

Computer Systems HW #1

Due: Friday, January 25 (3 PM in ITT 305 mailbox or under my office door, ITT 313)

Common MIPS Instructions (and psuedo-instructions)		
Type of Instruction	MIPS Assembly Language	Register Transfer Language Description
Memory Access (Load and Store)	lw \$4, Mem	$\$4 \leftarrow [\text{Mem}]$
	sw \$4, Mem	$\text{Mem} \leftarrow \$4$
	lw \$4, 16(\$3)	$\$4 \leftarrow [\text{Mem at address in } \$3 + 16]$
	sw \$4, 16(\$3)	$[\text{Mem at address in } \$3 + 16] \leftarrow \$4$
Move	move \$4, \$2	$\$4 \leftarrow \2
	li \$4, 100	$\$4 \leftarrow 100$
Load Address	la \$5, mem	$\$4 \leftarrow \text{load address of mem}$
Arithmetic Instruction (reg. operands only)	add \$4, \$2, \$3	$\$4 \leftarrow \$2 + \$3$
	mul \$10, \$12, \$8	$\$10 \leftarrow \$12 * \$8 \quad (\text{32-bit product})$
	sub \$4, \$2, \$3	$\$4 \leftarrow \$2 - \$3$
Arithmetic with Immediates (last operand must be an integer)	addi \$4, \$2, 100	$\$4 \leftarrow \$2 + 100$
	mul \$4, \$2, 100	$\$4 \leftarrow \$2 * 100 \quad (\text{32-bit product})$
Conditional Branch	bgt \$4, \$2, LABEL (bge, blt, ble, beq, bne)	Branch to LABEL if $\$4 > \2
Unconditional Branch	j LABEL	Always Branch to LABEL

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sumPos = 0;
sumNeg = 0;
for i = 0 to length-1 do
    if numbers[i] < 0 then
        sumNeg = sumNeg + numbers[i]
    else
        sumPos = sumPos + numbers[i]
    end if
end for

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1. Write MIPS Assembly Language code for the above algorithm that sums the array's elements.

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.data
numbers: .word 20, 30, 10, 40, 50, 60, 30, 25, 10, 5
length: .word 10
sumPos: .word 0
sumNeg: .word 0

.text
.globl main
main:

```

2. Compare zero-, one-, two-, three-address, and the load & store machines by writing programs to compute

$$X = A * B + C * D;$$

$$Y = (A + B) / (X - C);$$

for each of the five machines. The instructions available for use are as follows:

3 Address	2 Address	1 Address (Accumulator machine)	0 Address (Stack machine)
MOVE (X \leftarrow Y)	MOVE (X \leftarrow Y)	LOAD M	PUSH M
		STORE M	POP M
ADD (X \leftarrow Y + Z)	ADD (X \leftarrow X + Y)	ADD M	ADD
SUB (X \leftarrow Y - Z)	SUB (X \leftarrow X - Y)	SUB M	SUB
MUL (X \leftarrow Y * Z)	MUL (X \leftarrow X * Y)	MUL M	MUL
DIV (X \leftarrow Y / Z)	DIV (X \leftarrow X / Y)	DIV M	DIV
		Notes: “SUB M” performs AC = AC - M “DIV M” performs AC = AC / M	Notes: “SUB” performs POP T POP T2 T3 = T2 - T PUSH T3 “DIV” performs POP T POP T2 T3 = T2 / T PUSH T3

Load/Store Architecture - operands for arithmetic operations must be from/to registers. For example, to perform the high-level statement “Z = X - Y” we need the code:

LOAD R2, X

LOAD R3, Y

SUB R4, R2, R3

STORE R4, Z

3. Assume 8-bit opcodes, 32-bit absolute addressing, 5-bit register numbers, and 32-bit operands. Compute the number of bits needed in programs from question 2 by completing the following table.

	3 Address	2 Address	1 Address	0 Address	Load & Store
Number of bits needed to store the program					
Number of bits of data transferred to and from memory					
Total number of bits read and written while the program executes					