

```
(define 2nd-max  
  (lambda (lon)  
    (second (sort lon >))))
```

Short, sweet, and  $O(n \log n)$

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- 1 Find the largest value in the list.
- 2 Remove that item from the list.
- 3 Find the largest value in what's left.

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- 1 Find the largest value in the list.
- 2 Remove that item from the list.
- 3 Find the largest value in what's left.

```
(define 2nd-max
  (lambda (lon)
    (apply max                ; step 3
           (remove            ; step 2
              (apply max lon) ; step 1
              lon))))
```

A little longer, and only  $O(n)$ .  
But it makes three passes...

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Our argument contains at least two numbers:

(<number> <number> . <list-of-numbers>)

With the usual definition for a list:

<list-of-numbers>

::= ()

| (<number> . <list-of-numbers>)

So we will want an interface procedure.

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If we could write a loop, we might...

- Create two local variables, **largest** and **2nd-largest**.
- Initialize the variables using the first two items in the list.
- **Then** look at each item in the rest of the list to see if it is greater than either of the two variables and, if so, update the variables.

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(6 1 2 -3 9 4 -1 2 8 1 2 4)

largest 6

2nd-largest 1

rest (**2** -3 9 4 -1 2 8 1 2 4)

.

(6 1 2 -3 9 4 -1 2 8 1 2 4)

largest	6	<- replace w/ (max 6 2)
2nd-largest	1	
rest	( <b>2</b> -3 9 4 -1 2 8 1 2 4)	

.

(6 1 2 -3 9 4 -1 2 8 1 2 4)

largest                   6  
2nd-largest           1    <- replace w/ (max 1 2)  
rest                    (2 -3 9 4 -1 2 8 1 2 4)

.

(6 1 2 -3 9 4 -1 2 8 1 2 4)

largest 6

2nd-largest 2

rest (-3 9 4 -1 2 8 1 2 4)

.

IN PYTHON

```
def second_max(lst):  
    largest = max(lst[0], lst[1])  
    second  = min(lst[0], lst[1])  
  
    for n in lst[2:]:  
        larger  = max(largest, n)  
        smaller = min(largest, n)  
  
        largest = larger  
        second  = max(second, smaller)  
  
    return second
```

•

```
(define 2nd-max
  (lambda (lon)
    (2nd-max-tr
     (max (first lon) (second lon))
     (min (first lon) (second lon))
     (rest (rest lon)))))
```

This is order  $O(n)$  and makes only **one** pass!  
But, man, 2nd-max-tr is ugly...

.

Use our interface procedure as inspiration:

```
(define 2nd-max-tr
  (lambda (largest 2nd-largest lon)
    (if (null? lon)
        2nd-largest
        (2nd-max-tr
         ; -- handle (first lon)
         new value of largest
         new value of second
         ; handle (rest lon)
         new value of lon          ))))
```

.

```
(define 2nd-max-tr
  (lambda (largest 2nd-largest lon)
    (if (null? lon)
        2nd-largest
        (2nd-max-tr
         (max largest (first lon))
         (max 2nd-largest
          (min largest (first lon)))
         (rest lon))))))
```

This is order  $O(n)$ , makes only **one** pass,  
and uses only one stack frame.

This is dandy solution, IMHO.

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How might we compare these solutions?

- length of the code
- space used at run-time
- time used at run-time
- time to create the program
- ...
- complexity of the code
- ...
- familiarity

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learned a new language in order to ...

... learn a new way to think about languages

... learn a new style of programming

... learn patterns of recursive programs

now: use all three to ...

... learn how programming languages work

.

## Static Properties of Variables

A property is **static** when its value can be determined by looking at the text of a program.

A property is **dynamic** when the program must be executed in order to determine its value.

Compilers can use static properties of a program to detect errors and to improve program performance.

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## A little language

$\langle \text{exp} \rangle ::= \langle \text{varref} \rangle$   
          |  $(\text{lambda } (\langle \text{var} \rangle) \langle \text{exp} \rangle)$   
          |  $(\langle \text{exp} \rangle \langle \text{exp} \rangle)$

.

free and bound variables

```
int sumOfSquares( int m, int n )  
{  
    // m and n are bound  
    // to formal parameters  
    return m*m + n*n;  
}
```

•

A variable **is bound** or **occurs bound** in an expression if it refers to the formal parameter in the expression.

A variable **is free** or **occurs free** in an expression if it is not bound.

.

$\langle \text{exp} \rangle ::= \langle \text{varref} \rangle$   
          |  $(\text{lambda } (\langle \text{var} \rangle) \langle \text{exp} \rangle)$   
          |  $(\langle \text{exp} \rangle \langle \text{exp} \rangle)$

Free and bound variables in this language:

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This is **not** a combinator:

```
(define sum-of-applications  
  (lambda (f x y)  
    (+ (f x) (f y))))
```

•

Quiz 1

60 points total

=> 54     A

=> 48     B

=> 42     C

=> 36     D

quiz average = 45

"What is my course grade?"

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pair  
-----

list  
-----

type predicate

pair?

list?

access procedures

car  
cdr

first  
rest

constructor

cons

list

.