

## 1. In the 0-1 Knapsack problem:

A thief breaks into a jewelry store carrying a knapsack that will break if its weight limit ( $W$ ) is exceeded. The thief wants to maximize the total value in the knapsack without exceeding its weight limit  $W$ .

Consider the following 0-1 Knapsack problem with four items and a knapsack weight limit of  $W=10$  oz.

Item, $i$	Weight, $w_i$	Profit, $p_i$	Profit/Weight
1	4 oz.	\$40	\$10/oz.
2	7 oz.	\$63	\$9/oz.
3	5 oz.	\$25	\$5/oz.
4	3 oz.	\$12	\$4/oz.

To solve a problem using backtracking, you need to answer the following questions:

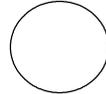
a) What should the state-space tree look like? (i.e., What would the "for each child  $c$ " loop iterate over?)

(Hint: consider alternate state-space tree which might be better for exploring subsets)

**Global Problem-Instance Information**

	1	2	3	4
(weights) w:	4	7	5	3
	1	2	3	4
(profits) p:	40	63	25	12
n:	4	W:	10	

**Starting Node** - original call to start recursive backtracking



b) What state information is needed at each node?

c) What criteria can be used to determine if a parent node (p) is NOT promising?

d) What information is needed by our promising function?

2. Since any subset is potentially the best solution, consider customizing the backtracking optimization template "checknode" (p. 228) for the 0-1 Knapsack problem. Use a single, "global", current-node state which is updated before we go down to the child (via a recursive call) and undone when we backtrack to the parent.

```
checknode( treeNode p ) {  
    treeNode c;  
    if p is better than best solution  
        best = p                # remember as the best solution  
    end if  
    if promising(p) then        # p is "promising" if it could lead to a better solution  
        for each child c of p do # each c represents a possible choice  
            checknode(c)        # follow a branch down the tree  
        end for  
    end if  
} // end checknode
```