</td>
<td></td>
</tr>
<tr>
<td>bottom</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
```

a) Complete the big-oh notation for the stack methods assuming this Python list implementation: ("n" is the # items)

<table>
<thead>
<tr>
<th>push(item)</th>
<th>pop()</th>
<th>peek()</th>
<th>size()</th>
<th>isEmpty()</th>
<th><strong>init</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td>Big-oh</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

b) Which operations should have what preconditions?
3. The text's alternative stack implementation also using a Python list is:

```python
class Stack:
    def __init__(self):
        self.items = []

    def isEmpty(self):
        return self.items == []

    def push(self, item):
        self.items.insert(0, item)

    def pop(self):
        return self.items.pop(0)

    def peek(self):
        return self.items[0]

    def size(self):
        return len(self.items)
```

Since an array is used to implement a Python list, the alternate Stack implementation using a list:

<table>
<thead>
<tr>
<th>&quot;Abstract&quot; Stack</th>
<th>&quot;alternate&quot; Stack Object</th>
<th>list Object</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>items:</td>
<td>0 1 2 3</td>
</tr>
<tr>
<td>top</td>
<td></td>
<td>c b a</td>
</tr>
<tr>
<td>c</td>
<td></td>
<td></td>
</tr>
<tr>
<td>b</td>
<td></td>
<td></td>
</tr>
<tr>
<td>a</td>
<td>bottom</td>
<td></td>
</tr>
</tbody>
</table>

a) Complete the big-oh notation for the "alternate" Stack methods: ("n" is the # items)

<table>
<thead>
<tr>
<th></th>
<th>push(item)</th>
<th>pop()</th>
<th>peek()</th>
<th>size()</th>
<th>isEmpty()</th>
<th><strong>init</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td>Big-oh</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

4. How could we use a stack to check if a word is a palindrome (e.g., radar, toot)?

5. How could we check to see if we have a balanced string of nested symbols? ("(()[]) {} []")
1. The Node class (in node.py) is used to dynamically create storage for a new item added to the stack. The LinkedStack class (in linked_stack.py) uses this Node class. Conceptually, a LinkedStack object would look like:

```
class Node:
    def __init__(self, initdata):
        self.data = initdata
        self.next = None
    def getData(self):
        return self.data
    def getNext(self):
        return self.next
    def setData(self, newdata):
        self.data = newdata
    def setNext(self, newnext):
        self.next = newnext

class LinkedStack(object):
    """Link-based stack implementation."""

    def __init__(self):
        self._top = None
        self._size = 0
    def push(self, newItem):
        """Inserts newItem at top of stack."""
        self._top = newItem
        self._size += 1
    def pop(self):
        """Removes and returns the item at top of the stack.
        Precondition: the stack is not empty."""
        item = self._top
        self._top = self._top.next
        self._size -= 1
        return item
    def peek(self):
        """Returns the item at top of the stack.
        Precondition: the stack is not empty."""
        return self._top.getData()
    def size(self):
        """Returns the number of items in the stack."""
        return self._size
    def isEmpty(self):
        return self._size == 0
    def __str__(self):
        """Items strung from top to bottom."""
```

a) Complete the push, pop, and __str__ methods.

b) Stack methods big-oh’s?
( Assume ”n” items in stack)
- constructor __init__:
- push(item):
- pop()
- peek()
- size()
- isEmpty()
- __str__()}
A FIFO queue is basically what we think of as a waiting line.

The operations/methods on a queue object, say myQueue are:

<table>
<thead>
<tr>
<th>Method Call on myQueue object</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>myQueue.dequeue()</td>
<td>Removes and returns the front item in the queue.</td>
</tr>
<tr>
<td>myQueue.enqueue(myItem)</td>
<td>Adds myItem at the rear of the queue</td>
</tr>
<tr>
<td>myQueue.peek()</td>
<td>Returns the front item in the queue without removing it.</td>
</tr>
<tr>
<td>myQueue.isEmpty()</td>
<td>Returns True if the queue is empty, or False otherwise.</td>
</tr>
<tr>
<td>myQueue.size()</td>
<td>Returns the number of items currently in the queue</td>
</tr>
<tr>
<td>str(myQueue)</td>
<td>Returns the string representation of the queue</td>
</tr>
</tbody>
</table>

2. Complete the following table by indicating which of the queue operations should have preconditions. Write "none" if a precondition is not needed.

<table>
<thead>
<tr>
<th>Method Call on myQueue object</th>
<th>Precondition(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>myQueue.dequeue()</td>
<td></td>
</tr>
<tr>
<td>myQueue.enqueue(myItem)</td>
<td></td>
</tr>
<tr>
<td>myQueue.peek()</td>
<td></td>
</tr>
<tr>
<td>myQueue.isEmpty()</td>
<td></td>
</tr>
<tr>
<td>myQueue.size()</td>
<td></td>
</tr>
<tr>
<td>str(myQueue)</td>
<td></td>
</tr>
</tbody>
</table>

3. The textbook's Queue implementation uses a Python list:

```python
class Queue:
    def __init__(self):
        self.items = []
    def isEmpty(self):
        return self.items == []
    def enqueue(self, item):
        self.items.insert(0, item)
    def dequeue(self):
        return self.items.pop()
    def peek(self):

    def size(self):
        return len(self.items)
    def __str__(self):
```

a) Complete the `peek`, and `str` methods

b) What are the Queue methods big-oh's?
   (Assume "n" items in the queue)
   - constructor `__init__`
   - `isEmpty`
   - `enqueue(item)`
   - `dequeue`
   - `peek`
   - `size`
   - `str`
3. An alternate queue implementation using a linked structure (LinkedQueue class) would look like:

```
"Abstract Queue"
  'w' 'x' 'y'
  front  rear

LinkedQueue Object
  _front: 
  _size: 3
  _rear: 

Node Objects
  data 'w' next
   data 'x' next
   data 'y' next
```

a) Draw on the picture and number the steps for the enqueue method of the "normal" case (non-empty queue)

b) Write the enqueue method code for the "normal" case:

```
// Code for enqueue method
LinkedQueue.enqueue(item) {
  // Implementation
}
```

c) Starting with the empty queue below, draw the resulting picture after your "normal" case code executes.

```
empty LinkedQueue Object
  _front:
  _size: 0
  _rear:
```

d) Fix your "normal" case code to handle the "special case" of an empty queue.
e) Draw on the above picture and number the steps for the `dequeue` method of the "normal" case (non-empty queue).

f) Write the `dequeue` method code for the "normal" case:

```
// Code goes here
```

g) What "special case(s)" does the `dequeue` method code need to handle?

h) Draw the picture for each special case and number the steps for the `dequeue` method in the "special" case(s)

```
// Picture goes here
```

i) Combine the "normal" and special case(s) code for a complete `dequeue` method.

```
// Combined code goes here
```

j) Complete the big-oh notation for the `LinkedQueue` methods: ("n" is the # items)

<table>
<thead>
<tr>
<th></th>
<th><code>init</code></th>
<th>enqueue(item)</th>
<th>dequeue()</th>
<th>peek()</th>
<th>size()</th>
<th>isEmpty()</th>
<th><strong>str</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td>Big-oh</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Data Structures (CS 1520)  Lab 3  Name: 

Objective: To understand FIFO (First-In-First-Out) queue implementations in Python including being able to determine the big-oh of each operation.

To start the lab: Download and unzip the file at: [www.cs.uni.edu/~sienup/cs1520s19/labs/lab3.zip](http://www.cs.uni.edu/~sienup/cs1520s19/labs/lab3.zip)

Part A: The textbook's QueueText implementation in lab3/queue_text.py uses a Python list

```
"Abstract Queue"

front  'w'  'x'  'y'  rear

```

<table>
<thead>
<tr>
<th>Callback</th>
<th><em>init</em></th>
<th>enqueue(item)</th>
<th>dequeue()</th>
<th>peek()</th>
<th>size()</th>
<th>isEmpty()</th>
<th><strong>str</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td>Big-oh</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

a) Complete the big-oh notation for the above QueueText implementation: ("n" is the # items)

b) Explain your big-oh answer for enqueue(item).

c) Explain your big-oh answer for dequeue()

d) Run the timeQueue.py file which times 100,000 enqueues followed by 100,000 dequeues.
Time for 100,000 enqueues:  
Time for 100,000 dequeues:

e) Why do the enqueues take so much more time?

After answering the above questions, raise your hand and explain your answers.

Part B:

a) Complete the QueueAlt implementation in lab3/queue_alt.py uses a Python list

```
"Abstract Queue"

front  'w'  'x'  'y'  rear

```

<table>
<thead>
<tr>
<th>Callback</th>
<th><em>init</em></th>
<th>enqueue(item)</th>
<th>dequeue()</th>
<th>peek()</th>
<th>size()</th>
<th>isEmpty()</th>
<th><strong>str</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td>Big-oh</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

b) Complete the big-oh notation for the above QueueAlt implementation: ("n" is the # items)

<table>
<thead>
<tr>
<th>Callback</th>
<th><em>init</em></th>
<th>enqueue(item)</th>
<th>dequeue()</th>
<th>peek()</th>
<th>size()</th>
<th>isEmpty()</th>
<th><strong>str</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td>Big-oh</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

c) Run the timeQueueAlt.py file which times 100,000 enqueues followed by 100,000 dequeues.
Time for 100,000 enqueues:  
Time for 100,000 dequeues:

After completing the QueueAlt class, answering the above questions, raise your hand and demonstrate your code.
**Part C:** Consider the `LinkedQueue` implementation in `lab3/linked_queue.py` which uses a linked structure that looks like:

![LinkedQueue Diagram]

a) Modify the above picture and number the steps for the `enqueue` method's "normal" case (non-empty queue)

b) Complete the `enqueue` method code for the "normal" case in the `lab3/linked_queue.py` file

c) Starting with the empty queue below, draw the resulting picture after your "normal" case code executes.

```
empty LinkedQueue Object
```

d) Fix your "normal" case code to handle the "special case" of an empty queue.

```
LinkedQueue Object
```

e) Modify the above picture and number the steps for the `dequeue` method's "normal" case (non-empty queue)

f) Complete the `dequeue` method code for the "normal" case in the `lab3/linked_queue.py` file

g) What "special case(s)" does the `dequeue` method code need to handle?

h) Complete the big-oh notation for the `LinkedQueue` methods: ("n" is the # items)

<table>
<thead>
<tr>
<th></th>
<th><strong>init</strong></th>
<th>enqueue(item)</th>
<th>dequeue()</th>
<th>peek()</th>
<th>size()</th>
<th>isEmpty()</th>
<th><strong>str</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td>Big-oh</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

i) Run the `timeLinkedQueue.py` file which times 100,000 enqueues followed by 100,000 dequeues.

Time for 100,000 enqueues:

Time for 100,000 dequeues:

After thoroughly testing your linked implementation, raise you hand and demonstrate your queue.
A Deque (pronounced "Deck") is a linear data structure which behaves like a double-ended queue, i.e., it allows adding or removing items from either the front or the rear of the Deque.

1. One possible implementation of a Deque would be to use a Python list to store the Deque items such that
   - the rear item is always stored at index 0,
   - the front item is always stored at the highest index (or -1)

   Deque Object
   items: [ ]
   List Object
   0 1 2 3
   'a' 'b' 'c' 'd'
   (rear) (front)

   a) Complete the __init__ method and determine the big-oh, $O()$, for each Deque operation, assuming the above implementation. Let $n$ be the number of items in the Deque.

<table>
<thead>
<tr>
<th>Method</th>
<th>Time Complexity</th>
</tr>
</thead>
<tbody>
<tr>
<td>isEmpty</td>
<td>$O(1)$</td>
</tr>
<tr>
<td>addFront</td>
<td>$O(n)$</td>
</tr>
<tr>
<td>removeFront</td>
<td>$O(n)$</td>
</tr>
<tr>
<td>addRear</td>
<td>$O(n)$</td>
</tr>
<tr>
<td>removeRear</td>
<td>$O(n)$</td>
</tr>
</tbody>
</table>

   b) Write the methods for the addRear and removeRear operation.

   ```python
   def addRear(self, newItem):
   def removeRear(self):
   ```

2. An alternative implementation of a Deque would be a linked implementation as in:

   LinkedDeque Object
   _rear: -------
   _front: -------
   _size: 4

   data next data next data next
   'a' -> 'b' -> 'c' -> 'd'

   a) Complete the __init__ method and determine the big-oh, $O()$, for each Deque operation assuming the above linked implementation. Let $n$ be the number of items in the Deque.

<table>
<thead>
<tr>
<th>Method</th>
<th>Time Complexity</th>
</tr>
</thead>
<tbody>
<tr>
<td>isEmpty</td>
<td>$O(1)$</td>
</tr>
<tr>
<td>addFront</td>
<td>$O(n)$</td>
</tr>
<tr>
<td>removeFront</td>
<td>$O(n)$</td>
</tr>
<tr>
<td>addRear</td>
<td>$O(n)$</td>
</tr>
<tr>
<td>removeRear</td>
<td>$O(n)$</td>
</tr>
</tbody>
</table>

   b) Suggest an improvement to the above linked implementation of the Deque to speed up some of its operations.
3. An alternative implementation of a Deque would be a doubly-linked implementation as in:

```python
from node import Node

class Node2Way(Node):
    def __init__(self, initdata):
        Node.__init__(self, initdata)
        self.previous = None

    def getPrevious(self):
        return self.previous

    def setPrevious(self, newprevious):
        self.previous = newprevious
```

DoublyLinkedDeque Object

3. An alternative implementation of a Deque would be a doubly-linked implementation as in:

DoublyLinkedDeque Object

- rear: ____________
- front: ____________
- size: _________

a) Determine the big-oh, $O()$, for each Deque operation assuming the above doubly-linked implementation. Let $n$ be the number of items in the Deque.

<table>
<thead>
<tr>
<th>isEmpty</th>
<th>addFront</th>
<th>removeFront</th>
<th>addRear</th>
<th>removeRear</th>
<th>size</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

4. A **priority queue** has the same operations as a regular queue, except the items are NOT returned in the FIFO (first-in, first-out) order. Instead, each item has a priority that determines the order they are removed. A hospital emergency room operates like a priority queue -- the person with the most serious injury has highest priority even if they just arrived.

a) Suppose that we have a priority queue with integer priorities such that the smallest integer corresponds to the highest priority. For the following priority queue, which item would be dequeued next?

```
priority queue:

40  10  79
|      |      |
|      |      |
| 30   |      |
|      |      |
```

b) To implement a priority queue, we could use an **unordered Python list**. If we did, what would be the big-oh notation for each of the following methods: (justify your answer)

- enqueue:
- dequeue:

c) To implement a priority queue, we could use a **Python list order by priorities** in descending order. If we did, what would be the big-oh notation for each of the following methods: (justify your answer)

- enqueue:
- dequeue
1. Section 6.6 discusses a very "non-intuitive", but powerful list/array-based approach to implement a priority queue, call a binary heap. The list/array is used to store a complete binary tree (a full tree with any additional leaves as far left as possible) with the items being arranged by heap-order property, i.e., each node is ≤ either of its children. An example of a min heap "viewed" an a complete binary tree would be:

![Binary Heap Diagram]

Python List actually used to store heap items

<table>
<thead>
<tr>
<th></th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
<th>8</th>
<th>9</th>
<th>10</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>6</td>
<td>15</td>
<td>10</td>
<td>114</td>
<td>20</td>
<td>20</td>
<td>50</td>
<td>300</td>
<td>125</td>
<td>117</td>
<td></td>
</tr>
</tbody>
</table>

a) For the above heap, the list/array indexes are indicated in [ ]'s. For a node at index i, what is the index of:
- its left child if it exists:
- its right child if it exists:
- its parent if it exists:

b) What would the above heap look like after inserting 13 and then 37? (show the changes on above tree)

General Idea of insert(newItem):
- append newItem to the end of the list (easy to do, but violates heap-order property)
- restore the heap-order property by repeatedly swapping the newItem with its parent until it percolates to correct spot

c) What is the big-oh notation for inserting a new item in the heap?

d) Complete the code for the percUp method used by insert.

```python
class BinHeap:
    def __init__(self):
        self.heapList = [0]
        self.currentSize = 0

    def percUp(self, currentIdx):
        parentIdx =
        while

    def insert(self,k):
        self.heapList.append(k)
        self.currentSize = self.currentSize + 1
        self.percUp(self.currentSize)
```
2. Now let us consider the \texttt{delMin} operation that removes and returns the minimum item.

![Heap Diagram]

Python List actually used to store heap items: 

<table>
<thead>
<tr>
<th></th>
<th>6</th>
<th>15</th>
<th>10</th>
<th>114</th>
<th>20</th>
<th>20</th>
<th>50</th>
<th>300</th>
<th>125</th>
<th>117</th>
</tr>
</thead>
</table>

a) What item would \texttt{delMin} remove and return from the above heap?

b) What is the quickest way to fill the hole left by \texttt{delMin}?

c) What new problem does this cause?

**General Idea of \texttt{delMin}()**:

- Remember the minimum value so it can be returned later (easy to find - at index 1)
- Copy the last item in the list to the root, delete it from the right end, decrement size
- Restore the heap-order property by repeatedly swapping this item with its smallest child until it \textit{percolates down} to the correct spot
- Return the minimum value

d) What would the above heap look like after \texttt{delMin}? (show the changes on above tree)

e) Complete the code for the \texttt{percDown} method used by \texttt{delMin}.

```python
class BinHeap:
    
    def minChild(self, i):
        if i*2+1 < self.currentSize: # if only left child
            return i*2
        else:
            if self.heapList[i*2] < self.heapList[i*2+1]:
                return i*2+
            else:
                return i*2+1

    def delMin(self):
        retval = self.heapList[1]
        self.currentSize = self.currentSize - 1
        self.heapList[1] = self.percDown(1)
        return retval

def percDown(self, currentIndex):
    return
```

f) What is the big-oh notation for \texttt{delMin}?
Data Structures (CS 1520)

Once we have a working BinHeap, then implementing the PriorityQueue class using a BinHeap is a piece of cake:

```python
>>> q = PriorityQueue()
>>> print(q)
[]
>>> q.enqueue(5)
>>> q.enqueue(1)
>>> q.enqueue(7)
>>> print(q)
[1, 5, 7]
>>> q.dequeue()
1
>>> print(q)
[5, 7]
```

3. A “list” is a generic term for a sequence of items in a linear arrangement. Unlike stacks, queues and deques access to list items is not limited to either end, but can be from any position in the list. The general terminology of a list is illustrated by:

```
"Abstract view of a list" [0] [1] [2] [3] ← index/position in the list
\[ \begin{array}{c}
\uparrow \\
\text{\textbackslash w} \\
\text{\textbackslash a} \\
\text{\textbackslash y} \\
\text{\textbackslash c} \\
\text{\textbackslash head} \\
\text{\textbackslash tail} \\
\end{array} \]
```

There are three broad categories of list operations that are possible:
- **index-based operations** - the list is manipulated by specifying an index location, e.g.,
  myList.insert(3, item)  # insert item at index 3 in myList
- **content-based operations** - the list is manipulated by specifying some content (i.e., item value), e.g.,
  myList.remove(item)  # removes the item from the list based on its value
- **cursor-based operations** - a cursor (current position) can be moved around the list, and it is used to identify list items to be manipulated, e.g.,
  myList.first()  # sets the cursor to the head item of the list
  myList.next()  # moves the cursor one position toward the tail of the list
  myList.remove()  # deletes the second item in the list because that’s where the cursor is currently located

The following table summarizes the operations from the three basic categories on a list, L:

<table>
<thead>
<tr>
<th>Index-based operations</th>
<th>Content-based operations</th>
<th>cursor-based operations</th>
</tr>
</thead>
<tbody>
<tr>
<td>L.insert(index, item)</td>
<td>L.add(item)</td>
<td>L.hasNext()</td>
</tr>
<tr>
<td>item = L[index]</td>
<td>L.remove(item)</td>
<td>L.next()</td>
</tr>
<tr>
<td>L[index] = newValue</td>
<td>L.search(item)</td>
<td>L.hasPrevious()</td>
</tr>
<tr>
<td>L.pop(index)</td>
<td>i = L.index(item)</td>
<td>L.previous()</td>
</tr>
<tr>
<td></td>
<td></td>
<td>L.first()</td>
</tr>
<tr>
<td></td>
<td></td>
<td>L.last()</td>
</tr>
<tr>
<td></td>
<td></td>
<td>L.insert(item)</td>
</tr>
<tr>
<td></td>
<td></td>
<td>L.replace(item)</td>
</tr>
<tr>
<td></td>
<td></td>
<td>L.remove()</td>
</tr>
</tbody>
</table>

Built-in Python lists are unordered with a mixture of index-based and content-based operations. We know they are implemented using a contiguous block of memory (i.e., an array). The textbook talks about an unordered list ADT, and a sorted list ADT which is more content-based. Both are implemented using a singly-linked list.

a) Why would a singly-linked list be a bad choice for implementing a cursor-based list ADT?
Objective: To understand priority queue implementations in Python including being able to determine the big-oh of each operation.

To start the lab: Download and unzip the file at: www.cs.unl.edu/~finnup/cs1520s19/labs/lab4.zip

Part A:

a) Suppose that we have a priority queue with integer priorities such that the smallest integer corresponds to the highest priority. For the following priority queue, which item would be dequeued next?

```
priority queue:
  40 10  79
  30 13  5
```

b) The `ListPriorityQueue` implementation in `lab4/list_priority_queue.py` uses an unordered Python list.

```
ListPriorityQueue Object
  _items: [ ]

List Object
  0 1 2 3 4 5
  40 30 10 13 5 79
```

What would be the big-oh notation for each of the following methods: (justify your answer)
- enqueue:
- dequeue:

c) The `SortedListPriorityQueue` implementation in `lab4/sorted_list_priority_queue.py` uses a Python list order by priorities in descending order.

```
SortedListPriorityQueue Object
  _items: [ ]

List Object
  0 1 2 3 4 5
  79 40 30 13 10 5
```

What would be the big-oh notation for each of the following methods: (justify your answer)
- enqueue:
- dequeue

d) Why would it be a bad idea to implement a priority queue using a Python list order by priorities in reverse (ascending) order? (HINT: What is the big-oh notations for enqueue and dequeue?)

```
SortedListPriorityQueue Object
  _items: [ ]

List Object
  0 1 2 3 4 5
  5 10 13 30 40 79
```

Answer the above questions, then raise your hand. Explain your answers to an instructor or TA.
Part B: (Lecture 7 and) Section 6.6 discusses a very “non-intuitive”, but powerful list/array-based approach to implement a priority queue, call a binary heap. The list/array is used to store a complete binary tree (a full tree with any additional leaves as far left as possible) with the items being arranged by heap-order property, i.e., each node is ≤ either of its children. An example of a min heap "viewed" an a complete binary tree would be:

```
  10
 /   \
 25   15
 / \
 40  30
 / \
50 85
```

Python List actually used to store heap items:

<table>
<thead>
<tr>
<th></th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
<th>8</th>
<th>9</th>
<th>10</th>
</tr>
</thead>
<tbody>
<tr>
<td>10</td>
<td></td>
<td></td>
<td>25</td>
<td>15</td>
<td>50</td>
<td>40</td>
<td>30</td>
<td>110</td>
<td>300</td>
<td>85</td>
<td>45</td>
</tr>
</tbody>
</table>

a) For the above heap, the list/array indexes are indicated in [ ]’s. For a node at index i, what is the index of:

- its left child if it exists:
- its right child if it exists:
- its parent if it exists:

Recall the General Idea of insert(newItem):

- append newItem to the end of the list (easy to do, but violates heap-order property)
- restore the heap-order property by repeatedly swapping the newItem with its parent until it percolates up to the correct spot

b) What would the above heap look like after inserting 18 and then 27? (show the changes on above tree)

c) What is the big-oh notation for inserting a new item in the heap?

Now let us consider the delMin operation that removes and returns the minimum item. Recall the General Idea of delMin():

- remember the minimum value so it can be returned later (easy to find - at index 1)
- copy the last item in the list to the root, delete it from the right end, decrement size
- restore the heap-order property by repeatedly swapping this item with its smallest child until it percolates down to the correct spot
- return the minimum value

d) What would the above heap look like after delMin? (show the changes on above tree)

Answer the above questions, then raise your hand. Explain your answers to an instructor or TA.
**Part C**: Run the `lab4/timePriorityQueues.py` program that enqueues 20,000 random integers followed by dequeueing all 20,000 integers from various priority queues discussed above. Complete the following timing table from the output of `timePriorityQueues.py`.

<table>
<thead>
<tr>
<th>Priority Queue Implementation</th>
<th>Execution Time in Seconds</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Enqueuing 20,000 Random ints</td>
</tr>
<tr>
<td>Unssorted Python list</td>
<td></td>
</tr>
<tr>
<td>Sorted Python list in descending order</td>
<td></td>
</tr>
<tr>
<td>“Reverse” sorted Python list in ascending order</td>
<td></td>
</tr>
<tr>
<td>Binary heap stored in a Python list</td>
<td></td>
</tr>
</tbody>
</table>

b) Why does it take more time to enqueue 20,000 items in the “unsorted” Python list version than dequeue 20,000 in the sorted Python list version?

c) Why does it take more time to dequeue 20,000 items in the heap version than enqueue 20,000 in the heap version?

d) Why is the heap implementation of the priority queue considered “better” than the other three?

After you have answered the above questions, raise your hand and explain your answers.

If you have extra time, complete previous labs (1 to 3) or work on homework!
1. The textbook’s unordered list ADT uses a singly-linked list implementation. I added the _size and _tail attributes:

```
UnorderedList Object
  _head
  _size 4
  _tail

data   next   data   next   data   next   data   next
'w'    'a'    'y'    'c'
```

a) The `search(targetItem)` method searches for `targetItem` in the list. It returns `True` if `targetItem` is in the list; otherwise it returns `False`. Complete the `search(targetItem)` method code:

```python
class UnorderedList:
    def search(self, targetItem):
```

b) The textbook’s unordered list ADT does not allow duplicate items, so operations `add(item)`, `append(item)`, and `insert(pos, item)` would have what precondition?

c) Complete the `append(item)` method including a check of it’s precondition(s):

```python
def append(self, item):
```

d) Why do you suppose I added a _tail attribute?
Data Structures (CS 1520)  Lecture 8  Name:

e) The textbook's `remove(item)` and `index(item)` operations "Assume the item is present in the list." Thus, they would have a precondition like "Item is in the list." When writing a program using an UnorderedList object (say `myGroceryList = UnorderedList()`), how would the programmer check if the precondition is satisfied?

```python
itemToRemove = input("Enter the item to remove from the Grocery list: ")
if
    myGroceryList.remove(itemToRemove)
```

f) The `remove(item)` and `index(item)` methods both need to look for the item. What is inefficient in this whole process?

g) Modify the `search(targetItem)` method code in (a) to set additional data attributes to aid the implementation of the `remove(item)` and `index(item)` methods.

h) Write the `index(item)` method including a check of its precondition(s).

```python
def index(self, item):
```

i) Write the `remove(item)` method including a check of its precondition(s).

```python
def remove(self, item):
```
1. The textbook's **ordered list** ADT uses a singly-linked list implementation. I added the `_size`, `_tail`, `_current`, `_previous`, and `_currentIndex` attributes:

![Diagram of OrderedList Object]

The `search(targetItem)` method searches for `targetItem` in the list. It returns `True` if `targetItem` is in the list; otherwise it returns `False`. Additionally, it has the side-effects of setting `_current`, `_previous`, and `_currentIndex`. The complete `search(targetItem)` method code for the `OrderedList` is:

```python
class OrderedList:
    def search(self, targetItem):
        if self._current != None and self._current.getData() == targetItem:
            return True

        self._previous = None
        self._current = self._head
        self._currentIndex = 0
        while self._current != None:
            if self._current.getData() == targetItem:
                return True
            elif self._current.getData() > targetItem:
                return False
            else: # inch-worm down list
                self._previous = self._current
                self._current = self._current.getNext()
                self._currentIndex += 1
        return False
```

a) What's the purpose of the "elif self._current.getData() > targetItem:" check?

Consider the `add(item)` method with the precondition: `item` is not in the list.

b) Write the precondition check at the start of the `add(item)` method.

c) Suppose you are adding the item value of 's'. Update the above picture for this "normal" case, and number the steps in the drawing.

d) What special cases need to be considered for the `add` method?
2. A recursive function is one that calls itself. Complete the recursive code for the `countDown` function that is passed a starting value and proceeds to count down to zero and prints "Blast Off!!!".

Hint: The `countDown` function, like most recursive functions, solves a problem by splitting the problem into one or more simpler problems of the same type. For example, `countDown(10)` prints the first value (i.e., 10) and then solves the simpler problem of counting down from 9. To prevent "infinite recursion", if-statement(s) are used to check for trivial base case(s) of the problem that can be solved without recursion. Here, when we reach a `countDown(0)` problem we can just print "Blast Off!!!".

```python
*** File: countDown.py ***

def main():
    start = eval(input("Enter count down start: "))
    print("nCount Down:")
    countDown(start)

def countDown(count):

main()
```

Program Output:

```
Enter count down start: 10
Count Down:
10
9
8
7
6
5
4
3
2
1
Blast Off!!!
```

(a) Trace the function call `countDown(5)` on paper by drawing the run-time stack and showing the output.

(b) What do you think will happen if your call `countDown(-1)`?

(c) Why is there a limit on the depth of recursion?
3. The non-recursive \_\_str\_ method for \texttt{OrderedList} object below would return: \texttt{"(head) a c m (tail)"}

```python
def \_\_str\_ (self):
    resultStr = \"(head) \"
    current = self.\_head
    while current != None:
        resultStr += str(current.getData()) + \" \"
        current = current.getNext()
    return resultStr + \"(tail)\"
```

We can thing of building the string for the list as \texttt{\"a \" + (string for the rest of the list)\"}

(a) Complete the recursive \texttt{strHelper} function in the \_\_str\_ method for our \texttt{OrderedList} class.

```python

def \_\_str\_ (self):
    \"\"\" Returns a string representation of the list with a space between each item. \"\"\"

    def strHelper(current):
        \# Start of \_\_str\_ method execution
        return \"(head) \" + strHelper(self.\_head) + \"(tail)\"
```

4. Some mathematical concepts are defining by recursive definitions. One example is the Fibonacci series:

\begin{align*}
0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, ...
\end{align*}

After the second number, each number in the series is the sum of the two previous numbers. The Fibonacci series can be defined recursively as:

\begin{align*}
def \text{fib} (n):
    \text{Fib}_0 &= 0 \\
    \text{Fib}_1 &= 1 \\
    \text{Fib}_n &= \text{Fib}_{n-1} + \text{Fib}_{n-2} \text{ for } n \geq 2.
\end{align*}

(a) Complete the recursive function:

(b) Draw the \texttt{call tree} for \texttt{fib(5)}.
c) On my office computer, the call to fib(40) takes 22 seconds, the call to fib(41) takes 35 seconds, and the call to fib(42) takes 56 seconds. How long would you expect fib(43) to take?

d) How long would you guess calculating fib(100) would take on my office computer?

e) Why do you suppose this recursive fib function is so slow?

f) What is the computational complexity? $O(\quad)$

g) How might we speed up the calculation of the Fibonacci series?

5. A VERY POWERFUL concept in Computer Science is dynamic programming. Dynamic programming solutions eliminate the redundancy of divide-and-conquer algorithms by calculating the solutions to smaller problems first, storing their answers, and looking up their answers if later needed instead of recalculating them.

We can use a list to store the answers to smaller problems of the Fibonacci sequence.

To transform from the recursive view of the problem to the dynamic programming solution you can do the following steps:

1) Store the solution to smallest problems (i.e., the base cases) in a list
2) Loop (no recursion) from the base cases up to the biggest problem of interest. On each iteration of the loop we:
   - solve the next bigger problem by looking up the solution to previously solved smaller problem(s)
   - store the solution to this next bigger problem for later usage so we never have to recalculate it

a) Complete the dynamic programming code:

```python
def fib(n):
    """Dynamic programming solution to find the nth number in the Fibonacci seq."""

    # List to hold the solutions to the smaller problems
    fibonacci = []

    # Step 1: Store base case solutions
    fibonacci.append( 0 )
    fibonacci.append( 1 )

    # Step 2: Loop from base cases to biggest problem of interest for position in range(
    fibonacci.append(
    )

    # return nth number in the Fibonacci sequence
    return
```

Running the above code to calculate fib(100) would only take a fraction of a second.

b) One tradeoff of simple dynamic programming implementations is that they can require more memory since we store solutions to all smaller problems. Often, we can reduce the amount of storage needed if the next larger problem (and all the larger problems) don’t really need the solution to the really small problems, but just the larger of the smaller problems. In fibonacci when calculating the next value in the sequence how many of the previous solutions are needed?
Data Structures (CS 1520)  Lab 5  Name:__________

Objective: To understand recursion by writing simple recursive solutions.

To start the lab: Download and unzip the file at: www.cs.uni.edu/~fienup/cs1520s19/labs/lab5.zip

Part A: Recall: We modified the textbook's ordered list ADT that uses a singly-linked list implementation by adding the _size, _tail, _current, _previous, and _currentIndex attributes:

```python
# NON-RECURSIVE CODE WE ARE REPLACING
def search(self, targetItem):
    if self._current != None and \
        self._current.getData() == targetItem:
        return True
    self._current = self._head
    self._currentIndex = 0
    while self._current != None:
        if self._current.getData() == targetItem:
            return True
        elif self._current.getData() > targetItem:
            return False
        else: # inch-worm down list
            self._previous = self._current
            self._current = self._current.getNext()
            self._currentIndex += 1
    return False
```

def searchHelper():
    # Recursive helper function that moves down the linked list.
    # It has no parameters, but uses self._current, self._previous, self._currentIndex."
    # ADD CODE HERE

    # START OF SEARCH - DO NOT MODIFY BELOW CODE
    if self._current != None and \
        self._current.getData() == targetItem:
        return True
    self._previous = None
    self._current = self._head
    self._currentIndex = 0
    return searchHelper()  # Returns the result of searchHelper

a) What are the base case(s) for the searchHelper that halt the while-loop of the non-recursive search code?

b) What are the recursive case(s) for the searchHelper that replaces the while-loop of the non-recursive search code?

c) Complete the recursive searchHelper function in the search method of our OrderedList class in ordered_linked_list.py. Test it with the listTester.py program.

Raise your hand when done. Demonstrate and explain your code to an instructor.
Part B: Recall that Lecture 7 and Section 6.6 discussed a very “non-intuitive”, but powerful list/array-based approach to implement a priority queue, call a binary heap. The list/array is used to store a complete binary tree (a full tree with any additional leaves as far left as possible) with the items being arranged by heap-order property, i.e., each node is ≤ either of its children. An example of a min heap “viewed” an a complete binary tree would be:

Recall the General Idea of insert(newItem):

- append newItem to the end of the list (easy to do, but violates heap-order property)
- restore the heap-order property by repeatedly swapping the newItem with its parent until it percolates up to the correct spot

Recall the General Idea of delMin():

- remember the minimum value so it can be returned later (easy to find - at index 1)
- copy the last item in the list to the root, delete it from the right end, decrement size
- restore the heap-order property by repeatedly swapping this item with its smallest child until it percolates down to the correct spot
- return the minimum value

Originally, we used iteration (i.e., a loop) to percolate up (see percUp) and percolate down (see percDown) the tree. (textbook code below)

```python
## NON-RECURSIVE CODE WE ARE REPLACING
def percUp(self,i):
    while i // 2 > 0:
        if self.heapList[i] < self.heapList[i // 2]:
            tmp = self.heapList[i // 2]
            self.heapList[i // 2] = self.heapList[i]
            self.heapList[i] = tmp
        i = i // 2

## RECURSIVE CODE
def percDown(self,i):
    while (i * 2) <= self.currentSize:
        mc = self.minChild(i)
        if self.heapList[i] > self.heapList[mc]:
            tmp = self.heapList[i]
            self.heapList[i] = self.heapList[mc]
            self.heapList[mc] = tmp
        i = mc
```

For part B, I want you to complete the recursive percUpRec and recursive percDownRec methods in binHeap.py. Run the binHeap.py file to test your code.

Raise your hand when done. Demonstrate and explain your code to an instructor.

(If you have extra time, work on previous labs or homeworks!)
1. Consider the coin-change problem: Given a set of coin types and an amount of change to be returned, determine the fewest number of coins for this amount of change.

a) What "greedy" algorithm would you use to solve this problem with US coin types of \(\{1, 5, 10, 25, 50\}\) and a change amount of 29-cents?

b) Do you get the correct solution if you use this algorithm for coin types of \(\{1, 5, 10, 12, 25, 50\}\) and a change amount of 29-cents?

2. One way to solve this problem in general is to use a divide-and-conquer algorithm. Recall the idea of **Divide-and-Conquer** algorithms.

Solve a problem by:
- dividing it into smaller problem(s) of the same kind
- solving the smaller problem(s) recursively
- use the solution(s) to the smaller problem(s) to solve the original problem

a) For the coin-change problem, what determines the size of the problem?

b) How could we divide the coin-change problem for 29-cents into smaller problems?

c) If we knew the solution to these smaller problems, how would be able to solve the original problem?
3. After we give back the first coin, which smaller amounts of change do we have?

![Diagram: Original Problem with 29 cents, 1-cent coin, 5-cent coin, 10-cent coin, 25-cent coin, 50-cent coin connections to possible smaller problems]

4. If we knew the fewest number of coins needed for each possible smaller problem, then how could determine the fewest number of coins needed for the original problem?

5. Complete a recursive relationship for the fewest number of coins.

\[
\text{FewestCoins}(\text{change}) = \begin{cases} 
\min( \text{FewestCoins}( \text{coin} \in \text{CoinSet and coin} \leq \text{change} ) ) + 1 & \text{if change} \not\in \text{CoinSet} \\
1 & \text{if change} \in \text{CoinSet}
\end{cases}
\]

6. Complete a couple levels of the recursion tree for 29-cents change using the set of coins \( \{ 1, 5, 10, 25, 50 \} \).
Test 1 will be Thursday February 21st in class. It will be closed-book and notes, except for one 8.5” x 11” sheet of paper (you can use front and back) containing any notes that you want AND the Python Summary handout. The test will cover the following topics (and maybe more).

**Chapter 1. Introduction**
Python control structures: if, while, for,
Python built-in data structures: list, dictionary, string
Preconditions, postconditions, and raising exceptions to enforce the precondition
Defining classes (e.g., Die) including inheritance, overriding methods

**Chapter 2. Algorithm Analysis**
Machine dependent measures of performance: program running time and instruction count
Machine independent measures of performance: big-oh, orders of complexity: constant O(1), logarithmic O(log n), linear O(n), “n log n”/log linear O(n log n), quadratic O(n^2), cubic O(n^3), exponential O(2^n)
Complexity analysis of an algorithm to determine its big-oh notation
Implementation of Python lists as an array of object references with implications on operation big-oh (e.g., pop() is O(1) while pop(0) is O(n), etc.)

**Chapter 3. Basic Data Structures**
General concept of a stack: LIFO, top and bottom
Stack Operations: pop, push, peek, size, isEmpty, __str__
Stack Implementations: Python list to store stack items and linked list of Nodes to store stack items including big-oh of operations
Stack Applications: general idea of using a stack to do parentheses matching and palindrome checking

General concept of a queue: FIFO, front and rear
Queue Operations: enqueue, dequeue, peek, size, isEmpty, __str__
Queue Implementations: Python list to store queue items and linked list of Nodes to store stack items including big-oh of operations

General concept of a deque: double ended queue, front and rear
Deque Operations: addFront, addRear, removeFront, removeRear, size, isEmpty, __str__
Deque Implementations: Python list to store deque items, singly-linked list of Nodes to store deque items, and doubly-linked list of Nodes (e.g, Node2Way) including big-oh of operations
Deque Applications: general idea of using deque to do palindrome checking

General concept of a list: head, tail, index
Categories of List operations: index-based, content-based, cursor-based
Unordered List operations and implementation with a singly-linked list of Nodes including big-oh of operations
Ordered List operations and implementation with a singly-linked list of Nodes including big-oh of operations

**Section 6.6. Priority Queue with Binary Heaps**
General concept of a priority queue: remove highest priority next
Priority Queue Operations: enqueue, dequeue, peek, size, isEmpty, __str__
Priority Queue Implementations: Python list unordered, Python list ordered by priority, Binary Heap including big-oh of operations
Binary Heap implementation: insert, findMin, delMin, isEmpty, size, buildHeap including big-oh of operations, and binary-tree diagrams after insert and delMin operations
Question 1. (4 points) Consider the following Python code.

```python
for j in range(n):
    i = 1
    while i < n:
        print(i, j)
        i = i * 2
```

What is the big-oh notation $O(\cdot)$ for this code segment in terms of $n$?

Question 2. (4 points) Consider the following Python code.

```python
for i in range(n):
    k = n
    while k > 1:
        k = k // 2
        print(k)
    for j in range(n):
        print(i, j)
```

What is the big-oh notation $O(\cdot)$ for this code segment in terms of $n$?

Question 3. (4 points) Consider the following Python code.

```python
def main(n):
    for i in range(n):
        doSomething(n)
def doSomething(n):
    for j in range(n*n):
        doMore(n)
def doMore(n):
    for k in range(n*n):
        print(k)
main(n)
```

What is the big-oh notation $O(\cdot)$ for this code segment in terms of $n$?

Question 4. (6 points) Suppose a $O(n^4)$ algorithm takes 10 second when $n = 1000$. How long would the algorithm run when $n = 10,000$?

Question 5. (8 points) Why should a method/function having a "precondition" raise an exception if the precondition is violated?
Question 6. A FIFO queue allows adding a new item at the rear using an enqueue operation, and removing an item from the front using a dequeue operation. One possible implementation of a queue would be to use a built-in Python list to store the queue items such that
- the **front** item is always stored at index 0,
- the rear item is always at index `len(self.items) - 1` or -1

![Diagram of a queue and a Python list object](image)

a) (6 points) Complete the expected big-oh $O()$, for each Queue operation, assuming the above implementation. Let $n$ be the number of items in the queue.

<table>
<thead>
<tr>
<th>isEmpty</th>
<th>enqueue(item)</th>
<th>dequeue</th>
<th>peek - returns front item without removing it</th>
<th><strong>str</strong></th>
<th>size</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

b) (8 points) Complete the method for the **dequeue** operation, **including the precondition check to raise an exception if it is violated**.

```python
def dequeue(self):
    """Removes and returns the Front item of the Queue
    Precondition: the Queue is not empty.
    Postcondition: Front item is removed from the Queue and returned"
```

c) (8 points) Complete the method for the **__str__** operation,

```python
def __str__(self):
    """Returns a string representation of items from front to rear."
    strResult = "[front) "
```
Question 7. Consider the binary heap approach to implement a priority queue. A Python list is used to store a complete binary tree (a full tree with any additional leaves as far left as possible) with the items being arranged by heap-order property, i.e., each node is ≤ either of its children. An example of a min heap "viewed" as a complete binary tree would be:

```
<table>
<thead>
<tr>
<th>Python List actually used to store heap items</th>
</tr>
</thead>
<tbody>
<tr>
<td>12 23 34 25 60 90 120 44 28 31 96 84 98</td>
</tr>
</tbody>
</table>
```

a) (3 points) For the above heap, the list indexes are indicated in [ ]'s. For a node at index i, what is the index of:
- its left child if it exists:
- its right child if it exists:
- its parent if it exists:

b) (7 points) What would the above heap look like after inserting 40 and then 20 (show the changes on above tree)

Now consider the delMin operation that removes and returns the minimum item.

c) (2 point) What item would delMin remove and return from the above heap?

d) (7 points) What would the heap look like after delMin? (show the changes on tree in the middle of the page)

e) (6 points) Performing 20,000 inserts into an initially empty binary heap takes 0.23 seconds. Now, if we perform 20,000 delMin operations, it takes 0.39 seconds. Explain why 20,000 delMin operations take more time than the 20,000 insert operations?
Fall 2018

Data Structures - Test 1

Name: ___________

Question 8. The `Node2Way` and `Node` classes can be used to dynamically create storage for each new item added to a Deque using a doubly-linked implementation as in:

![DoublyLinkedDeque Object Diagram]

DoublyLinkedDeque Object

- front:
- rear:
- size:

a) (6 points) Complete the big-oh expected $O(\_\_\_\_\_\_\_\_)$ for each DoublyLinkedDeque operation, assuming the above implementation. Let $n$ be the number of items in the DoublyLinkedDeque.

<table>
<thead>
<tr>
<th>Operation</th>
<th>$O(________)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>isEmpty</td>
<td>$________$</td>
</tr>
<tr>
<td>addRear</td>
<td>$________$</td>
</tr>
<tr>
<td>removeRear</td>
<td>$________$</td>
</tr>
<tr>
<td>addFront</td>
<td>$________$</td>
</tr>
<tr>
<td>removeFront</td>
<td>$________$</td>
</tr>
<tr>
<td><strong>str</strong></td>
<td>$________$</td>
</tr>
</tbody>
</table>

b) (16 points) Complete the `addRear` method for the above DoublyLinkedDeque implementation.

```python
class DoublyLinkedDeque(object):
    """ Doubly-linked list based deque implementation."""
    def __init__(self):
        self.size = 0
        self._front = None
        self._rear = None
    def addRear(self, newItem):
        """ Adds the newItem to the rear of the Deque. Precondition: none """
        def _init_(self, initdata):
            self.data = initdata
            self.next = None
        def getData(self):
            return self.data
        def getNext(self):
            return self.next
        def setNext(self, newnext):
            self.next = newnext
class Node:
    def __init__(self, initdata):
        self.data = initdata
        self.next = None
def getData(self):
    return self.data
def getNext(self):
    return self.next
def setNext(self, newnext):
    self.next = newnext
class Node2Way(Node):
    def __init__(self, initdata):
        Node.__init__(self, initdata)
        self.previous = None
    def getPrevious(self):
        return self.previous
def setPrevious(self, newprevious):
    self.previous = newprevious

from node import Node
```

c) (5 points) Why would using singly-linked nodes (i.e., only Node objects with data and next) to implement the Deque lead to poor performance (i.e., cause some Deque operations to have worse big-oh notations)? Justify your answer.